

Low Latency QoS Design for 5G Solution Juniper Validated Design (JVD)

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5G xHaul Low Latency Class of Service Juniper Validated Design (JVD)

Juniper Networks Validated Designs provide you with a comprehensive, end-to-end blueprint for deploying Juniper solutions in your network. These designs are created by Juniper's expert engineers and tested to ensure they meet your requirements. Using a validated design, you can reduce the risk of costly mistakes, save time and money, and ensure that your network is optimized for maximum performance.

About This Document

This document explains a Juniper Validated Design (JVD) for building and deploying a 5G xHaul solution architecture using the Juniper ACX7000 series, MX series, and PTX series platforms. This JVD extends solutions that are provided previously with [5G CSR Seamless Segment Routing](#) and [5G Fronthaul Class of Service](#) for the focused delivery of differentiated services supporting ultra-low latency workloads. Comprehensive Quality of Service (QoS) in the 5G network architecture is mandatory to ensure reliable and efficient performance across diverse applications and services. QoS mechanisms prioritize traffic, manage bandwidth, and preserve latency budgets, ensuring critical applications receive the necessary resources and maintain overall network performance and user experience. The capabilities and performance of the 5G network become paramount as technological advancements are considered to support autonomous driving, remote surgery, industrial automation, and other mission-critical services.

The functional and performance aspects of Fronthaul services and Class of Service (CoS) operations are thoroughly analyzed to support multi-level traffic priorities with low latency queuing. After conducting rigorous testing, Juniper ACX7024, ACX7100, and ACX7509 provide scalable and flexible platforms for implementing CoS in 5G Fronthaul networks. These routers offer robust performance, high port density, and advanced traffic management capabilities, allowing seamless expansion and adaptation to evolving network requirements.

This JVD is based on 5G reference architectures. However, you can apply the CoS modeling, best practices, and performance results to many implementations.

Solution Benefits

5G radio access networks (RAN) introduce new requirements for the mobile backhaul (MBH) network infrastructure like number of nodes in the network, performance, and feature richness. This leads to growing network complexity. Juniper provides an end-to-end solution for the 5G xHaul network infrastructure, designed to support 4G MBH and 5G network infrastructure over the same physical network. This approach allows operators to smoothly transition from 4G to 5G without disrupting their existing services. The necessary changes can be gradually introduced to accommodate the evolving requirements for increased bandwidth, lower latency, enhanced synchronization, network slicing, and higher scale.

The 5G solution architecture requires robust QoS performance to deliver differentiated services and preserve low latency for critical traffic types. The QoS network architecture enables the following key benefits:

- Tailoring traffic priorities based on the application requirements, such as enhanced mobile broadband (eMBB), ultra-reliable low-latency communication (URLLC), and massive machine-type communications (mMTC)
- Ensuring network operators can design appropriate bandwidth, latency, and reliability parameters across different traffic classes
- Optimizing network resources to ensure efficient bandwidth allocation and traffic prioritization based on application requirements
- Enhancing user experience with consistent and reliable quality service across diverse applications like streaming, gaming, and remote work
- Implementing LLQ to reduce transmission delay.
- Providing faster response time for crucial real-time applications by assigning preferential treatment over low-priority traffic
- Using highly reliable networks to minimize the risk of data loss and transmission errors, ensuring stable and predictable communication for critical services
- Implementing traffic prioritization at multiple levels ensures urgent services, such as emergency communications and public safety applications receive immediate attention and appropriate resources

These benefits collectively enable 5G networks to meet the diverse demands of modern applications and services, ensuring resilient, efficient, and flexible network performance.

This JVD examines CoS operations and performance requirements needed to ensure the integrity of critical 5G Fronthaul traffic flows between O-RU to O-DU emulated devices (see **Figure 1**), as facilitated

by ACX7024 (CSR), ACX7100-32C (HSR) and ACX7509 (HSR). Additional validation includes MX304 as the Services Edge with PTX10001-36MR as Core for supporting end-to-end MBH traffic flows.

To facilitate the JVD objectives, a comprehensive CoS network model is created to support end-to-end 5G xHaul multiservice traffic types. The CoS model aligns with O-RAN multiple Priority Queue structure, establishing three main components: Low Latency Queues, Shaped Priority Queues, and Weighted Fair Queues. This model is adjustable beyond the featured use cases to meet additional customer requirements.

Juniper's ACX7000 series platforms are specifically designed for use in the 5G Fronthaul segment. They provide necessary connectivity and routing capabilities at the cell sites and hub sites to enable seamless communication between the RAN and core network. The ACX Series routers offer advanced features tailored for MBH and xHaul applications, designed to handle the distinct challenges of high-volume traffic, low latency, and strict QoS requirements associated with 4G and 5G deployments.

Through rigorous testing, KPIs are observed and shared to ensure that the provided solution is compliant with stringent fronthaul SLAs. The featured Juniper Networks fronthaul platforms, including ACX7024, ACX7100-48L, ACX7100-32C, and ACX7509, are excellent choices for access and aggregation purposes. These platforms deliver enhanced performance, low latency, and high bandwidth with the necessary scale and advanced feature set.

Use Case and Reference Architecture

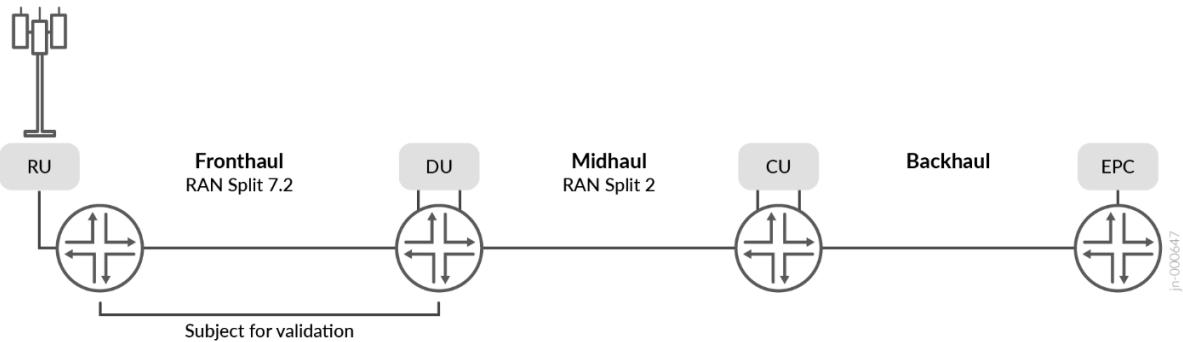
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5G O-RAN is an open and disaggregated radio access network architecture that enables interoperability, flexibility, and innovation among vendors and operators. The fronthaul is the most demanding segment of the xHaul architecture, requiring high performance and functionality to ensure delivery and preservation of ultra-low latency workloads introduced in the 5G network. One of the key features of 5G O-RAN is the support for multiple classes of service that provide differentiated QoS and network slicing for various use cases and applications.

The solution is based on the 5G xHaul network and reference architecture defined by O-RAN Alliance Working Group 9 in the [O-RAN Xhaul Packet Switched Architectures and Solutions](#) technical specification [O-RAN.WG9.XPSAAS-R003-v08.00], which establishes the service aspects and requirements for the xHaul plane.

Figure 1: xHaul Network



The xHaul consists of the Fronthaul (FH), Midhaul (MH), and Backhaul (BH) segments that connect to the O-RAN radio units (O-RUs), distributed units (O-DUs), and centralized units (O-CUs), respectively.

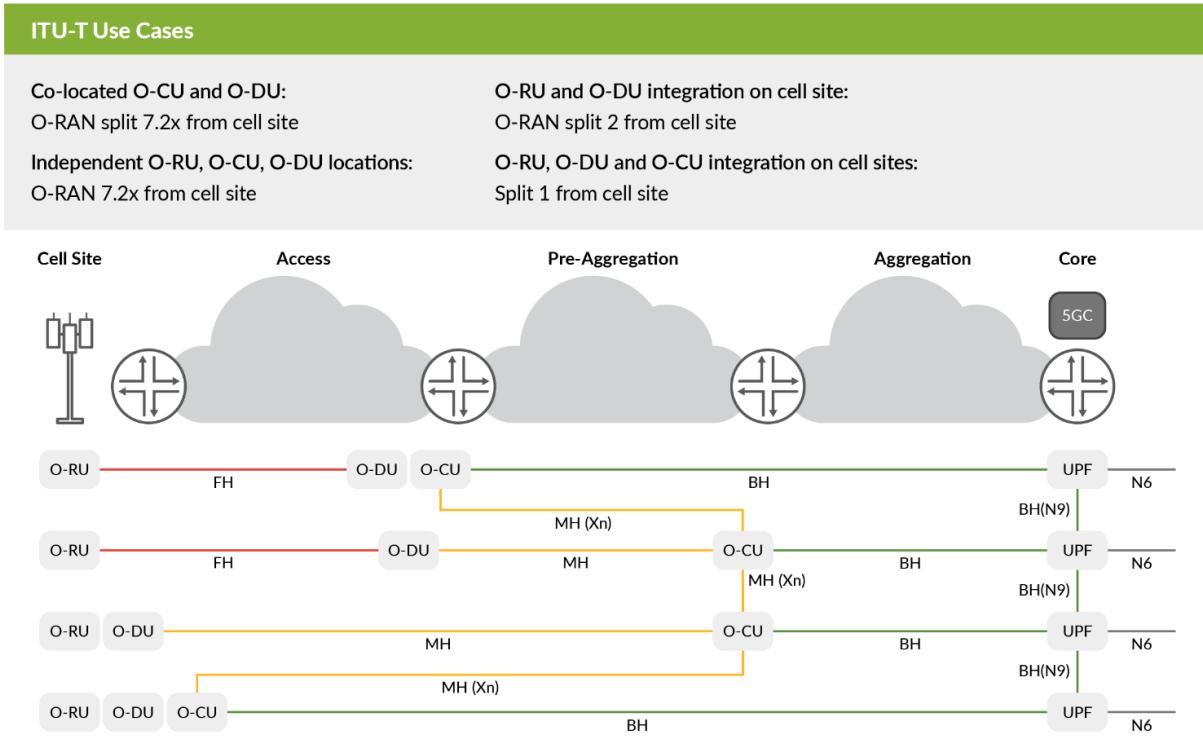
The fronthaul network segment is leveraged to enable Layer 2 connectivity between O-RU and O-DU, also referred to as RU and DU in the diagram ([Figure 1 on page 4](#)) of the RAN for control, data (eCPRI), management traffic flows, and to provide time and frequency synchronization between RAN elements of the fronthaul network.

The advancement of the RAN involves different architectures for 4G, including distributed, centralized, and virtualized infrastructures, which need to coexist with the 5G disaggregated O-RAN. These diverse ecosystems provide flexibility for the placement of components such as O-DU and O-CU.

NOTE: This JVD does not cover all possible scenarios. However, it closely aligns with O-RAN split 7.2x, where the O-RU connects to the CSR and the O-DU is co-located with O-CU within the HSR infrastructure. Additional insertion points can be implemented to support disaggregation between the midhaul and backhaul segments by extending appropriate services.

[Figure 2 on page 5](#) summarizes the deployment scenarios for the RAN according to the ITU-T for simultaneous support of 4G and 5G as proposed by the O-RAN Alliance.

Figure 2: RAN Deployment Scenarios for Simultaneous Support of 4G and 5G



5G infrastructure must support stringent latency budgets between O-RU and O-DU. To achieve this goal, the number of Transport Network Elements (TNEs) facilitating the fronthaul segment is kept to a minimum, in most cases limited to one or two hops. For this reason, spine-and-leaf topologies are ideal. However, ring architectures might be leveraged if the given latency budget can be achieved. For transport requirements, O-RAN WG9 technical specification (O-RAN.WG9.XTRP-REQ-v01.00) provides guidelines for a maximum of one-way latency budget of $100\mu\text{s}$ for the fronthaul. This budget includes latency incurred by the fiber run ($\sim 4.9\mu\text{s}/\text{km}$) and transit nodes. The objective for this JVD is to deliver a solution with latency performance within the acceptable range from O-RU to O-DU, approximately $\leq 10\mu\text{s}$ per node. The ACX7000 series platforms typically support $\leq 6\mu\text{s}$ for low latency workloads, including under network congestion scenarios.

To meet Service Level Objectives (SLOs), 5G transport network QoS must be capable of supporting the spectrum of fronthaul traffic characteristics and mobile backhaul service applications. The referenced CoS scheduling model is constructed to support service types defined by 3GPP, as shown in the following table. These service hierarchies might be realized as part of network slicing architectures, as explained in the O-RAN WG1 technical specification (O-RAN.WG1.Slicing-Architecture-R003-v11.00). While network slicing is outside the scope of this validation, the underlying service attributes are important considerations in developing a viable CoS model.

Table 1: 3GPP Scheduling Model

| Service Type | Characteristics |
|--------------|--|
| eMBB | Slice suitable for the handling of 5G enhanced Mobile Broadband. |
| URLLC | Slice suitable for the handling of ultra-reliable low latency communications |
| V2X | Slice suitable for the handling of V2X services. |
| MIoT / mMTC | Slice suitable for the handling of massive IoT and/or massive machine type communication |

The flow-based 5G QoS model provides more granular control and flexibility compared to the bearer-based 4G LTE QCI model. Instead of assigning a single QoS value to an entire bearer, 5G allows differentiated flows within the same bearer to deliver individual QoS requirements. This enables more fine-grained QoS management and allows different applications and services to receive tailored treatment based on their specific needs. It also enables the network to adjust QoS parameters for individual flows dynamically, optimizing resource allocation and providing a better user experience.

When transitioning from the 4G LTE QCI model to the 5G QoS Identifier (5QI) model, there exist some similarities in defining and handling the traffic. However, in 5G, new categories are introduced specifically for flows that require very low delay and high bandwidth. These are crucial for user-plane and control-plane traffic streams between O-RU and O-DU in the 5G Fronthaul segment. The O-RAN architecture separates user-plane and control-plane traffic to enhance the efficiency and performance of 5G networks. The user-plane traffic includes the transmission of high-speed user data with low latency, while the control-plane traffic manages the signaling and coordination necessary to support the UPF. Both traffic streams are transported using the eCPRI protocol over the Open Fronthaul (OFH) interface, ensuring a standardized and efficient communication pathway between the O-RU and O-DU.

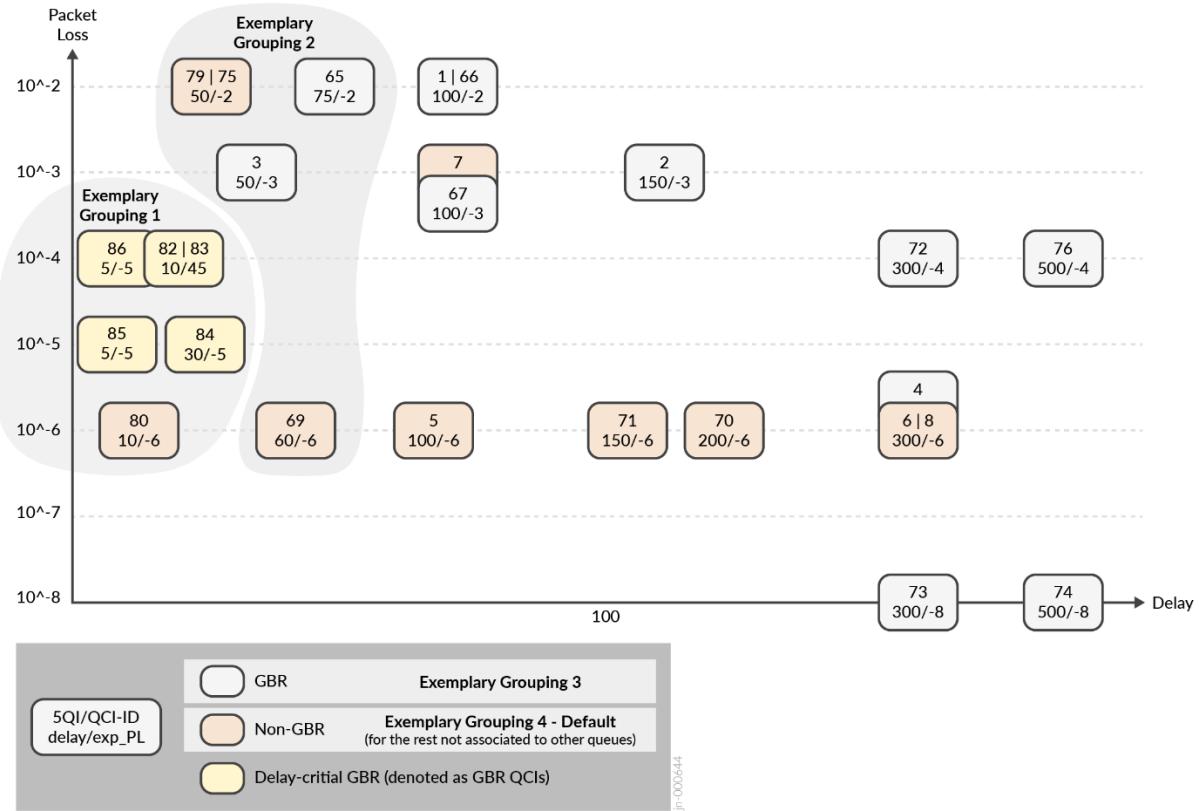
5G incorporates several new delay-critical Guaranteed Bit Rate (GBR) flow categories to meet the strict requirements of certain 5G applications that demand ultra-reliable low latency communications (URLLC). The proper categorization with performance capabilities of the network becomes paramount with technological advancements to support mission-critical services like autonomous driving, remote surgery, and industrial automation.

With high bandwidth and extremely low latency characteristics, 5G fronthaul access devices must assign the highest priority to the eCPRI-based flows supporting user and control traffic streams to ensure optimal performance and reliability. This means that the Service Providers might need to rethink their existing CoS designs, where the highest priority queues are typically reserved for voice or video traffic.

3GPP defines 5QI to QoS characteristics mapping, which can be referenced to build an appropriate CoS model. Resources are classified as delay-critical GBR, GBR, or non-GBR. For more information on 3GPP 5QI models, see details in [ETSI TS-23.501](#).

The O-RAN 5QI/QCI Exemplary Grouping model is shown in [Figure 3 on page 7](#) and defined in O-RAN [O-RAN.WG9.XPSAAS-R003-v08.00], groups common QCI and 5CI QoS flow characteristics into four exemplary groups based on delay budget.

Figure 3: O-RAN Exemplary 5QI Grouping



QoS schemas vary between mobile operators; the JVD does not attempt to qualify a specific design as the recommendation. The main goal is to substantiate deterministic behaviors based on critical and noncritical traffic flows across a range of services delivered by the xHaul network. The CoS solution should further prove capable of supporting the Exemplary model explained here, along with the Time Sensitive Network (TSN) profile and transport CoS characteristics.

Time Sensitive Networking (TSN) Fronthaul Profiles

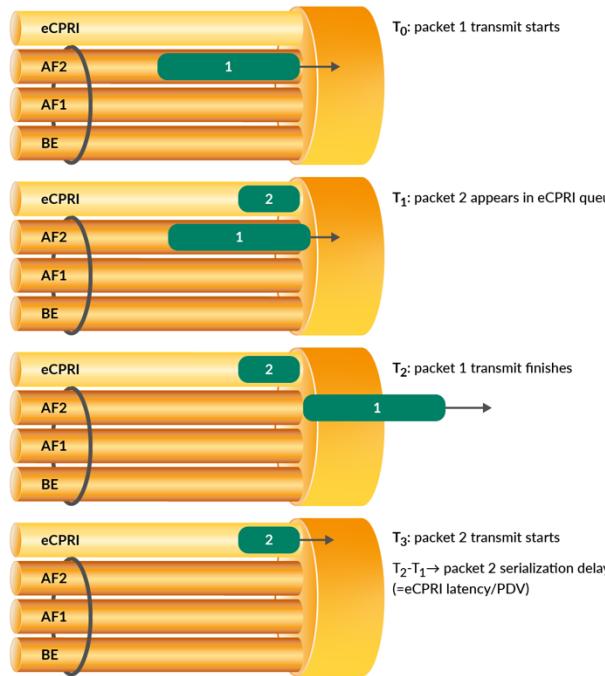
The IEEE [802.1CM](#) standard formalizes Time-Sensitive Networking (TSN) for fronthaul and ensures that fronthaul networks in 5G and advanced mobile networks meet stringent latency, jitter, and reliability requirements. It applies TSN principles, like precise synchronization and traffic shaping, to maintain low-latency and low-jitter communication between remote radio heads and baseband units. This standard

helps mobile operators enhance network performance, support new high-precision applications, and ensure scalable and flexible infrastructure for future advancements. The standard applies to both CPRI and eCPRI traffic flows.

The JVD goals align with TSN Profile A, outlined by 802.1CM section 8.1, which follows a strict-priority model where the user plane is assigned a high-priority class while control and management plane data are assigned a low-priority class. Fronthaul traffic is always prioritized over non-fronthaul traffic, with a maximum frame size recommendation of 2000 octets for both.

O-RAN defines TSN Profile A within its architecture to ensure deterministic communication in the fronthaul network, connecting remote radio units (RRUs) with centralized baseband units (BBUs). Profile A is recommended for ultra-low latency and jitter, with precise synchronization and traffic shaping, suitable for the most demanding 5G applications. [Figure 4 on page 8](#), defined by the O-RAN WG9, illustrates the packet behavior in TSN Profile A.

Figure 4: TSN Profile A



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802.1CM additionally defines TSN Profile B with express frame preemption. O-RAN.WG9.XPSAAS-R003-v08.00 illustrates the minimal benefit for port speeds >1Gbps. ["Table 2 " on page 9](#) outlines queuing delay differences based on port speed and calculates associated fiber reach. Each profile includes a reference frame size influencing queuing delay. The frame size with Profile B is smaller due to preemption, but only when frames are larger than 155 bytes. Frames below 155 bytes are not preempted.

Table 2: Queuing Delay of Non-Fronthaul Data

| Port Speed | TSN Profile A: Queuing delay of 2,020 bytes | TSN Profile B: Queuing delay of 155 bytes | Difference (μs) | Difference (Fiber) |
|------------|--|--|--------------------|-----------------------|
| 1 Gbps | 16.160 μs | 1.240 μs | 14.920 μs | 3,045 m |
| 10 Gbps | 1.616 μs | 0.124 μs | 1.492 μs | 304 m |
| 25 Gbps | 0.646 μs | 0.050 μs | 0.596 μs | 122 m |
| 50 Gbps | 0.323 μs | 0.025 μs | 0.298 μs | 61 m |
| 100 Gbps | 0.162 μs | 0.012 μs | 0.149 μs | 30 m |
| 200 Gbps | 0.081 μs | 0.006 μs | 0.075 μs | 15 m |
| 400 Gbps | 0.040 μs | 0.003 μs | 0.037 μs | 8 m |

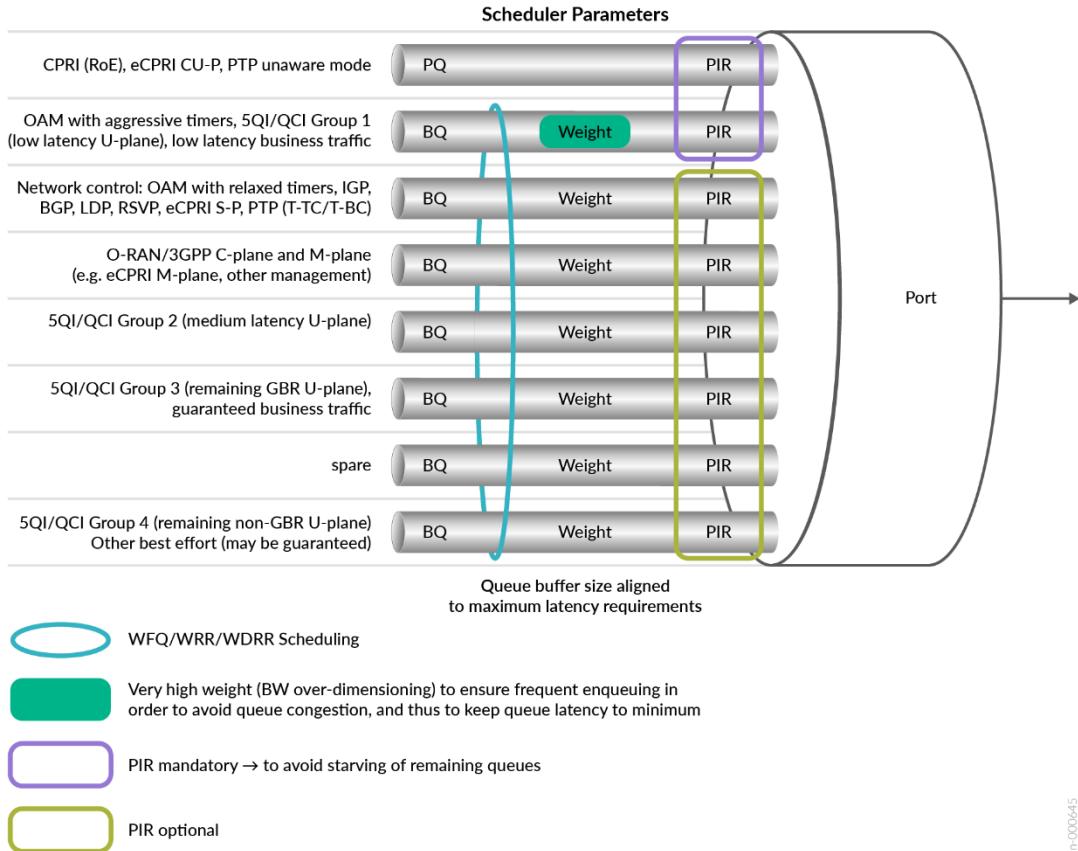
From 10Gbps onwards, the benefits of frame preemption in Profile B diminish to the point of nominal return. The recommended Profile A is leveraged in the JVD solution and includes multiple priority queues, with each having preemptive capability (between queues). The LLQ supports dequeuing prioritization over lower priority frames to ensure delivery of delay-sensitive traffic.

O-RAN Priority Queuing Models

O-RAN/3GPP proposes two common QoS profiles supporting a minimum of 6 queues to accommodate transport network requirements. These profiles include a single Priority Queue (PQ) or multiple Priority Queues.

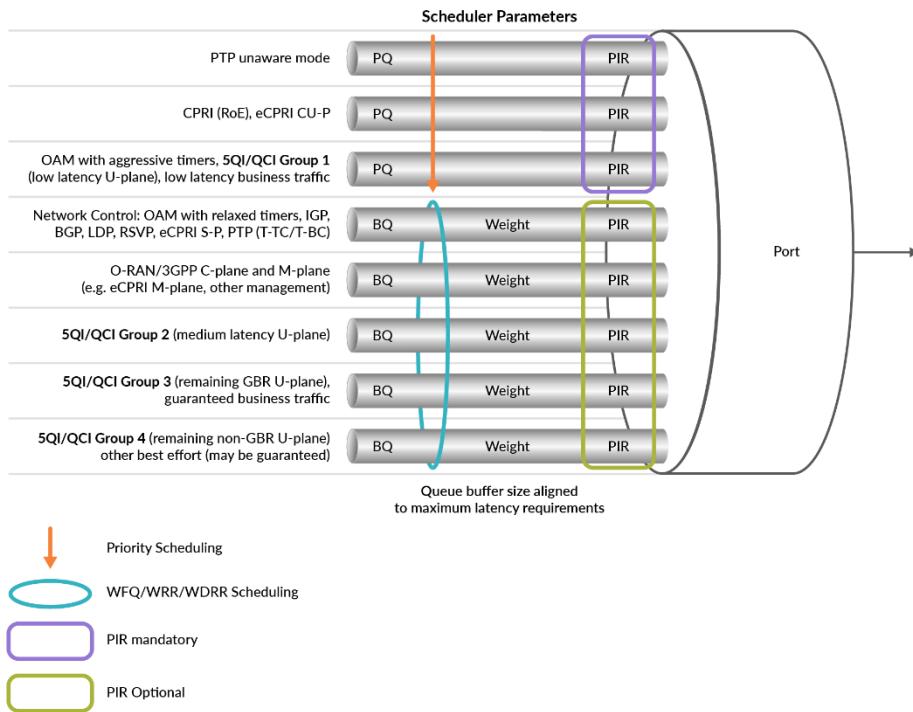
In the single PQ model ([Figure 5 on page 10](#)), the priority queue is defined for handling ultra-low latency flows such as PTP and eCPRI. This queue is serviced ahead of all other queues. Lower priority queues are then serviced as weighted fair queuing (WFQ).

Figure 5: Single Priority Queue



In the second O-RAN transport core QoS model (Figure 6 on page 11), a hierarchy supporting multiple priority queues is utilized with the ability to dequeue in priority order. Priority queues must be rate-limited to prevent starving queues of lower priorities.

Figure 6: Multiple Priority Queue



The previous JVDs, 5G validated designs for [5G CSR Seamless Segment Routing](#) and [5G Fronthaul Class of Service](#) leveraged the Single Priority Queue model. The strict-high queue is reserved for the most critical traffic flows, such as eCPRI between O-RU and O-DU and is shaped to prevent starving lower priority queues.

As per Junos OS Evolved Release 23.3R1, the ACX7000 family supports six priorities: low-latency, strict-high, high, medium-high, medium-low, and low. Each priority queue is capable of preempting lower priority queues. A major focus for this JVD featuring LLQ is to validate the Multiple Priority Queue model.

Recommended Latency Budgets

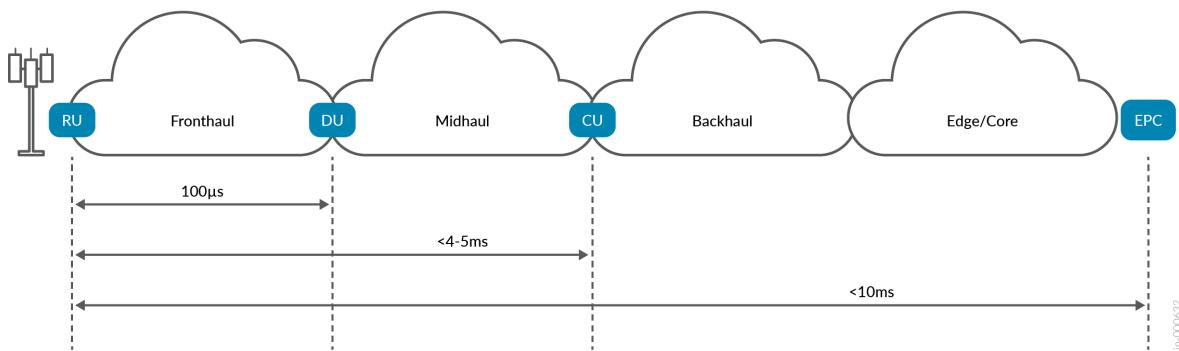
5G xHaul infrastructure defines strict latency budgets, particularly in the fronthaul segment. The total budget factors include fiber length, transit network devices, and transport design. O-RAN mandates a maximum of 100 μ s fronthaul one-way latency.

Typically, this equates to $\leq 10\mu$ s per device without accounting for fiber distances. Operators commonly seek solutions to reduce further per-device transit latency in the range of $\sim 6\mu$ s, which can help to

extend the total network reach. This is a massive paradigm shift from the earlier 4G architecture requirements. The end-to-end xHaul RU to EPC is expected to be ≤ 10 milliseconds.

When leveraging a multiple PQ model, we recommend placing G.8275.2 PTP unaware as the highest possible priority. However, O-RAN discourages deployments with PTP unaware in favor of G.8275.1 PTP profile, using boundary clock mode on each transit node. PTP aware does not carry such precise delay requirements, allowing this traffic to be placed in a low-priority queue. The highest priority traffic is therefore given to eCPRI. Traffic classes used in the JVD are explained in the "[Solution Architecture](#)" on page 23 section.

Figure 7: 5G xHaul Latency Budget



The transport architecture must be capable of preserving delay budget integrity while guaranteeing traffic priorities. The following table [O-RAN.WG9.XPSAAS-v08.00] defines suitable per-hop latency and per-hop packet delay variation budgets for critical traffic types and is an important reference point in developing the 5G JVD QoS model and contributing traffic classes. This table includes the alignment of traffic classes based on *Exemplary* grouping.

Table 3: Recommended Latency/PDV Budgets

| Traffic Type | Packet Size | Per-Hop Latency | Per-Hop PDV |
|---------------------------------------|-------------|------------------|-------------|
| PTP (unaware mode) ¹ | ~100 bytes | constant average | ~0.5 µs) |
| CPRI (RoE) | ~1500 bytes | ~1-20 µs | ~1-20 µs |
| eCPRI CU-plane | ~1500 bytes | ~1-20 µs | ~1-20 µs |
| OAM with aggressive timers | ~100 bytes | ~1 ms | ~1 ms |
| 5QI/QCI Group 1 (low latency U-plane) | variable | ~1 ms | ~1 ms |

Table 3: Recommended Latency/PDV Budgets (Continued)

| Traffic Type | Packet Size | Per-Hop Latency | Per-Hop PDV |
|---|-------------|-----------------|-------------|
| Low latency business traffic | variable | ~1 ms | ~1 ms |
| Network Control: OAM with relaxed timers, IGP, BGP, LDP, RSVP, PTP aware mode (T-TC/T-BC) | variable | ~5 ms | ~1-3 ms |
| O-RAN/3GPP C-plane and M-plane) | variable | ~5 ms | ~1-3 ms |
| 5QI/QCI Group 2 (medium latency U-plane) | variable | ~5 ms | ~1-3 ms |
| 5QI/QCI Group 3 (remaining GBR U-plane) | variable | ~10 ms | ~5 ms |
| Guaranteed business traffic | variable | ~10 ms | ~5 ms |
| 5QI/QCI Group 4 (remaining non-GBR U-plane) | variable | ~10-50 ms | ~5-25 ms |
| Other traffic types (best effort) | | | |

(1) PTP unaware mode will not be covered and is no longer recommended.

Regulatory Interests

- ETSI 3GPP Forum (3rd Generation Partnership Project)
- O-RAN Alliance
- Metro Ethernet Forum (MEF)
- IEEE P802.1CM
- *eCPRI Specification*

Validation Framework

IN THIS SECTION

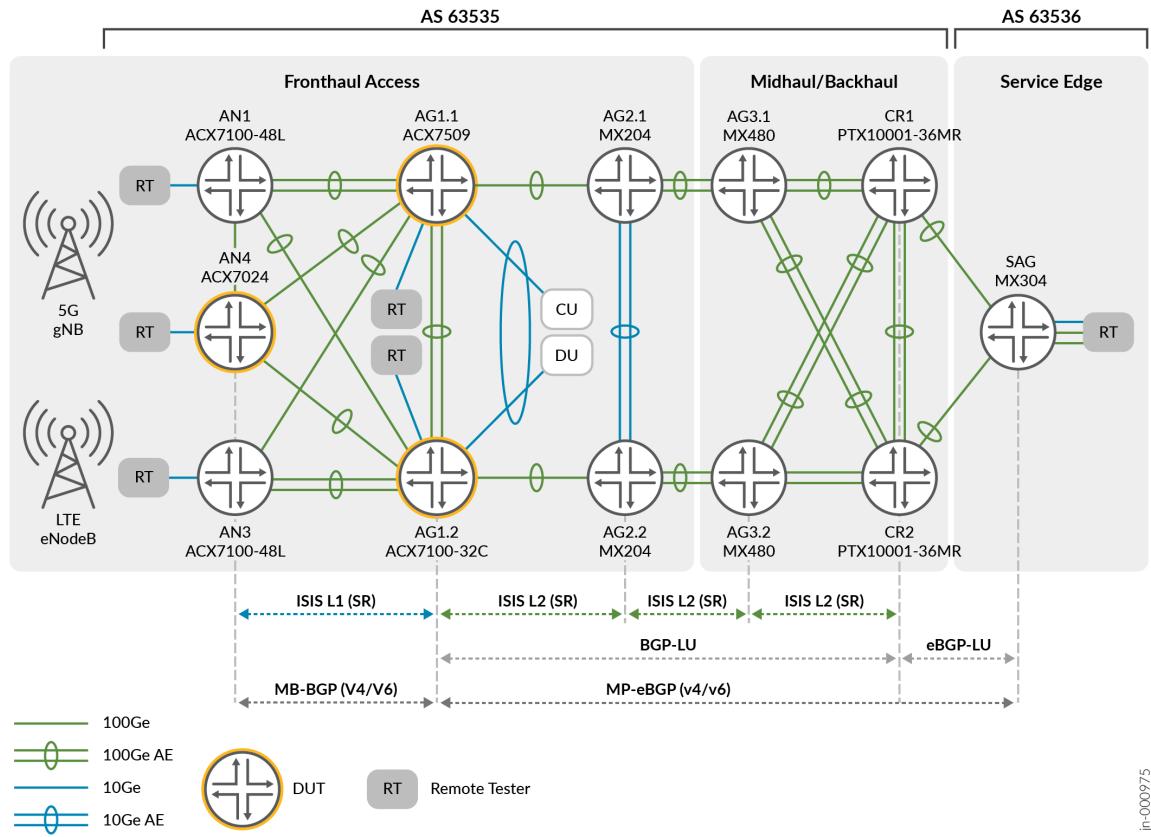
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Two xHaul network deployment scenarios are under consideration, standalone 5G Fronthaul network infrastructure and a joint deployment of *the traditional* L3VPN 4G MBH and 5G xHaul running on top of the same physical network infrastructure.

While operators are rolling out new 5G networks, you need to maintain an existing 4G infrastructure for an extended period. The same physical infrastructure must serve the purposes of L3MBH and 5G networks. The same 5G CSR must allow simultaneous connectivity to the 5G gNB and 4G eNB cell towers providing L2 fronthaul and L3 MBH services. The lab topology aligns with the co-located O-DU and O-CU split 7.2x from the cell site, however the architecture allows for any functional split with flexible insertion points across the midhaul/backhaul portion of the network.

Topology involves Seamless MPLS with BGP-LU at border nodes and IS-IS SR underlay across multiple domains. Service overlay includes EVPN, L3VPN, VPLS, L2VPN, and L2Circuit. The end-to-end CoS model that this JVD uses might be extended or evolved to support architectures and use cases beyond what this document proposes.

Figure 8: Lab Topology



Test Bed

Figure 9 on page 16 explains connectivity for building the physical 5G fronthaul CoS LLQ JVD infrastructure. The fronthaul segment leverages a spine-and-leaf topology, enabling connections between access nodes supporting cell site router (CSR) functionality to pre-aggregation nodes supporting hub site router (HSR) functions. The access segment leverages 100GbE speeds for the proposed topology, supporting up to 400GbE port speeds. CSR aggregates O-RU traffic supporting 5G, including legacy 4G workloads. The selected access nodes for the CSR role include the ACX7024 (AN4) and ACX7100-48L (AN1, AN3) platforms. HSR aggregates traffic from multiple CSRs, including additional emulated access segments and supports connectivity into the O-DU and O-CU relative to the functional split model. The selected aggregation nodes for the HSR role include the ACX7509 (AG1.1) and the ACX7100-32C (AG1.2). Additional helper nodes build out the aggregation and core segments with the MX204 (AG2), MX480 (AG3), and PTX10001-36MR (CR1, CR2) devices. Finally, the MX304 platform supports the Services Aggregation Gateway (SAG) role.

Figure 9: Test Bed

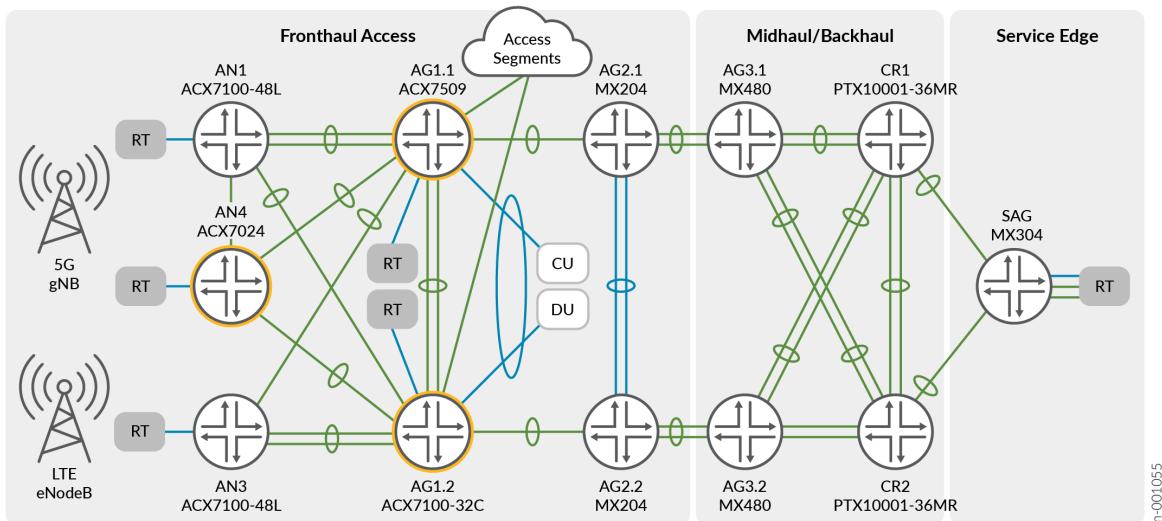


Figure 10: ACX7024 Devices Under Test



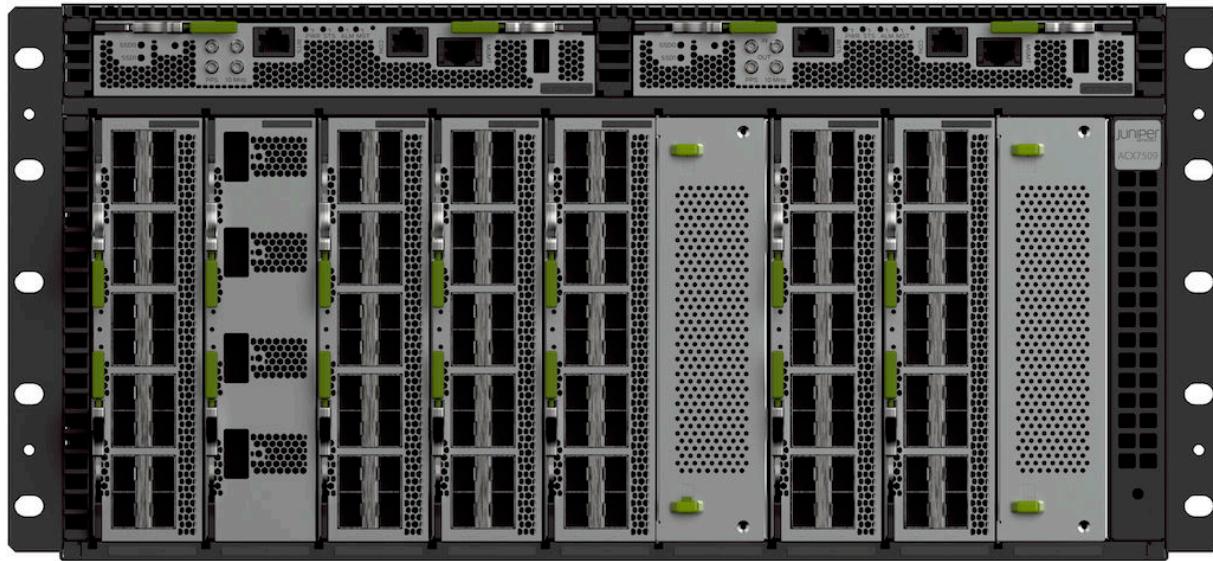
Figure 11: ACX7100-48L Devices Under Test



Figure 12: ACX7100-32C Devices Under Test



Figure 13: ACX7509 Devices Under Test



Platforms and Devices Under Test (DUT)

To review the software versions and platforms on which this JVD was validated by Juniper Networks, see the [Validated Platforms and Software](#) section in this document.

Test Objectives

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This JVD is a cross-functional collaboration between Juniper Solution Architects and Test teams to develop coherent multidimensional solutions for domain-specific use cases. The JVD team comprises technical leaders in the industry with a wealth of experience supporting complex customer use cases.

Using this JVD, you can significantly reduce the risk of costly mistakes while saving time and money in the deployment of network solutions. A JVD-based network provides benefits such as a more stable network with fewer bugs and a shorter time to resolution if any bugs are discovered. The validation process ensures that the network is optimized for maximum performance, leading to a better user experience for enterprise, service provider operation, and network service consumers. Furthermore, the design concepts deployed are formulated around best practices, leveraging relevant technologies to deliver the scope of the solution. KPIs are identified as part of an extensive test plan that focuses on functionality, performance integrity, and service delivery. With JVDs, you can shorten the time to market when implementing new network solutions, reducing the lead time to generate revenue from new services.

The key test goals of the JVD initiative include:

- Test iterative multidimensional use cases
- Optimize best practices and address solution gaps
- Validate overall solution integrity and resilience
- Support configuration and design guidance
- Deliver practical, validated, and deployable solutions

Test Goals

The main test objective is to ensure that the advanced CoS features of the ACX7000 series platforms meet the requirements for reliable and high-quality performance in 5G Fronthaul. The secondary test objective is to confirm that the CoS behaviors are consistent throughout the mobile backhaul topology and across different devices.

An intrinsic test goal of JVD is to identify any limitations, anomalies, and problem reports (PRs) that are discovered during the validation stages. These issues are addressed and fixed whenever possible, and any code changes are revalidated as part of the test execution to ensure the solution functions as intended.

The CoS operational goals are summarized by functional category. Additional details of each category are available in the ["Solution Architecture" on page 23](#) section.

Classification

- Behavior Aggregate (BA) classification is based on the received or pre-marked 802.1p, DSCP, or EXP packet headers.
- Fixed Classification maps all interface traffic to a specific forwarding class.
- Multifield classifier matches received packet attributes and maps to a forwarding class. A multifield classifier is included to match eCPRI, PTPoE, and OAM traffic or received 802.1p bits and maps to an appropriate forwarding class.
- Host-outbound exception traffic is assigned to a specified forwarding class and queue.

Queuing and Scheduling

- Eight forwarding classes and eight queues are used in the validation.
- (LLQ allows delay-sensitive data to be given preferential treatment over other traffic.
- Multi-level priority queues ensure differentiated services and applications can be appropriately prioritized across the network topology.
- Scheduler percentages are honored based on custom shaping rates (port level).
- Unused bandwidth is made available to other queues as required (up to shaping-rate) and proportional to the configured transmit-rate.
- Queueing and scheduling implementation distinctions for ACX7000 platforms:
 - Six priority levels are supported, with each higher priority pre-empting the lower: low-latency, strict-high, high, medium-high, medium-low, and low.
 - As of Junos OS Evolved Release 24.3R1, ACX supports eight priority levels for port QoS.
 - LLQ priority is serviced ahead of and pre-empts all other queues.
 - All priority queues with equal priority are serviced as round-robin.
 - Transmit-rate is not utilized on priority queues and might be shaped (PIR) to prevent starving lower priority queues.
 - Low-priority queues are serviced as WFQ round-robin in accordance with the transmit rate.
 - Only low-priority queues might use the transmit-rate configuration.
 - LLQ is given a dedicated VOQ to Egress Queue (EGQ) to ensure latency preservation.
 - When shaping rate is used, the amount is deducted from the total port speed. The remainder port speed is used to calculate transmit-rate percentage.
 - Fair Adaptive Dynamic Threshold (FADT) shared buffer allocation is updated on demand.

- Queueing and scheduling implementation distinctions for MX Trio-based platforms:
 - Five traffic priority levels are supported: strict-high, high, medium-high, medium-low, and low.
 - Queues in-profile (operating within the configured bandwidth rate) function within the guaranteed region where only strict-high queues can starve.
 - Guaranteed queues are serviced as priority queue deficit weighted round-robin (PQ-DWRR).
 - Queues out-of-profile (operating over the configured bandwidth rate) function in the excess priority region. Equal priority queues are serviced as WRR, with transmit-rate being the weight. Excess region priority is configurable, including excess transmit rate or weight.
 - An out-of-profile high-priority queue cannot starve an in-profile low-priority queue. Only strict-high operates without any excess region (therefore might be shaped to prevent starving other queues).
 - Buffer can be locked with manual configuration.

Rewrite

- Traffic is mapped to the correct queue(s), and customer or default rewrite rules are honored.
- Ingress traffic is marked as 802.1p or DSCP and rewritten for EXP classification at egress.
- Ingress traffic with EXP marking is rewritten to 802.1p or DSCP at egress.
- The rewrite operation affects the outermost tag for dual-tagged frames, preserving the inner tag.
- Preserve QoS codepoints end to end for inner and outer tags (wherever intended).
- Multiple rewrite types are supported per interface.

Shaping and Rate Limiting

- Port-level shaper honors existing scheduler transmit rates based on the provisioned shaper bandwidth.
- PQs allow rate-limiting to prevent starving low-priority queues.
- The port shaper inherits and appropriately scales the configured scheduling characteristics.
- Shaped queues do not exceed the configured shaping rate.

Latency Budget

- LLQ has the highest priority. In scenarios where the network is not congested and the line rate is below 100%, the LLQ adheres to the following latency budgets:
 - O-RU-to-O-DU latency averages $\leq 10\mu$ per device.

- RU-to-SAG latency \leq 10ms (expected \leq 100 μ s).
- eCPRI Type 5 One Way Delay measurement is approximately \leq 10 μ s per device.
- LLQ functionality ensures that the latency budget for delay-sensitive traffic is preserved in congestion and non-congestion conditions.

Key 5G Validation Attributes

- The solution ensures an end-to-end latency and stable jitter for eCPRI packets delivered between RU and DU nodes.
- Prioritize critical eCPRI fronthaul traffic during congestion and non-congestion conditions.
- Preserve latency budget of eCPRI fronthaul traffic during congestion and non-congestion conditions.
- Maintain traffic priorities across shared links.
- Maintain traffic priorities within and between VPN services that share common links.
- Allow transmission of eCPRI Type 7 event indication errors during network congestion.
- Maintain consistent CoS and resiliency across negative stress conditions (restart control and data plane daemons, add/delete configurations, and so on.).
- CSR allows all types of UNI interfaces to interconnect with 5G gNB and 4G eNB as described.
- HSR supports the proposed connectivity of O-DU and O-CU, including redundancy options.
- Transport and connectivity services are supported for 4G MBH.

Test Flow Characteristics

The JVD includes the following traffic patterns.

Accepted Frame Types

- Untagged (port-based)
- Tagged (802.1Q)
- Double-tagged (802.1ad)

Accepted Ether Types

- 0x8100 802.1Q (C-TAG or used with Single-Tagged flows)
- 0x88a8 802.1ad (S-TAG on Q-in-Q flows)

Test Non-Goals

Non-goals represent functions outside the scope of this JVD or might be covered previously.

- Custom drop profiles weighted random early detection (WRED)
- Temporal transmit rate or buffer (elastic buffer is used)
- Hierarchical CoS and H-Policing
- 802.1CM TSN Profile B with frame pre-emption (802.1Qbu)
- VLAN manipulation operations (covered in previous CoS JVD)
- Underlay or Overlay validation other than those specified in the solution goals section
- Failover and convergence scenarios (covered in previous 5G JVDs)
- End-to-End timing and synchronization distribution: synchronous Ethernet, IEEE1588v2
- G.8275.1 PTP Aware with per-node boundary clock
- SLA Monitoring: RFC 2544, Y.1564, TWAMP, and active assurance
- Telemetry, management, and automation
- Network or Link Slicing
- Flex-Algo, transport classes, BGP classful transport (covered in previous 5G JVDs)

Failure Scenarios

The validation covers the following failure scenarios:

- Link congestion
- Queue congestion without traffic discarding
- Queue congestion with traffic discarding
- Single and multiple queue congestion conditions
- Injected failure events (including eCPRI type)
- Process restart

Solution Architecture

IN THIS SECTION

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- [Service Carve Out | 28](#)
- [Class of Service Architecture | 29](#)
- [Test Topologies for Latency Measurement | 31](#)
- [O-RAN and eCPRI Emulation | 35](#)
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- [Low Latency Queuing | 38](#)
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The solution architecture deploys spine-leaf access fronthaul topology, midhaul/backhaul ring topologies are combined to include aggregation and core roles with the services gateway comprising the complete xHaul infrastructure. While the QoS implementation is generally transferrable across multiple network designs, the major components of building the end-to-end JVD solution include the following key attributes.

- 5G xHaul MBH reference architecture
- Seamless MPLS across xHaul IGP domains (Inter-AS and Inter-Domain BGP-LU)

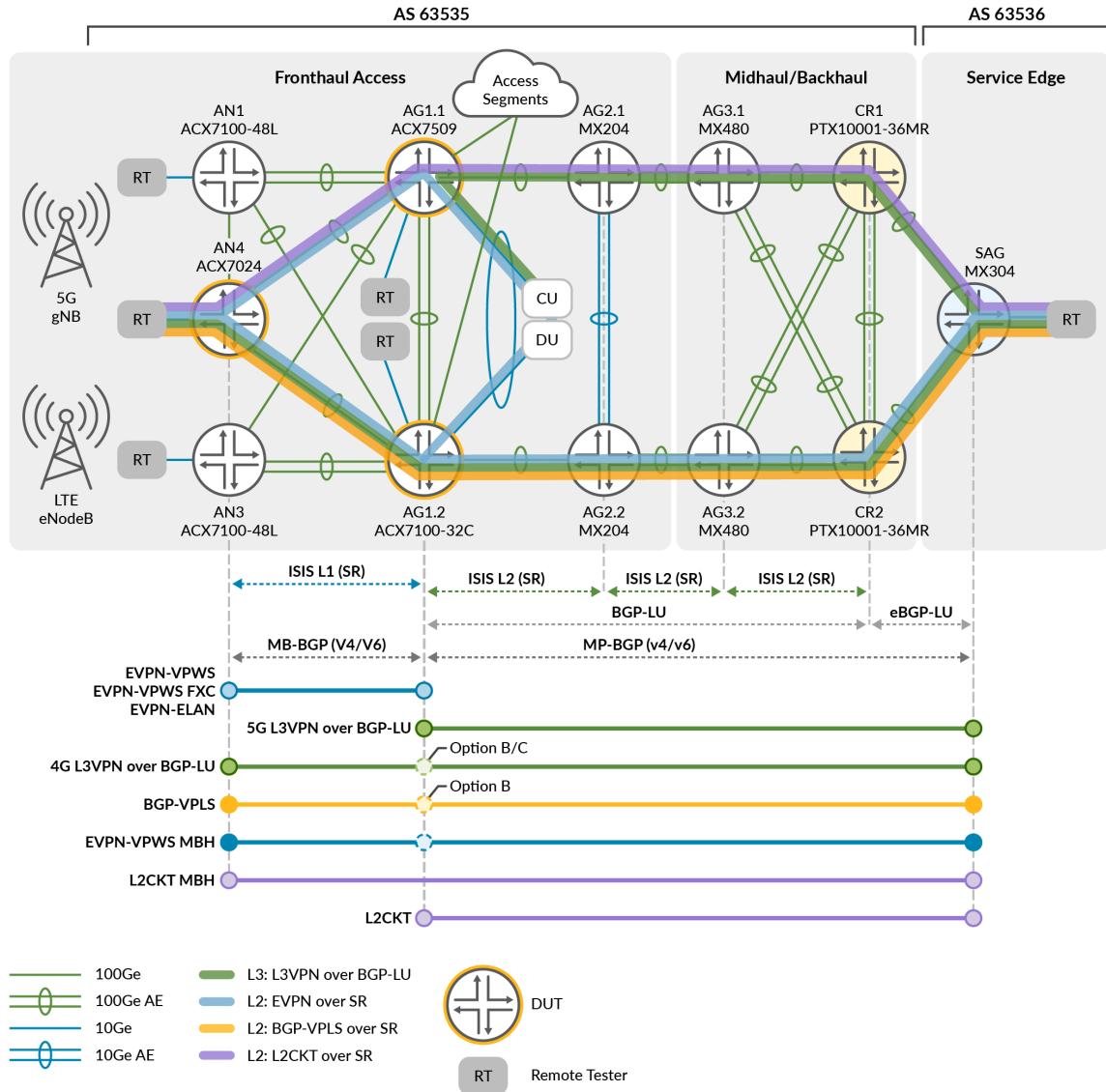
- Segment Routing L-IS-IS
- Fast failover and detection mechanisms TI-LFA, BFD, Microloop Avoidance, OAM, and so on.
- Redundant route reflectors
- Community-based route optimizations
- Inter-AS Option B/C
- EVPN-VPWS and Flexible Cross Connect (FXC) with A/A Multihoming
- EVPN-ELAN with A/A Multihoming and EVPN Virtual Gateway Address (VGA) IRB
- BGP-VPLS single-homed
- L2Circuit MBH
- L3VPN

[Figure 14 on page 25](#) describes an end-to-end service architecture for the joint 4G/5G solution where hub site AG1.1 and AG1.2 routers also support aggregation of the L3 MBH network, collapsing 4G Pre-Aggregation and 5G HSR roles. VPN PE functions are supported in parallel to Inter-AS Option-B and Option-C procedures. AG1.1/AG1.2 nodes provide O-DU connectivity and additional access segment insertion points to emulate network attachments with increased scale.

The network underlay comprises interdomain Seamless MPLS SR with BGP Labeled Unicast (BGP-LU). Access nodes (AN) are placed into an IS-IS Level 1 domain with adjacencies to L1/L2 HSR (AG1) nodes, where the Level 2 domain extends from aggregation (AG2, AG3) to core (CR) segments. Seamless MPLS is achieved by enabling BGP-LU at border nodes. TI-LFA loose mode node redundancy is enabled within each domain instantiation. The environment is scaled by two sets of route reflectors at CR1 and CR2, with westward HSR (AG1) clients. AG1.1 and AG1.2 serve as the redundant route reflectors for the access fronthaul segment. Multi-Protocol BGP peering between SAG to HSR (AG1) supports inter-AS Option-B solutions.

Services overlay incorporates modern and legacy VPNs, including EVPN, L3VPN, BGP-VPLS, and L2Circuit, with enhancements for Flow Aware Transport Pseudowire Label (FAT-PW) and Ethernet OAM where applicable.

Figure 14: Solution Architecture



Overall, the types of traffic flows represented in the topology include:

- 5G fronthaul layer 2 eCPRI between O-RU to O-DU
- 5G midhaul and backhaul layer 3 IP packet flows between 5G O-CU to UPF/5GC
- Layer 3 IP between 5G supporting open Fronthaul Management Plane and Midhaul/Backhaul Control and User plane
- 4G L3MBH IP packet flows between 4G CSR and EPC (SAG)

- 4G L2 wholesale MBH flows between CSR (AN) to EPC (SAG)

The split 7.2x includes layer 2 EVPN 5G Fronthaul traffic flows from O-RU to O-DU and L3VPN from O-CU to UPF at SAG. The architecture supports additional functional splits by simply enabling required connectivity from AG2 or AG3 segments.

Service Profiles

The following profiles explain the VPN services included in the JVD profile and how these services correlate to the 4G and 5G use cases. It is not an absolute mapping; operators might select different VPN technologies to support services across the xHaul. However, EVPN-VPWS is the primary delivery mechanism for critical 5G fronthaul flows, with L3VPN servicing C/U and M plane communications across the fronthaul, midhaul, and backhaul segments.

Additional VPN services are shown to represent potential points of connectivity, such as O-RU emulation or attachment of access regions. This might include inter-region VPN services (not requiring transport back to SAG) or connectivity into additional Telco Cloud complexes, which might be facilitated at the HSR.

Table 4: xHaul Use Cases

| Use Case | Service Overlay Mapping | Endpoints |
|--------------|--|-------------------|
| 5G Fronthaul | Fronthaul CSR to HSR EVPN-VPWS Single-Homing with E-OAM Performance Monitoring and FAT-PW | AN4 - AG1.1/AG1.2 |
| 5G Fronthaul | Fronthaul CSR-HSR EVPN-VPWS with active-active Multihoming | AN4 - AG1.1/AG1.2 |
| 5G Fronthaul | Fronthaul CSR-HSR EVPN-ELAN with active-active Multihoming | AN4 - AG1.1/AG1.2 |
| 5G Fronthaul | Fronthaul CSR-HSR EVPN-VPWS Flexible Cross Connect (FXC) Single-Homing with E-OAM Performance Monitoring | AN4 - AG1.1/AG1.2 |
| 5G Fronthaul | Fronthaul CSR-HSR EVPN-VPWS Flexible Cross Connect (FXC) with active-active Multihoming | AN4 - AG1.1/AG1.2 |
| 5G Fronthaul | Fronthaul CSR-HSR L3VPN for M-Plane | AN4 - AG1.1/AG1.2 |

Table 4: xHaul Use Cases (*Continued*)

| Use Case | Service Overlay Mapping | Endpoints |
|------------|--|-------------------|
| 5G Midhaul | EVPN IRB anycast gateway with L3VPN Multihoming DU/HSR to SAG | AG1.1/AG1.2 - SAG |
| 5G Midhaul | Bridge Domain IRB anycast static MAC/IP with L3VPN Multihoming DU/HSR to SAG | AG1.1/AG1.2 - SAG |
| L2 MBH | End-to-End L2Circuit CSR to SAG with FAT-PW | AN4 - SAG |
| L2 MBH | End-to-End Single-Homing EVPN-VPWS CSR to SAG with E-OAM and FAT-PW | AN4 - SAG |
| L2 MBH | End-to-End Single-Homing BGP-VPLS CSR to SAG with E-OAM and FAT-PW | AN4 - SAG |
| L3 MBH | End-to-End L3VPN CSR to SAG | AN4 - SAG |

Traffic Types

The following table prioritizes examples of traffic types and explains how these are mapped with the associated forwarding classes.

Table 5: Traffic Profiles

| Forwarding Class | Priority | Traffic Examples |
|------------------|-------------|---|
| Q7: FC-SIGNALING | strict-high | OAM aggressive timers, O-RAN/3GPP C-plane |
| Q6: FC-LLQ | low-latency | CPRI RoE, eCPRI C/U-Plane ≤2000 bytes |
| Q5: FC-REALTIME | medium-high | 5QI/QCI Group 1 low-latency U-plane, low latency business, Interactive video, low latency voice |
| Q4: FC-HIGH | low | 5QI/QCI Group 2 medium latency U-plane data |
| Q3: FC-CONTROL | high | Network control: OAM relaxed timers, IGP, BGP, PTP aware mode, and so on. |

Table 5: Traffic Profiles (*Continued*)

| Forwarding Class | Priority | Traffic Examples |
|--------------------|---------------|---|
| Q2: FC-MEDIUM | low | 5QI/QCI Group 3 remainder GBR U-plane guaranteed business data. Video on-demand. O-RAN/3GPP M-plane, e.g., eCPRI M-plane, other management, software upgrades |
| Q1: FC-LOW | low | high latency, guaranteed low priority data |
| Q0: FC-BEST-EFFORT | low-remainder | remainder non-GBR U-plane data |

Service Carve Out

Fronthaul traffic flows are appropriately prioritized with objective latency budgets and low delay and jitter tolerance. As a best practice, delayed critical eCPRI traffic is assigned the highest priority queue for handling low-latency workloads. Midhaul and backhaul services can be considered with lower requirements but might include URLLC use cases, such as delay-sensitive real-time video. These aspects are considered in the JVD.

The VPN services required to complete the essential 5G xHaul communications include:

- EVPN-VPWS for fronthaul C/U Plane
- L3VPN for fronthaul management plane
- L3VPN for midhaul/backhaul control plane
- L3VPN for midhaul/backhaul data plane

The environment is enhanced with additional VPN services at scale, providing performance and emulated expansion of additional network segments, e.g., CSR supporting many O-RUs.

Table 6: Service Classification

| VPN Service | Segment | Classification Type | Forwarding Classes |
|-------------|-----------|---------------------|--------------------|
| EVPN-VPWS | Fronthaul | Fixed | FC-LLQ |

Table 6: Service Classification (Continued)

| VPN Service | Segment | Classification Type | Forwarding Classes |
|-------------|-----------|---------------------|---|
| EVPN-VPWS | Fronthaul | Multifield | FC-LLQ, FC-SIGNALING, FC-CONTROL, FC-REALTIME, FC-HIGH, FC-BEST-EFFORT |
| EVPN-FXC | Fronthaul | BA | FC-SIGNALING, FC-CONTROL, FC-REALTIME, FC-HIGH, FC-MEDIUM, FC-LOW, FC-BEST-EFFORT |
| EVPN-ELAN | Fronthaul | BA | FC-SIGNALING, FC-CONTROL, FC-REALTIME, FC-HIGH, FC-MEDIUM, FC-LOW, FC-BEST-EFFORT |
| L3VPN | Fronthaul | BA | FC-SIGNALING, FC-CONTROL, FC-REALTIME, FC-HIGH, FC-MEDIUM, FC-LOW, FC-BEST-EFFORT |
| L3VPN | Midhaul | Fixed | FC-REALTIME |
| L3VPN | Midhaul | BA | FC-LLQ, FC-SIGNALING, FC-CONTROL, FC-REALTIME, FC-HIGH, FC-MEDIUM, FC-LOW, FC-BEST-EFFORT |
| L3VPN | Midhaul | Multifield | FC-REALTIME, FC-HIGH, FC-MEDIUM, FC-LOW |
| L2Circuit | MBH | Fixed | FC-HIGH, FC-MEDIUM, FC-LOW |
| EVPN-VPWS | MBH | BA | FC-SIGNALING, FC-CONTROL, FC-REALTIME, FC-HIGH, FC-MEDIUM, FC-LOW, FC-BEST-EFFORT |
| BGP-VPLS | MBH | BA | FC-HIGH, FC-MEDIUM, FC-LOW |
| L3VPN | MBH | BA | FC-SIGNALING, FC-CONTROL, FC-REALTIME, FC-HIGH, FC-MEDIUM, FC-LOW, FC-BEST-EFFORT |

Class of Service Architecture

Flows are sent through access nodes (Figure 15 on page 30) toward O-DU or SAG with Layer 2 (802.1p) or Layer 3 (DSCP) marking at positions defined as *classification* and rewritten to EXP at egress across the SR-MPLS topology. Queue statistics are monitored to ensure that the intended classification and scheduling produce the expected results. Rewrite operations occur at the specified positions. Packet

captures are taken to verify that DSCP, 802.1p, or EXP bits are correctly rewritten or preserved. In the inverse direction, flows sent through SAG are marked and validated once egressing the access nodes.

The 5G CoS LLQ JVD builds upon previous validated designs featuring ACX, MX, and PTX platform families to focus on delivering differentiated services requiring ultra-low latency. This profile incorporates capabilities of ACX7000 platforms in the fronthaul to dedicate queuing machinery to preserve low latency workloads. LLQ enables support of a more comprehensive O-RAN traffic profile for multi-level priority QoS.

MX304 as SAG and PTX10001-36MR core nodes, implements a corresponding configuration to support the intended use cases and diverse services.

Figure 15: Class of Service Points of Validation

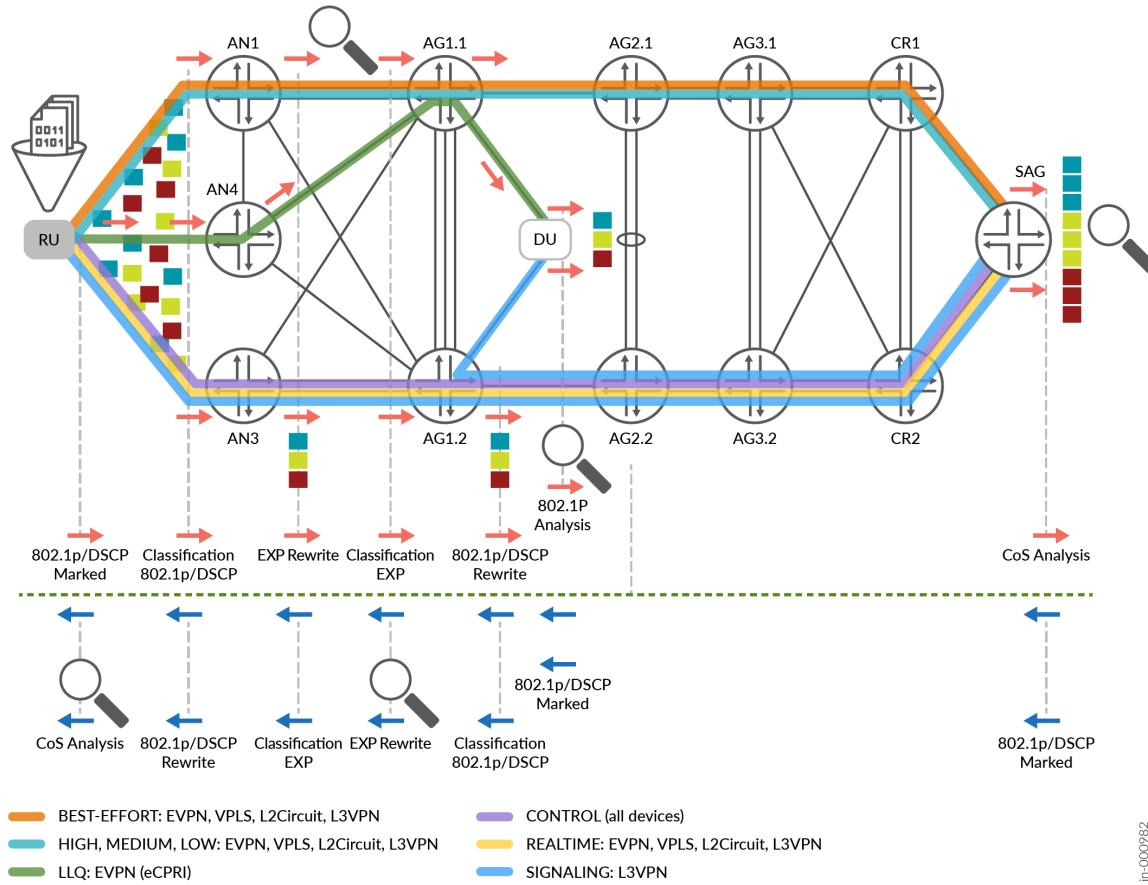


Figure 15 on page 30 is presented as a generalization of the end-to-end traffic, differentiating fronthaul vs backhaul and points of CoS operations to validate these critical functions, e.g., classification, scheduling, and rewrite operations. The diagram does not represent the precise per-flow path selection used in the JVD. The points of packet capture for further inspection are indicated.

The major objective is to examine predictable behaviors regarding how critical and noncritical traffic flows are handled across 5G xHaul network services. CoS functionalities are validated across EVPN, L3VPN, BGP-VPLS, and L2Circuit VPNs, aligning with 5QI traffic classification. QoS implementation should exhibit deterministic functionality. The transport architecture must be capable of supporting the adaptation of existing and emerging mobile applications, preserving delay budget integrity while guaranteeing traffic priorities.

Test Topologies for Latency Measurement

The JVD proposes multiple topologies to measure and report on latency performance for each DUT and across fronthaul and backhaul network segments. This ensures to capture meaningful data on the device performance across different functional roles and topologies.

- Topology 1 validates the individual DUT performance acting in a CSR or HSR role. This topology provides most accurate performance for the individual device.
- Topology 2 measures performance over point-to-point EVPN-VPWS services between CSR and HSR. ACX7024 acts as CSR DUT with ACX7100-32C or ACX7509 as HSR DUTs. The test equipment emulates O-RU and O-DU.
- Topology 3 measures performance across EVPN-VPWS active-active multihoming connectivity with a single CSR DUT to a pair of HSRs leveraging All-Active ESI LAG toward DU (QFX5120).

Critical flows are crafted to emulate the fronthaul traffic patterns with burst and steady streams (validating each pattern) across a variety of packet sizes. The packet sizes include 64B – 2000B for fronthaul flows, with a maximum frame size of 2020 bytes, including up to 7-byte Preamble, 1-byte Start of Frame Delimiter (FSG), and 20-byte Inter Frame Gap (IFG). Each test scenario includes fronthaul priority flows where performance measurements are taken based on the frame size.

As per TSN Profile A, non-fronthaul traffic should include a maximum frame size of 2000 octets, though the validation includes scenarios with jumbo frames for comparison. Background traffic is generated with iMIX 64B to 2000B frame size from SAG (xHaul) toward O-DU or O-RU to create congestion scenarios measured at converging DUT egress points. Test equipment traffic excludes self-latency. However, it should be noted that where topology includes a physical DU (QFX Series Switch), latency is incurred and counted against the total as a transit next hop.

For more information on 5G specifications, see the ["Recommended Latency Budgets" on page 11](#) section.

Figure 16: Topology 1a ACX7024 Standalone DUT

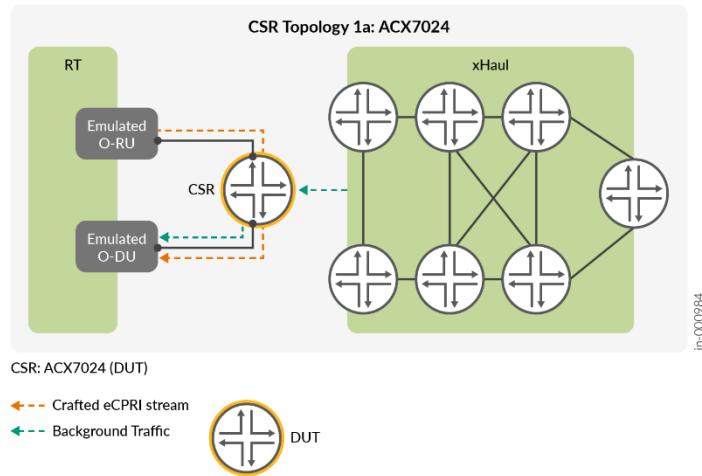


Figure 17: Topology 1b ACX7100-32C Standalone DUT

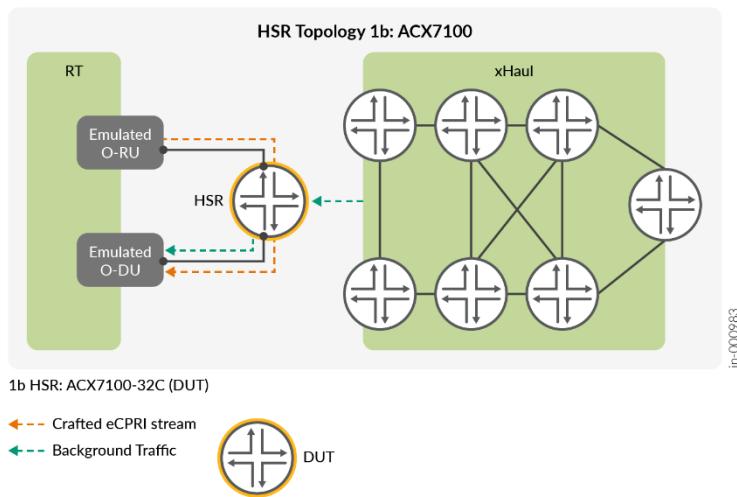
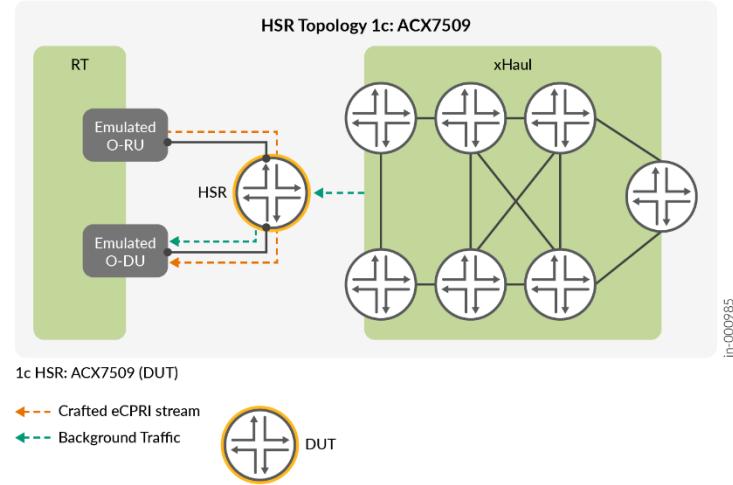


Figure 18: Topology 1c ACX7509 Standalone DUT



Topology 1 (a, b, c) is crucial to understanding the individual DUTs latency contribution. Flows of various packet sizes of burst and continuous traffic patterns are included. Background traffic includes fixed frame size and iMIX traffic flows. For more details, you can refer to the Test Report. The configuration is locally switched on DUT. The included DUTs are ACX7024 as the primary CSR, ACX7100-32C as HSR1, and ACX7509 as HSR2.

Figure 19: Topology 2a CSR and HSR

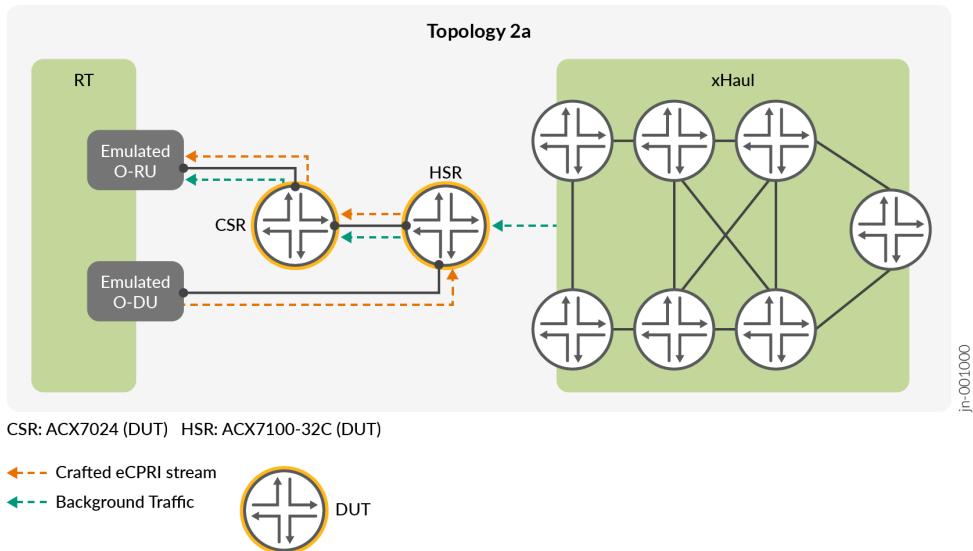
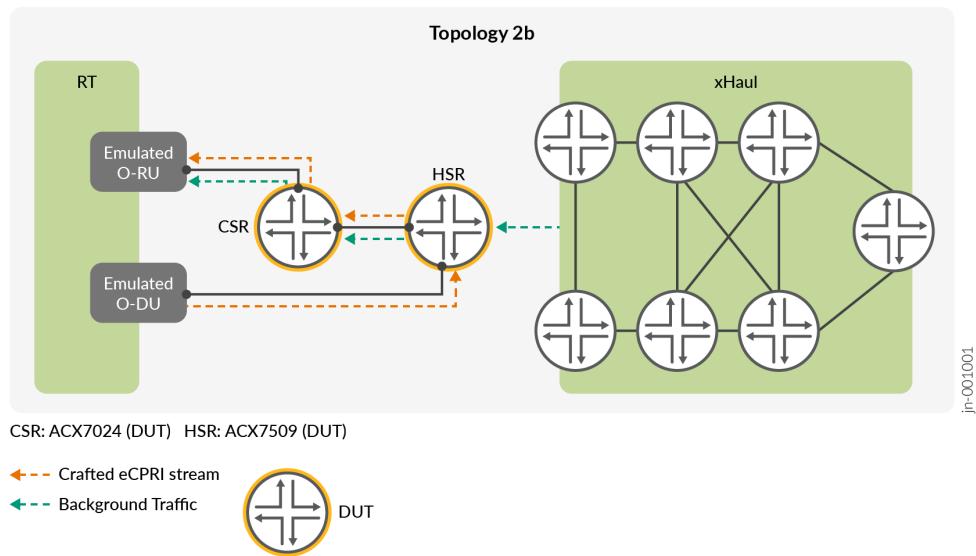
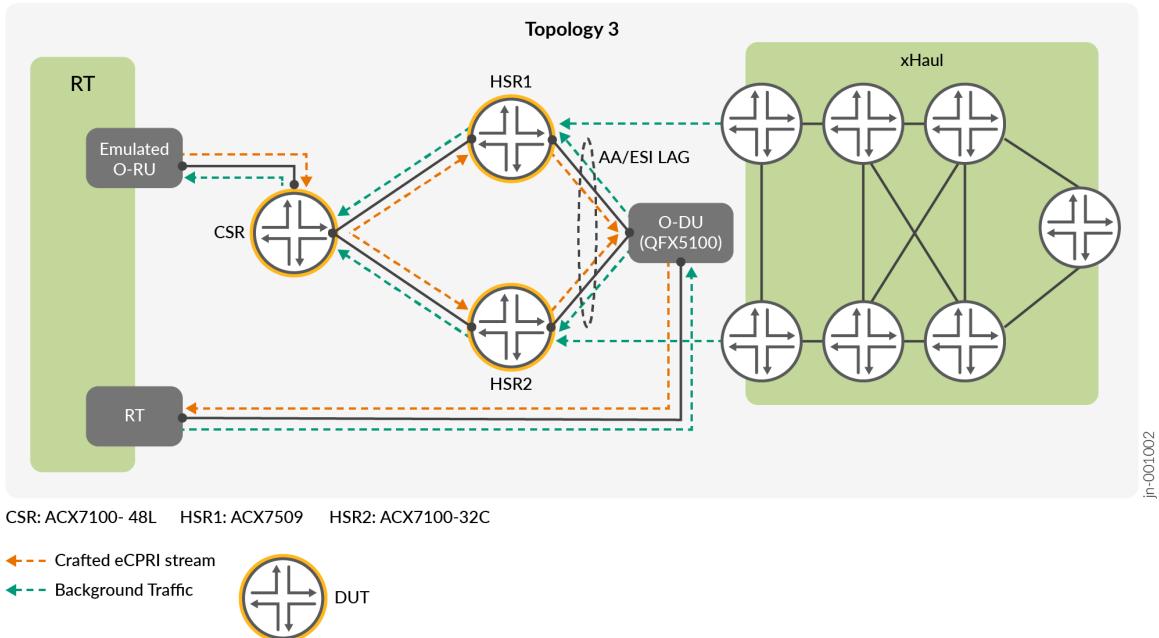


Figure 20: Topology 2b CSR and HSR



Topology 2 provides single-homed performance data in a 2-hop scenario without a physical RU or DU element (both emulated by the test center).

Figure 21: Topology 3 CSR to active-active HSR



Topology 3 provides multihomed performance data in a 3-hop scenario, which includes QFX Series Switches as the DU connecting to the all-active ESI LAG. The test center emulates RU.

O-RAN and eCPRI Emulation

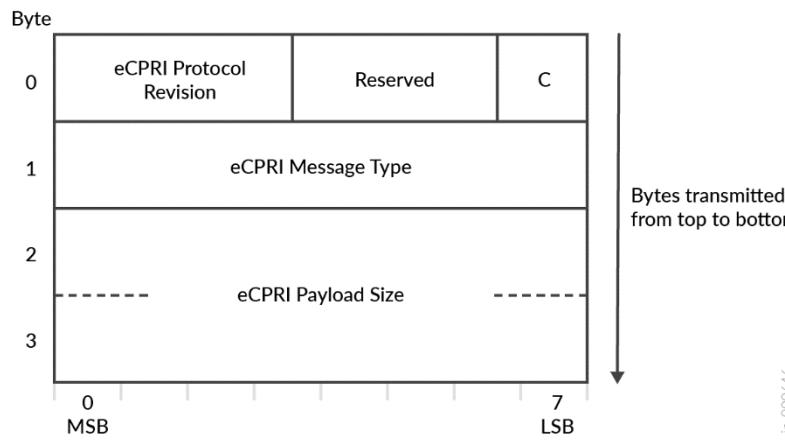
The O-RAN and eCPRI test scenarios include several functional permutations with emulated O-DU and Remote Radio Unit (RRU). Innovative steps are taken to validate the performance and assurance of 5G eCPRI communications, U-Plane message-type behaviors, and O-RAN conformance. These steps ensure that the featured DUTs are correctly and consistently transporting critical 5G services.

Results evolve across multiple topologies to capture a range of DUT performance characteristics in non-congested and congested scenarios.

The summary of O-RAN and eCPRI test scenarios include:

- eCPRI O-RAN Emulation: Leverages a standard IQ sample file to generate flows.
- O-RAN Conformance: Analyzes messages for conformance to O-RAN specification.
- Crafted eCPRI O-RAN: Produces variable eCPRI payload for comparing latency performance.
- eCPRI Services Validation: Emulates user plane messages, performing functional and integrity analysis.
- eCPRI Remote Memory Access (Type 4): Performs read or write from or to a specific memory address on the opposite eCPRI node and validates expected success or failure conditions.
- eCPRI Delay Measurement Message (Type 5): Estimates one-way delay between two eCPRI ports.
- eCPRI Remote Reset Message (Type 6): One eCPRI node requests a reset of another node. It validates expected sender or receiver operations.
- eCPRI Event Indication Message (Type 7): Either side of the protocol indicates to the other side that an event has occurred. An event is either a raised or ceased fault or a notification. Validation confirms events raised as expected.

Figure 22: eCPRI User Plane Common Header Format



VLAN Operations

The ACX7000 series supports a comprehensive set of VLAN manipulation operations. For more information on the 80 VLAN combinations tested across L2Circuit, BGP-VPLS, EVPN-VPWS, and EVPN-ELAN services, see the validated design for [5G Mobile xHaul CSR](#).

The following table summarizes VLAN operations supported by ACX EVO-based platforms.

Table 7: ACX-EVO Supported VLAN Operations

| IFL Tag Type | Input Map Operation | Output Map Operation | Eth Bridge | Vlan Bridge | Supported IFD Tagging |
|--------------|---------------------|----------------------|------------|-------------|-----------------------|
| UT | None | None | Yes | No | None |
| UT | Push | Pop | Yes | No | None |
| UT | Push-Push | Pop-Pop | Yes | No | None |
| ST | None | None | No | Yes | VLAN and Flex Tagging |
| ST | Push | Pop | No | Yes | VLAN and Flex Tagging |

Table 7: ACX-EVO Supported VLAN Operations (Continued)

| IFL Tag Type | Input Map Operation | Output Map Operation | Eth Bridge | Vlan Bridge | Supported IFD Tagging |
|--------------|---------------------|----------------------|------------|-------------|-----------------------|
| ST | Swap | Swap | No | Yes | VLAN and Flex Tagging |
| ST | Pop | Push | No | Yes | VLAN and Flex Tagging |
| ST | Push-Swap | Swap-Pop | No | Yes | VLAN and Flex Tagging |
| DT | None | None | No | Yes | VLAN and Flex Tagging |
| DT | Pop | Push | No | Yes | VLAN and Flex Tagging |
| DT | Swap | Swap | No | Yes | VLAN and Flex Tagging |
| DT | Swap-Swap | Swap-Swap | No | Yes | VLAN and Flex Tagging |
| DT | Pop-Swap | Swap-Push | No | Yes | VLAN and Flex Tagging |
| DT | Pop-Pop | Push-Push | No | Yes | VLAN and Flex Tagging |
| Native ST | None | None | No | Yes | Flexible Tagging |
| Native ST | Push | Pop | No | Yes | Flexible Tagging |
| Native ST | Swap | Swap | No | Yes | Flexible Tagging |
| Native ST | Pop | Push | No | Yes | Flexible Tagging |
| Priority ST | Push | Pop | No | Yes | VLAN and Flex Tagging |

Table 7: ACX-EVO Supported VLAN Operations (Continued)

| IFL Tag Type | Input Map Operation | Output Map Operation | Eth Bridge | Vlan Bridge | Supported IFD Tagging |
|--------------|---------------------|----------------------|------------|-------------|-----------------------|
| Priority ST | Swap | Swap | No | Yes | VLAN and Flex Tagging |
| Priority ST | Pop | Push | No | Yes | VLAN and Flex Tagging |

Test scenarios include the following VLAN operations:

- Pop s-tag in outer position (on top of 802.1Q tagged frames)
- Push s-tag (moving c-tag to inner position)
- Swap outer tag
- Classification based on 802.1p priority code point (PCP) bits
- Rewrite outer VLAN tag 802.1p bits
- Preservation of c-tag PCP bits
- Translate or rewrite the VLAN tag of 802.1Q frames
- Multiple untagged, single-tagged, and dual-tagged operations

For single-tagged and dual-tagged operations, the classification of incoming frames is done based on outer 802.1Q Ethernet header PCP bits. All traffic types (single-tagged, dual-tagged, or untagged) are classified into appropriate forwarding classes, determined by VPN service type, with EVPN (eCPRI + critical fronthaul flows) given the highest priority and lowest latency.

Low Latency Queuing

Junos OS Evolved Release 23.3R1 introduces the LLQ for ACX7000 platforms, enabling delay-sensitive data to be given preferential treatment over other traffic by allowing dequeuing of data so that priority traffic can be sent first.

With LLQ functionality, the queue is prioritized over all other priority queues to ensure latency is preserved. The Virtual Output Queues (VOQ) and Egress Queue (EGQ) priority hierarchy is as follows:

Latency Queues > Priority Queues > Low Queues

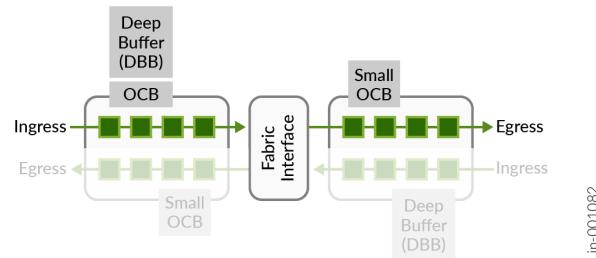
With Latency Queueing support, applications that require low latency queuing have the option to configure this priority through CLI: `set class-of-service schedulers priority low-latency`

When queues (other than LLQ) are congested, LLQ is expected to support $10\mu\text{s}$ average latency. The LLQ mechanics are capable of $\leq 6\mu\text{s}$ without congestion. Multiple factors might influence delay. The validation demonstrates LLQ capability exceeding expectations.

LLQ supports multiple queues of the same priority. However, we recommend to include no more than two LLQs per system to ensure latency integrity is preserved. Reasonably, the more latency queues that must perform round-robin distribution create delay. A PFE syslog warning is given if more than two LLQs are configured.

The ACX7000 family implements a VOQ architecture designed to avoid Head of Line Blocking (HOLB) and optimize buffering. This approach uses feedback mechanisms between the ingress traffic manager (ITM) and the egress traffic manager (ETM). Packets received at ingress VOQs are mapped to the destination egress port. The ITM generates a credit request and sends it to the ETM. Depending on the egress port and queue credit availability, the ETM generates appropriate credit and schedules the traffic with minimal buffering. Ingress VOQ bundles are associated with a specific egress port. Each bundle consists of eight VOQs. The VOQ connection forms a virtual representation of the queue and an egress port association called a Virtual Output Queue Identifier (VOQ ID).

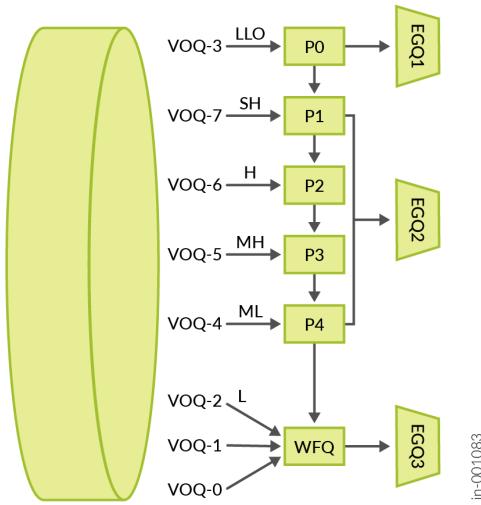
Figure 23: Virtual Output Queuing Structure



Packet buffering is performed at ingress with a large delay bandwidth buffer (DBB) and SRAM small on-chip buffer (OCB). An additional on-chip egress buffer is leveraged for packet serialization.

[Figure 24 on page 40](#) explains the ACX7000 scheduling hierarchy in Port QoS mode. Along with priority, a low delay scheduling profile is associated with VOQs configured as low latency. With Port QoS mode, three distinct dedicated Egress Queues (EGQ) are enabled to support LLQ, Priority Queues, and Low Queues.

Figure 24: Multi-Level Priority Hierarchy with LLQ Port QoS



EGQ1 is dedicated to handling VOQs configured with low latency. EGQ2 is dedicated to handling priority queues other than latency queues. EGQ3 is dedicated to handling low-priority traffic. The EGQ separation allows independent treatment to each category while minimizing the impact on each other. VOQ priority association is based on configurations.

Multi-Level Priority

Junos OS Evolved Release 23.3R1 introduces multi-level priorities for the ACX7000 series. Six priorities are supported: low-latency, strict-high, high, medium-high, medium-low, and low. All queues preempt the lower priority. Equal-priority queues perform round-robin distribution, and low-priority queues are WFQ. This support is extended to both port-level QoS and hierarchical QoS (HQoS).

NOTE: Only port-based QoS is covered in this validation.

Before Junos OS Evolved release 23.3R1, the ACX7000 family supported strict-high and low priority queues, with SH preempting low. In the new CoS architecture, each queue is given a priority level that is capable of preempting lower priority queues. These priority levels are:

- P0 (highest priority): Low-Latency
- P1: Strict-High
- P2: High

- P3: Medium-High
- P4: Medium-Low
- WFQ (lowest priority): Low

Priority Queues are designated from P0-P4 and leverage round-robin distribution for queues within the same priority level. Low priority is the only WFQ, where the transmit rate is the weight given for the round-robin distribution of multiple low-priority queues. With strict queue preemption, we recommend to shape the priority queues to prevent starving lower priority queues.

The functional behavior differs from TRIO-based MX architectures, which leverage guaranteed and excess regions to ensure the configured transmit rate for all priority queues, including low priority. PQ-DWRR is implemented in the guaranteed region. Equal priority queues are serviced as WRR, with the transmit-rate being the weight. Only the strict-high queue operates without an excess region and can starve low-priority queues.

NOTE: From Junos OS Evolved Release 24.3R1, ACX7000 supports eight priority levels, with the addition of low-high (P5) and low-medium (P6) for port QoS.

Forwarding Classes

ACX7000 supports eight forwarding classes (FC), with four being enabled by default. One or more FCs can be mapped up to the eight supported queues. Conversely, MX-series supports sixteen forwarding classes, which can be mapped across eight queues. For this JVD, a maximum of eight queues are utilized for each transit device, supporting O-RAN proposed traffic profiles.

The JVD CoS model most closely aligns with the O-RAN Multiple Priority Queue structure and includes the following queue assignments:

Table 8: Queue Assignments

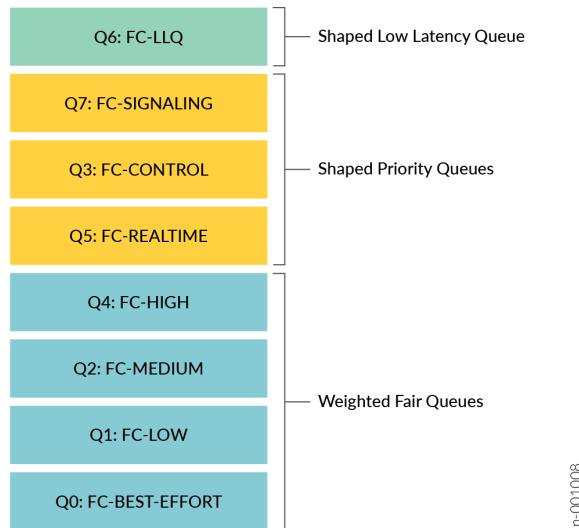
| QUEUE | Forwarding Class | Queue Characteristics |
|-------|------------------|---|
| 7 | FC-SIGNALING | strict-high priority and PIR shaped |
| 6 | FC-LLQ | low-latency priority and PIR shaped (eCPRI) |
| 5 | FC-REALTIME | medium-high priority and PIR shaped (voice & interactive video) |

Table 8: Queue Assignments (Continued)

| QUEUE | Forwarding Class | Queue Characteristics |
|-------|------------------|--|
| 4 | FC-HIGH | low priority WFQ guaranteed |
| 3 | FC-CONTROL | High priority and PIR shaped (network control) |
| 2 | FC-MEDIUM | low priority WFQ guaranteed |
| 1 | FC-LOW | low priority WFQ guaranteed |
| 0 | FC-BEST-EFFORT | low priority WFQ remainder |

Figure 25 on page 42 illustrates the CoS model proposed for the solution architecture. The CoS hierarchy can essentially be divided into three main components.

- Low Latency Queues
- Shaped Priority Queues
- Weighted Fair Queues

Figure 25: Differentiated CoS Queuing Model

Critical fronthaul flows are mapped to the LLQ or PQs, which allows a parallel queueing structure for non-fronthaul flows with WFQs. The high, medium, and low traffic types are given a guaranteed

committed information rate (CIR) while allowing a dynamic peak information rate (PIR) if bandwidth is available. To achieve this result, the transmit rates configured for low-priority queues (WFQ) must be calculated based on the appropriate port speed, excluding bandwidth allocated to shaped queues. For more information, see the ["Scheduling" on page 49](#) section. The best-effort (remainder) queue has no guarantee.

When implementing the 5G architecture, the LLQ PQ model might be contained within the fronthaul segment with WFQs designated for MBH. This is not a goal for the JVD. In general, fronthaul flows are mapped to the first three queues, as shown in [Figure 25 on page 42](#), with FC-REALTIME supporting URLLC end-to-end applications. Non-fronthaul (i.e., MBH) is always mapped to the lower four low-priority queues (WFQ). FC-CONTROL is applicable to all devices. Low-priority queues might still support some functions within the fronthaul, such as M-plane and software upgrades.

The following example displays the Forwarding Class configuration, which is the same across ACX, MX, and PTX platforms.

```
ACX-EVO Forwarding Class
set class-of-service forwarding-classes class FC-SIGNALING queue-num 7
set class-of-service forwarding-classes class FC-LLQ queue-num 6
set class-of-service forwarding-classes class FC-REALTIME queue-num 5
set class-of-service forwarding-classes class FC-HIGH queue-num 4
set class-of-service forwarding-classes class FC-CONTROL queue-num 3
set class-of-service forwarding-classes class FC-MEDIUM queue-num 2
set class-of-service forwarding-classes class FC-LOW queue-num 1
set class-of-service forwarding-classes class FC-BEST-EFFORT queue-num 0
```

Example traffic types per queue:

Table 9: Forwarding Class with Traffic Types

| QUEUE | Forwarding Class | Example Traffic Type |
|-------|------------------|---|
| 7 | FC-SIGNALING | OAM aggressive timers, O-RAN/3GPP C-plane |
| 6 | FC-LLQ | CPRI RoE, eCPRI C/U-Plane ≤2000 bytes |
| 5 | FC-REALTIME | 5QI/QCI Group 1 low-latency U-plane, low-latency business. Interactive video, low latency voice |
| 4 | FC-HIGH | 5QI/QCI Group 2 medium latency U-plane data |

Table 9: Forwarding Class with Traffic Types (Continued)

| QUEUE | Forwarding Class | Example Traffic Type |
|-------|------------------|---|
| 3 | FC-CONTROL | Network control: OAM relaxed timers, IGP, BGP, PTP aware mode, and so on |
| 2 | FC-MEDIUM | 5QI/QCI Group 3 remainder GBR U-plane guaranteed business data. Video-on-demand. O-RAN/3GPP M-plane, e.g., eCPRI M-plane, other management, software upgrades |
| 1 | FC-LOW | high latency, guaranteed low-priority data |
| 0 | FC-BEST-EFFORT | 5QI/QCI Group 4 – remainder non-GBR U-plane data |

Classification

Classification is performed at ingress. The following three styles of classification are included in the validation.

- Behavior Aggregate (BA) matches on the received layer 2 802.1p bits and/or layer 3 DSCP. In the event where both are received, DSCP takes priority. Behavior Aggregate is *packet-based*, where flows are pre-marked with layer 3 DSCP, layer 2 802.1Q Priority Code Points (PCP), or MPLS EXP.
- Fixed or Discrete Classification allows the forwarding class to map directly on the interface. Fixed classification is *context-based*, where all traffic arriving on a specific interface is mapped to one forwarding class.
- Multifield (MF) Classification allows matching criteria within the packet fields and maps to one or more forwarding classes.

Classification on transit nodes is performed at ingress and is matched upon outer label MPLS EXP bits. Core interfaces utilize BA classifiers. BA, Fixed, and MF-style classifiers are used at service interfaces (CE-facing).

When both BA and MF classifications are performed simultaneously, MF takes priority since BA is processed first and followed by MF processing. As a result, MF overrides BA results (assuming a match).

In the use cases where received tagged frames include trusted PCP bits, behavior aggregate classification is used. For the use cases where received frames are untrusted PCP bits, fixed or multifield classification is used.

The classification posture is the same across ACX, MX, and PTX platforms included in the JVD. ["Table 10" on page 45](#) describes the overall priority mapping.

Table 10: Classification Definitions

| QUEUE | Forwarding Class | 802.1p | DS/PCP | EXP |
|-------|------------------|--------|-----------------------|-----|
| 7 | FC-SIGNALING | 110 | CS5, CS6 | 110 |
| 6 | FC-LLQ | 100 | CS4, AF41, AF42, AF43 | 100 |
| 5 | FC-REALTIME | 101 | EF | 101 |
| 4 | FC-HIGH | 011 | CS3, AF31, AF32, AF33 | 011 |
| 3 | FC-CONTROL | 111 | CS7 | 111 |
| 2 | FC-MEDIUM | 010 | CS2, AF21, AF22, AF23 | 010 |
| 1 | FC-LOW | 001 | CS1, AF11, AF12, AF13 | 001 |
| 0 | FC-BEST-EFFORT | 000 | BE | 000 |

EXP Classifier

```
set class-of-service classifiers exp CL-MPLS import default
set class-of-service classifiers exp CL-MPLS forwarding-class FC-SIGNALING loss-priority low code-points 110
set class-of-service classifiers exp CL-MPLS forwarding-class FC-LLQ loss-priority low code-points 100
set class-of-service classifiers exp CL-MPLS forwarding-class FC-REALTIME loss-priority low code-points 101
set class-of-service classifiers exp CL-MPLS forwarding-class FC-HIGH loss-priority low code-points 011
set class-of-service classifiers exp CL-MPLS forwarding-class FC-CONTROL loss-priority low code-points 111
set class-of-service classifiers exp CL-MPLS forwarding-class FC-MEDIUM loss-priority low code-points 010
set class-of-service classifiers exp CL-MPLS forwarding-class FC-LOW loss-priority low code-points 001
set class-of-service classifiers exp CL-MPLS forwarding-class FC-BEST-EFFORT loss-priority low code-points 000
```

802.1P Classifier

```
set class-of-service classifiers ieee-802.1 CL-8021P import default
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-SIGNALING loss-priority low code-
points 110
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-LLQ loss-priority low code-points 100
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-REALTIME loss-priority low code-
points 101
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-HIGH loss-priority low code-points
011
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-CONTROL loss-priority low code-
points 111
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-MEDIUM loss-priority low code-points
010
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-LOW loss-priority low code-points 001
set class-of-service classifiers ieee-802.1 CL-8021P forwarding-class FC-BEST-EFFORT loss-priority low code-
points 000
```

DSCP IPv4 Classifier

```
set class-of-service classifiers dscp CL-DSCP import default
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-SIGNALING loss-priority low code-points cs6
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-SIGNALING loss-priority high code-points cs5
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LLQ loss-priority low code-points cs4
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LLQ loss-priority low code-points af41
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LLQ loss-priority medium-high code-points
af42
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LLQ loss-priority high code-points af43
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-REALTIME loss-priority low code-points ef
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-HIGH loss-priority low code-points cs3
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-HIGH loss-priority low code-points af31
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-HIGH loss-priority medium-high code-points
af32
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-HIGH loss-priority high code-points af33
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-CONTROL loss-priority low code-points cs7
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-MEDIUM loss-priority low code-points cs2
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-MEDIUM loss-priority low code-points af21
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-MEDIUM loss-priority medium-high code-
points af22
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-MEDIUM loss-priority high code-points af23
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LOW loss-priority low code-points cs1
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LOW loss-priority low code-points af11
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LOW loss-priority medium-high code-points
af12
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-LOW loss-priority high code-points af13
set class-of-service classifiers dscp CL-DSCP forwarding-class FC-BEST-EFFORT loss-priority high code-points
be
```

DSCP IPv6 Classifier

```
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 import default
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-SIGNALING loss-priority low code-points cs6
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-SIGNALING loss-priority high code-points cs5
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LLQ loss-priority low code-points cs4
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LLQ loss-priority low code-points af41
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LLQ loss-priority medium-high code-points af42
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LLQ loss-priority high code-points af43
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-REALTIME loss-priority low code-points ef
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-HIGH loss-priority low code-points cs3
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-HIGH loss-priority low code-points af31
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-HIGH loss-priority medium-high code-points af32
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-HIGH loss-priority high code-points af33
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-CONTROL loss-priority low code-points cs7
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-MEDIUM loss-priority low code-points cs2
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-MEDIUM loss-priority low code-points af21
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-MEDIUM loss-priority medium-high code-points af22
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-MEDIUM loss-priority high code-points af23
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LOW loss-priority low code-points cs1
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LOW loss-priority low code-points af11
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LOW loss-priority medium-high code-points af12
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-LOW loss-priority high code-points af13
set class-of-service classifiers dscp-ipv6 CL-DSCP-IPV6 forwarding-class FC-BEST-EFFORT loss-priority high code-points be
```

The following configuration is an example of multifield classification matching distinct traffic types to assign a forwarding class. The filter might be modified as required to adapt an appropriate match and/or prioritization requirements.

```
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term eCPRI from ether-type 0xAEFE
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term eCPRI then forwarding-class FC-LLQ
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term eCPRI then count eCPRI-in
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term PTPoE from destination-mac-address
01:1b:19:00:00:00/48
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term PTPoE from ether-type 0x88F7
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term PTPoE then forwarding-class FC-CONTROL
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term CFM from ether-type 0x8902
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term CFM then forwarding-class FC-SIGNALING
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term REALTIME from learn-vlan-1p-priority 5
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term REALTIME then forwarding-class FC-REALTIME
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term HIGH from learn-vlan-1p-priority 3
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term HIGH then forwarding-class FC-HIGH
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term MEDIUM from learn-vlan-1p-priority 2
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term MEDIUM then forwarding-class FC-MEDIUM
set firewall family ethernet-switching filter FF-5G-LLQ-CLASS term INTERNET then forwarding-class FC-BEST-
EFFORT
```

Scheduling

ACX7000 scheduling begins in the ingress pipeline and is realized as an egress function with a feedback loop mechanism. Once the packet is classified and arrives at the ingress VOQ, the Ingress Traffic Manager (ITM) issues a credit request to the Egress Traffic Manager (ETM). Depending on the egress port and queue availability, a credit request is granted by ETM, and packets are scheduled and finally arrive at the Egress Queue (EGQ).

The 8-Queues (configurable per port) are associated with the VOQ architecture rather than EGQ. The VOQs map into the two or three EGQs depending on the implementation. The scheduling architecture for ACX7000 platforms has changed from Junos OS Evolved Release 23.3R1. For more information about the updated EGQ hierarchy with LLQ, see the ["Low Latency Queuing" on page 38](#) section.

Schedulers assign traffic priorities and bandwidth characteristics to the forwarding classes. ACX7000 platforms support six traffic priorities (Eight priorities as per Junos OS Evolved Release 24.3R1). MX-series platforms support five traffic priorities.

Priorities are associated with VOQs, and VOQs map to EGQs. As a result, the mapping of priority level to EGQ assignment is shown here as the 8-priority level model.

Table 11: Priority to EGQ Mapping

| Priority Name | Priority Level | Egress Queue Assignment |
|---------------|----------------|-------------------------|
| Low-Latency | P0 | EGQ1 |
| Strict-High | P1 | EGQ2 |
| High | P2 | EGQ2 |
| Medium-High | P3 | EGQ2 |
| Medium-Low | P4 | EGQ2 |
| Low-High | P5 | EGQ2 |
| Low-Medium | P6 | EGQ2 |
| Low (WFQ) | P7 | EGQ3 |

As shown, multiple priority queues map to the same EGQ, so it is important to understand the nature of the ACX ingress pipeline. ACX efficient packet processing mechanisms prevent Head-of-Line-Blocking (HOLB) and avoid buffering packets that might be later dropped. Once a packet is scheduled, it should not be dropped.

There are important functional differences between the CoS operations of ACX and MX. All priority queues for ACX (low-latency, strict-high, high, med-high, and med-low) preempt the low-priority queues. PQs can be shaped to prevent the starving lower priority queues. By contrast, on MX TRIO-based platforms, queues operate in a guaranteed or excess region. Queues in profile (operating within the configured transmit rate) are serviced in the guaranteed region and are moved to the excess region when exceeding the transmit rate. Only strict-high priority has the authority to preempt all other priorities since it does not have an excess region and is therefore always guaranteed.

Let's now consider the behavior when the shaping rate is included for ACX7000 platforms on a priority queue. It is important to understand in case there is a combination of low-priority queues using transmit-rate with PQs using shaping-rate. The configured shaping-rate bandwidth applied to the queue is deducted from the port speed. This means the configured transmit-rate is based on the port speed minus shaping-rates configured.

The following table provides an example of how the shaping rate configuration influences the port speed, which is used to determine how much bandwidth is delegated based on the transmit rate percentage. The first row shows a port speed of 100GbE, and no queues are using the shaping rate. In this case, the 20% transmit rate on the low-priority queue allocates 20Gb.

The next row has 100GbE port speed, and one high-priority queue is given 30% of the port speed (shaping-rate). This brings the total remaining port speed to 70Gb. The 20% transmit rate for the low-priority queue is based on 70Gb, which accounts for 14Gb.

In the third row, the total shaping rate is 50% of the port speed, which allocates 10Gb based on the 20% transmit rate for the low-priority queue.

Finally, in row four, a total of 100% is allocated as shaping rate, leaving 0 for low priority. In all cases, unused bandwidth is available to be distributed across other queues. When bandwidth is available, the allocation is allowed to exceed the configured transmit rate. The shaping rate cannot be exceeded. Packets exceeding the configured shaping rate are dropped.

Table 12: Example Port Speed with Shaping-Rate

| Port Speed | High PQ Shaping-Rate | Med-High PQ Shaping-Rate | Updated Port Speed | Low Priority Transmit-Rate | Low Priority Bandwidth (Gbps) |
|------------|----------------------|--------------------------|--------------------|----------------------------|-------------------------------|
| 100GbE | - | - | 100 | 20% | 20 |
| 100GbE | 30% | - | 70 | 20% | 14 |
| 100GbE | 30% | 20% | 50 | 20% | 10 |
| 100GbE | 50% | 50% | 0 | 20% | 0 |

The previous generation ACX5448/ACX710 allows configurations of shaping-rate and transmit-rate at once. As a result, port speed does not change with the inclusion of the shaping rate. The ACX7000 series does not use this combination.

In the validated design, a 5G CoS model is created with four priority queues configured with shaping rate: Signaling, LLQ, Realtime, and Control. These queues carry the most critical traffic types. Under full congestion, these PQs are allowed to consume 80% of the total bandwidth, which is guaranteed.

The remaining 20% of the total port bandwidth is delegated across four WFQs. Although only 20% is allocated, recall that this is only if the PQ utilization is completely full. Leftover bandwidth will be available. For a moment, let's ignore the four priority queues and create a simple proportional bandwidth distribution model based on queues: High, Medium, Low, and Best-Effort. All these queues are treated the same with low-priority WFQ scheduling.

- High is given 40% of the total bandwidth
- Medium is given 30% of the total bandwidth
- Low is given 20% of the total bandwidth

- Best-Effort is given the remainder

The following table explains the WFQ CoS model, which exists in parallel to the PQ CoS model.

Table 13: Four WFQ Model

| Queue Name | Priority | Transmit Rate | Egress Queue |
|---------------|----------|-----------------|--------------|
| High | Low | 40% | EGQ3 |
| Medium | Low | 30% | EGQ3 |
| Low | Low | 20% | EGQ3 |
| Best Effort | Low | Remainder (10%) | EGQ3 |
| Total: | | 100% | |

All four queues are allowed to borrow from other queues if the bandwidth is available, with the weighted distribution (based on TR) favoring the queues in the order shown. High, Medium, and Low are all guaranteed some bandwidth, even if PQs are consuming all their configured rates. Only a small amount is remaining for Best-Effort, so this queue is not guaranteed.

This is one example of a potential 5G CoS model. Traffic patterns are quite different between operators. It is not possible to create one model that works perfectly for all situations. Thus, you must update the design to appropriately fit your use case.

A goal of the validation is to create a model realizing the O-RAN multiple PQ structure (transport core profile-B) with TSN Profile A, enabling PQs to preempt other PQs with LLQ supporting distinct dequeuing prioritization.

The following scheduler configuration is based on the ACX7000 family. Please contact your Juniper representative for the full configurations.

Schedulers

```

set class-of-service schedulers SC-SIGNALING shaping-rate percent 5
set class-of-service schedulers SC-SIGNALING buffer-size percent 5
set class-of-service schedulers SC-SIGNALING priority strict-high
set class-of-service schedulers SC-LLQ shaping-rate percent 40
set class-of-service schedulers SC-LLQ buffer-size percent 10
set class-of-service schedulers SC-LLQ priority low-latency
set class-of-service schedulers SC-REALTIME shaping-rate percent 30
set class-of-service schedulers SC-REALTIME buffer-size percent 20
set class-of-service schedulers SC-REALTIME priority medium-high
set class-of-service schedulers SC-HIGH transmit-rate percent 40
set class-of-service schedulers SC-HIGH buffer-size percent 30
set class-of-service schedulers SC-HIGH priority low
set class-of-service schedulers SC-CONTROL shaping-rate percent 5
set class-of-service schedulers SC-CONTROL buffer-size percent 5
set class-of-service schedulers SC-CONTROL priority high
set class-of-service schedulers SC-MEDIUM transmit-rate percent 30
set class-of-service schedulers SC-MEDIUM buffer-size percent 20
set class-of-service schedulers SC-MEDIUM priority low
set class-of-service schedulers SC-LOW transmit-rate percent 20
set class-of-service schedulers SC-LOW buffer-size percent 10
set class-of-service schedulers SC-LOW priority low
set class-of-service schedulers SC-BEST-EFFORT transmit-rate remainder
set class-of-service schedulers SC-BEST-EFFORT buffer-size remainder
set class-of-service schedulers SC-BEST-EFFORT priority low

```

Scheduler Map

```

set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-SIGNALING scheduler SC-SIGNALING
set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-LLQ scheduler SC-LLQ
set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-REALTIME scheduler SC-REALTIME
set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-HIGH scheduler SC-HIGH
set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-CONTROL scheduler SC-CONTROL
set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-MEDIUM scheduler SC-MEDIUM
set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-LOW scheduler SC-LOW
set class-of-service scheduler-maps SM-5G-SCHEDULER forwarding-class FC-BEST-EFFORT scheduler SC-BEST-EFFORT

```

The configuration at this point produces the following port-scheduling hierarchies. The VOQ ID is arbitrary and is assigned from a VOQ pool as part of the bundle association. Note the configuration results in the EGQ assignment.

Table 14: Custom Priority to EGQ Mapping

| Queue Priority | Forwarding Class | Queue | VOQ ID | Egress Queue Assignment |
|----------------|------------------|-------|--------|-------------------------|
| Low-Latency | FC-LLQ | 6 | 750 | EGQ1 |
| Strict-High | FC-SIGNALING | 7 | 751 | EGQ2 |
| High | FC-CONTROL | 3 | 747 | EGQ2 |
| Medium-High | FC-REALTIME | 5 | 749 | EGQ2 |
| Low | FC-HIGH | 4 | 748 | EGQ3 |
| Low | FC-MEDIUM | 2 | 746 | EGQ3 |
| Low | FC-LOW | 1 | 745 | EGQ3 |
| Low-Remainder | FC-BEST-EFFORT | 0 | 744 | EGQ3 |

Port Shaping

Queue shaping is explained in the previous section. Port Shaping is applied directly to the interface under the CoS hierarchy and adjusts the configured scheduler percentages based on the new port speed. For example, if a 10G shaper is configured on a 100GbE port, the reference port speed is 10G. A scheduler with a 50% transmit rate results in 5Gbps rather than 50Gbps.

Port Shaper

```
set class-of-service interface <x> shaping-rate <bps>
```

Rewrite Rules

Rewrite is performed at the egress path based on the protocol match. The following configurations are applicable to ACX, PTX, and MX platforms. While sending dual tags, the rewrite operations are

performed on the outer (S-TAG). It is generally desirable that inner C-TAG 802.1p bits are left unchanged and transmitted transparently.

802.1P Rewrite

```

set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-SIGNALING loss-priority low code-point 110
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-SIGNALING loss-priority high code-point 110
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-LLQ loss-priority low code-point 100
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-LLQ loss-priority medium-high code-point 100
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-LLQ loss-priority high code-point 100
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-REALTIME loss-priority low code-point 101
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-REALTIME loss-priority high code-point 101
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-HIGH loss-priority low code-point 011
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-HIGH loss-priority medium-high code-point 011
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-HIGH loss-priority high code-point 011
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-CONTROL loss-priority low code-point 111
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-CONTROL loss-priority high code-point 111
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-MEDIUM loss-priority low code-point 010
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-MEDIUM loss-priority medium-high code-point 010
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-MEDIUM loss-priority high code-point 010
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-LOW loss-priority low code-point 001
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-LOW loss-priority medium-high code-point 001
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-LOW loss-priority high code-point 001
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-BEST-EFFORT loss-priority low code-point 000
set class-of-service rewrite-rules ieee-802.1 RR-8021P forwarding-class FC-BEST-EFFORT loss-priority high code-point 000

```

```
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-SIGNALING loss-priority low code-point 110
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-SIGNALING loss-priority high code-point 110
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-LLQ loss-priority low code-point 100
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-LLQ loss-priority medium-high code-point
100
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-LLQ loss-priority high code-point 100
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-REALTIME loss-priority low code-point 101
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-REALTIME loss-priority high code-point 101
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-HIGH loss-priority low code-point 011
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-HIGH loss-priority medium-high code-point
011
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-HIGH loss-priority high code-point 011
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-CONTROL loss-priority low code-point 111
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-CONTROL loss-priority high code-point 111
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-MEDIUM loss-priority low code-point 010
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-MEDIUM loss-priority medium-high code-
point 010
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-MEDIUM loss-priority high code-point 010
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-LOW loss-priority low code-point 001
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-LOW loss-priority medium-high code-point
001
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-LOW loss-priority high code-point 001
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-BEST-EFFORT loss-priority low code-point
000
set class-of-service rewrite-rules exp RR-MPLS forwarding-class FC-BEST-EFFORT loss-priority high code-point
000
```

DSCP IPv4 Rewrite

```
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-SIGNALING loss-priority low code-point cs6
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-SIGNALING loss-priority high code-point
cs5
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-LLQ loss-priority low code-point af41
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-LLQ loss-priority medium-high code-point
af42
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-LLQ loss-priority high code-point af43
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-REALTIME loss-priority low code-point ef
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-REALTIME loss-priority high code-point ef
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-HIGH loss-priority low code-point af31
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-HIGH loss-priority medium-high code-point
af32
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-HIGH loss-priority high code-point af33
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-CONTROL loss-priority low code-point cs7
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-CONTROL loss-priority high code-point cs7
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-MEDIUM loss-priority low code-point af21
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-MEDIUM loss-priority medium-high code-
point af22
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-MEDIUM loss-priority high code-point af23
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-LOW loss-priority low code-point af11
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-LOW loss-priority medium-high code-point
af12
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-LOW loss-priority high code-point af13
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-BEST-EFFORT loss-priority low code-point
be
set class-of-service rewrite-rules dscp RR-DSCP forwarding-class FC-BEST-EFFORT loss-priority high code-point
be
```

DSCP IPv6 Rewrite

```
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-SIGNALING loss-priority low
code-point cs6
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-SIGNALING loss-priority high
code-point cs5
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-LLQ loss-priority low code-
point af41
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-LLQ loss-priority medium-high
code-point af42
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-LLQ loss-priority high code-
point af43
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-REALTIME loss-priority low code-
point ef
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-REALTIME loss-priority high
code-point ef
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-HIGH loss-priority low code-
point af31
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-HIGH loss-priority medium-high
code-point af32
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-HIGH loss-priority high code-
point af33
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-CONTROL loss-priority low code-
point cs7
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-CONTROL loss-priority high code-
point cs7
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-MEDIUM loss-priority low code-
point af21
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-MEDIUM loss-priority medium-
high code-point af22
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-MEDIUM loss-priority high code-
point af23
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-LOW loss-priority low code-
point af11
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-LOW loss-priority medium-high
code-point af12
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-LOW loss-priority high code-
point af13
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-BEST-EFFORT loss-priority low
code-point be
set class-of-service rewrite-rules dscp-ipv6 RR-DSCP-IPV6 forwarding-class FC-BEST-EFFORT loss-priority high
code-point be
```

Interface Class of Service

Once the CoS parameters are established, the application of the configuration is made to the interface itself under CoS hierarchy. Recall the three styles of classification: Behavior Aggregate, Fixed, and Multifield Classifier. Only BA and Fixed styles are applicable under this configuration stanza. BA matches on received layer 2 802.1p bits and/or layer 3 DSCP. Fixed classification maps all traffic received on the interface to a single forwarding class. Wildcard match conditions are used for interfaces and units where applicable. From Junos OS Evolved Release 23.2R1, multiple classifiers and rewrite rules are allowed on the same interface for ACX7000 platforms.

Scheduler Map

```
set class-of-service interfaces <ifd> scheduler-map SM-5G-SCHEDULER
BA Classifier
set class-of-service interfaces <ifd> unit <ifl> classifiers exp CL-MPLS
set class-of-service interfaces <ifd> unit <ifl> classifiers dscp CL-DSCP
set class-of-service interfaces <ifd> unit <ifl> classifiers ieee-802.1 CL-8021P
Fixed Classifier
set class-of-service interfaces <ifd> forwarding-class <fc name>
Rewrite Rules
set class-of-service interfaces <ifd> unit <ifl> rewrite-rules exp RR-MPLS
set class-of-service interfaces <ifd> unit <ifl> rewrite-rules dscp RR-DSCP
set class-of-service interfaces <ifd> unit <ifl> rewrite-rules ieee-802.1 RR-8021P
```

BA Classifier

```
set class-of-service interfaces <ifd> unit <ifl> classifiers exp CL-MPLS
set class-of-service interfaces <ifd> unit <ifl> classifiers dscp CL-DSCP
set class-of-service interfaces <ifd> unit <ifl> classifiers ieee-802.1 CL-8021P
```

Fixed Classifier

```
set class-of-service interfaces <ifd> forwarding-class <fc name>
```

Rewrite Rules

```
set class-of-service interfaces <ifd> unit <ifl> rewrite-rules exp RR-MPLS
set class-of-service interfaces <ifd> unit <ifl> rewrite-rules dscp RR-DSCP
set class-of-service interfaces <ifd> unit <ifl> rewrite-rules ieee-802.1 RR-8021P
```

Host Outbound Traffic

ACX7000 series platforms support host-outbound classification from Junos OS Evolved Release 23.3R1.

```
set class-of-service host-outbound-traffic forwarding-class <FC-NAME>
```

Buffer Management

ACX7000 series platforms leverage VOQ ingress buffer machinery based on guaranteed and dynamic buffers. Dynamic buffers are elastic in nature and leverage Fair Adaptive Dynamic Threshold (FADT) as the algorithm to manage the shared buffer pool.

When the default CoS is used, the ACX7K allocates the buffer as:

Table 15: Default Buffer Allocation

| Queue | Buffer |
|--------------|-----------|
| 0 | 95% |
| 3 | 5% |
| Other Queues | * Minimum |

** Based on port speed.*

The minimum buffer that can be allocated is based on the port speed:

- 10G = 2048
- 25G = 4096

- 40G = 4096
- 50G = 4096
- 100G = 8192

Buffer usage can be monitored primarily with these commands:

- show interface voq (from CLI)
- show cos voq buffer-occupancy ifd (from VTY)

Buffer Allocation

The CoS buffer configuration outlined in the ["Scheduling" on page 49](#) section proposes the following buffer allocation per queue for the ACX7000 platform.

```
set class-of-service schedulers SC-SIGNALING buffer-size percent 5
set class-of-service schedulers SC-LLQ buffer-size percent 10
set class-of-service schedulers SC-REALTIME buffer-size percent 20
set class-of-service schedulers SC-HIGH buffer-size percent 30
set class-of-service schedulers SC-CONTROL buffer-size percent 5
set class-of-service schedulers SC-MEDIUM buffer-size percent 20
set class-of-service schedulers SC-LOW buffer-size percent 10
set class-of-service schedulers SC-BEST-EFFORT buffer-size remainder
```

The following table explains the buffer allocation calculated for each queue utilized in the validation. For example, a 100GbE port with 1250KB of dedicated buffer can be used.

Table 16: Buffer Allocation Results

| Queue | Rate | Buffer Size | Priority | Calculation | Buffer |
|-------|-------------------|-------------|-------------|----------------|--------|
| Q7 | 5% shaped rate | 5% | strict-high | $1250000*0.05$ | 62500 |
| Q6 | 40% shaped rate | 10% | low-latency | $1250000*0.1$ | 125000 |
| Q5 | 30% shaped rate | 20% | medium-high | $1250000*0.2$ | 250000 |
| Q4 | 40% transmit rate | 30% | low | $1250000*0.3$ | 375000 |

Table 16: Buffer Allocation Results (Continued)

| Queue | Rate | Buffer Size | Priority | Calculation | Buffer |
|-------|-------------------|-------------|----------|----------------|--------|
| Q3 | 5% shaped rate | 5% | high | $1250000*0.05$ | 62500 |
| Q2 | 30% transmit rate | 20% | low | $1250000*0.2$ | 250000 |
| Q1 | 20% transmit rate | 10% | low | $1250000*0.1$ | 125000 |
| Q0 | remainder | remainder | low | $1250000*0$ | *8192 |

* remaining buffer is 0%, so minimum buffer programmed

Test Bed Configuration

Contact your Juniper account representative to obtain the full archive of the test bed configuration used for this JVD.

Results Summary and Analysis

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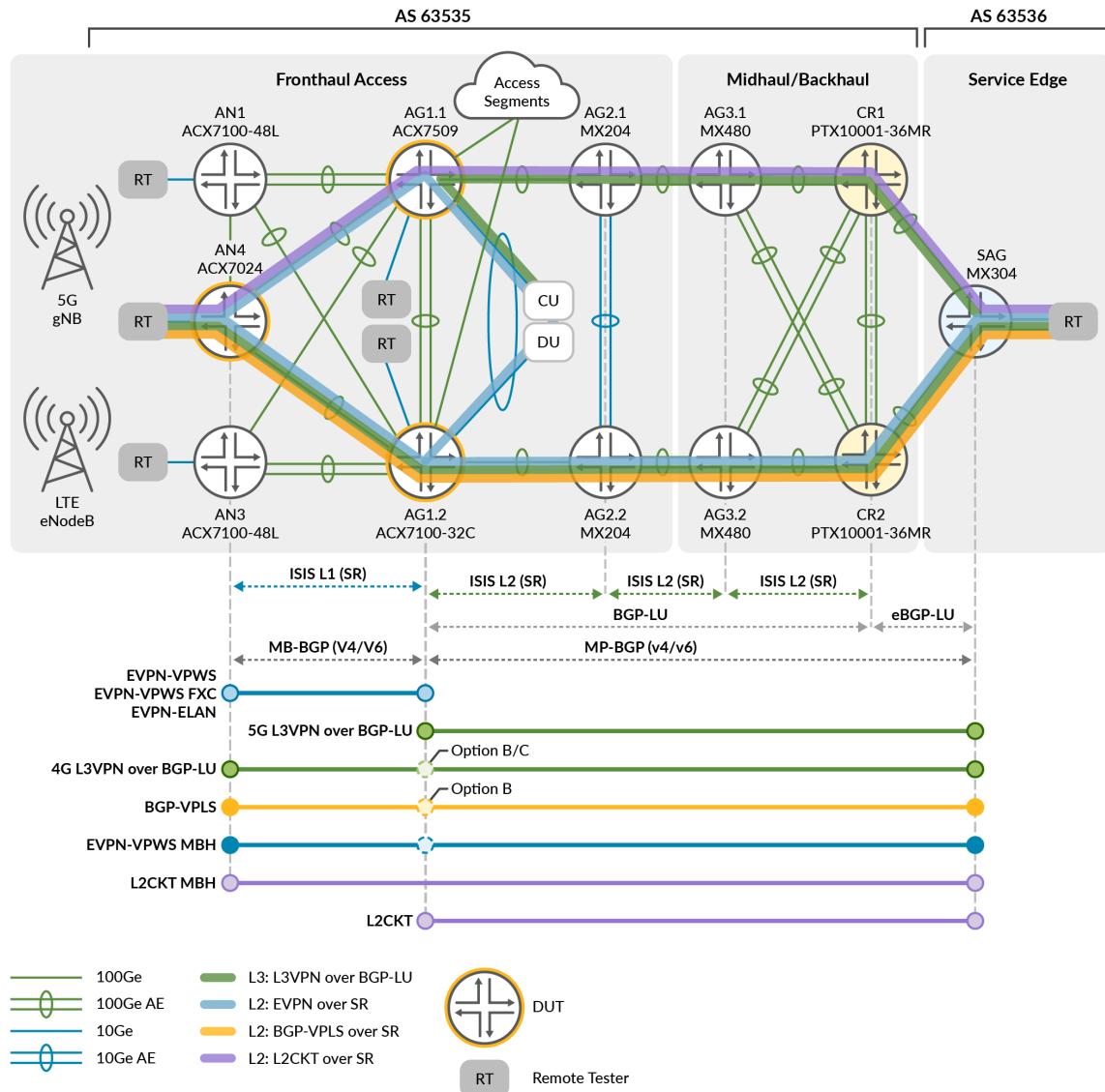
CoS for LLQ JVD extends previous solutions with the focused delivery of differentiated services supporting ultra-low latency workloads. 5G O-RAN is an open and disaggregated RAN architecture that aims to enable interoperability, flexibility, and innovation among different vendors and operators.

The fronthaul is the most demanding segment of the xHaul architecture, requiring high performance and functionality. Comprehensive QoS in the 5G network architecture is mandatory to ensure reliable and efficient performance across diverse applications and services. QoS mechanisms prioritize traffic, manage bandwidth, and preserve latency budgets, ensuring critical applications receive the necessary resources and maintain overall network performance and user experience.

During the validation process, a robust solution is demonstrated for 5G xHaul transport infrastructure using Seamless MPLS with Segment Routing and EVPN-VPWS, EVPN-FXC, EVPN-ELAN, VPLS, L2Circuit, and L3VPN services. The JVD generates reasonable multi-vector scale of Layer 2 and Layer 3 connectivity services, meeting the expectations of Mobile Network Operators (MNOs) and Metropolitan Area Network (MAN) operators for real network deployments while satisfying strict SLA requirements.

ACX7024 (CSR), ACX7100-32C (HSR), and ACX7509 (HSR) as the DUT routers have each successfully passed all 581 test cases curated to support the reference architecture.

Figure 26: Reference Architecture



Class of Service Operations

The CoS model differs between operators. The model created in the validation of this profile is shared in the following table. However, this model might be modified to meet the requirements for a given implementation based on the characteristics of the operator's traffic patterns and goals. In the table, the higher priority queues use a shaping rate (SR), and low-priority queues use a transmit rate (TR). For more information, see the "Scheduling" on page 49^{OBJ} section.

Table 17: Validated Scheduling Profile

| Forwarding Classes | Scheduling Parameters | | | | Classification and Rewrite | | | Traffic Profile | |
|--------------------|-----------------------|-------------|---------|--------|----------------------------|-----------------------|-----|---|--|
| | Queue | Priority | BW Rate | Buffer | 802.1p | DSCH | EXP | Type | |
| FC-SIGNALING | 7 | Strict-High | 5% SR | 5% | 110 | CS5, CS6 | 110 | OAM aggressive timers, O-RAN/3GPP C-plane | |
| FC-LLQ | 6 | Low-Latency | 40% SR | 10% | 100 | CS4, AF41, AF42, AF43 | 100 | CPRI RoE, eCPRI C/U-Plane ≤ 2000 bytes | |
| FC-REALTIME | 5 | Medium-High | 30% SR | 20% | 101 | EF | 101 | 5QI/QCI Group 1 low-latency U-plane, low-latency business. Interactive video, low latency voice | |
| FC-HIGH | 4 | Low | 40% TR | 30% | 011 | CS3, AF31, AF32, AF33 | 011 | 5QI/QCI Group 2 medium latency U-plane data | |
| FC-CONTROL | 3 | High | 5% SR | 5% | 111 | CS7 | 111 | Network control: OAM relaxed timers, IGP, BGP, PTP aware mode | |
| FC-MEDIUM | 2 | Low | 30% TR | 20% | 010 | CS2, AF21, AF22, AF23 | 010 | 5QI/QCI Group 3 remainder GBR U-plane guaranteed business data. Video-on-demand. O-RAN/3GPP M-plane, e.g., eCPRI M-plane, other management, software upgrades | |
| FC-LOW | 1 | Low | 20% TR | 10% | 001 | CS1, AF11, AF12, AF13 | 001 | high latency, guaranteed low-priority data | |

Table 17: Validated Scheduling Profile (*Continued*)

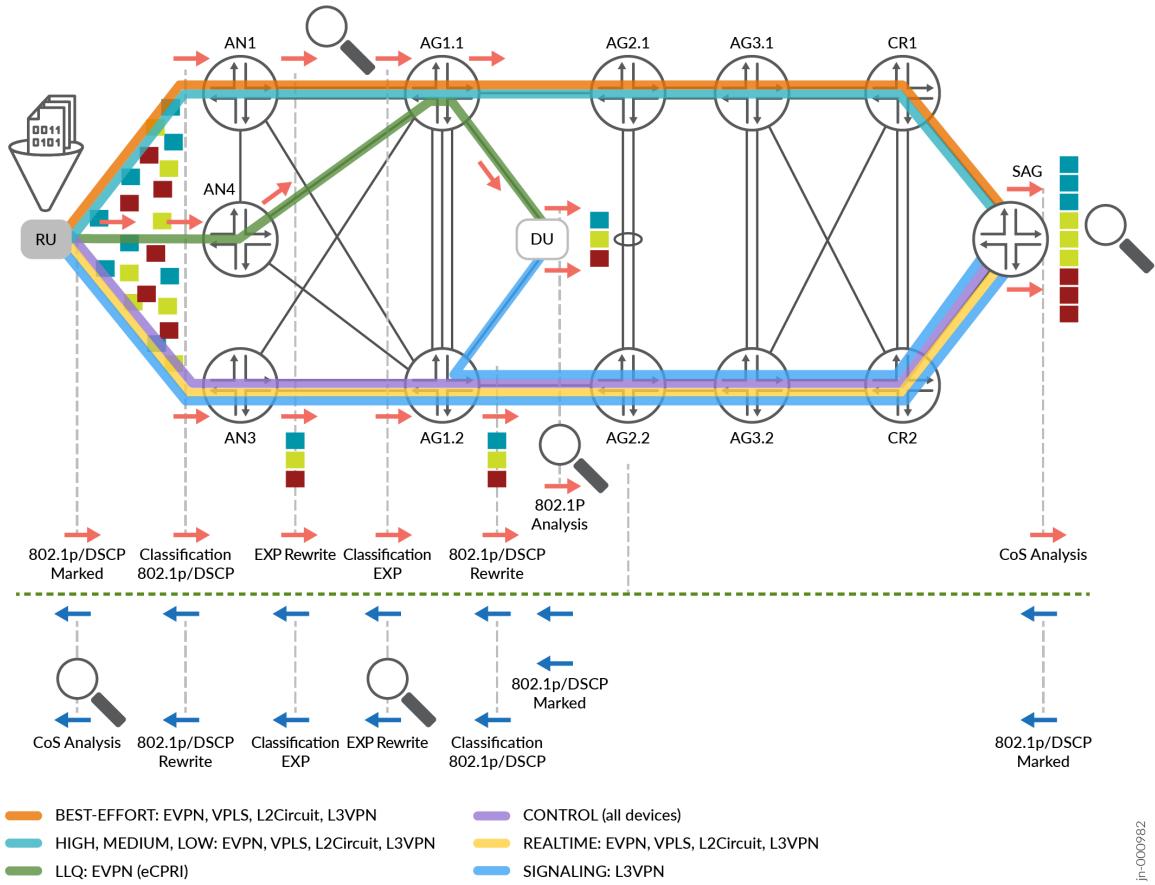
| Forwarding Classes | Scheduling Parameters | | | | Classification and Rewrite | | | Traffic Profile | |
|--------------------|-----------------------|----------|-----------|-----------|----------------------------|------|-----|--|--|
| | Queue | Priority | BW Rate | Buffer | 802.1p | DSCP | EX | Type | |
| FC-BEST-EFFORT | 0 | Low | remainder | remainder | 000 | BE | 000 | 5QI/QCI Group 4 – remainder non-GBR U-plane data | |

The validation includes three styles of ingress classification: fixed, behavior aggregate, and multifield classifier. BA is *packet-based*, where flows are pre-marked with Layer 3 DSCP, Layer 2 802.1Q Priority Code Points (PCP), or MPLS EXP. Fixed classification is *context-based*, where all traffic arriving on a specific interface is mapped to one forwarding class. Multifield (MF) Classification allows matching criteria within packet fields and mapping to one or more forwarding classes.

A crucial goal of the JVD is the validation of CoS functional operations, ensuring that traffic is properly matched (classified) and placed into an appropriate forwarding class to be scheduled for egress transmission. O-RAN or 3GPP proposes a minimum of six queues and a maximum of eight queues per interface. For the JVD, eight queues and associated forwarding classes are used to accommodate the traffic scheme requirements. In supporting *Profile-B*, multiple PQs are used, which are shaped (PIR) to prevent starving lower priority queues. In parallel, four WFQs are created with low priority and transmit rate bandwidth assigned by a percentage of the remaining port speed. A final low-priority queue is given the remainder without guaranteed bandwidth.

At egress, DSCP, 802.1p or EXP codepoints, and packet loss priorities (PLP) are rewritten based on the assigned forwarding class and rewrite-rule instruction. ACX supports rewriting the outer tag, which is the default and commonly preferred to preserve the inner (C-TAG) 802.1p bits for transparent transmission.

Figure 27: Class of Service Operations



The following table summarizes the three classification and rewrite operations, resulting in the preservation of the CoS priority bits. Services include EVPN-VPWS, EVPN Flexible Cross Connect (FXC), L2Circuit, VPLS, EVPN-ELAN, EVPN IRB Virtual Gateway Address (VGA) active-active ESI to L3VPN, and Layer 2 Bridge Domain with Anycast IRB to L3VPN. The table also share information about the services validated for each classification type. The final section of [Table 18 on page 69](#) involves test scenarios where specific traffic types (eCPRI, PTP, and OAM) are sent, and multifield classifiers are used to match the packet header fields to be given desired priority.

In addition, the validation ensures that the priority hierarchies are honored. Irrespective of whether fixed, BA, or multifield classifiers are used, traffic is validated to be mapped to the proper priority queue. The checkmarks under LLQ (low-latency), SH (strict-high), H (high), MH (medium-high), and L (low) confirm that the scheduler rates defined in ["Table 17 " on page 66](#) are honored. This ensures traffic rates can be guaranteed (CIR) regardless of the priority level with proper configuration. The dashed lines represent scenarios not applicable to the test conditions.

For example, Fixed classification is applied at the ingress interface irrespective of packet codepoints received. Traffic is verified to traverse the correct egress queue (Queue Match) and is confirmed again with the correct EXP ingress BA classification at the next hop.

In the table, Scheduler Rates Honored confirms the expected bandwidth allocation of queues utilizing shaping or transmit rate behaviors as outlined previously. Codepoints Rewritten validates whether an expected DSCP, 802.1p, or EXP rewrite operation is successful. This is dependent on the protocol type at the position of the network. Bits Preserved confirms that the expected codepoints of the outer or inner tag are recorded properly. Packet captures are taken to confirm.

All CoS functional test case scenarios passed successfully during the execution. For more information, see the Test Report or contact your Juniper representative.

Table 18: CoS Summarized Results

| Traffic Scenario | Ingress Classification Mapped to FC | | | Scheduler Rates Honored | | | | | | Codepoints Rewritten | | | Bits Preserved |
|----------------------------|-------------------------------------|--|-----|-------------------------|----|----|----|----|--------|----------------------|-----|-----|----------------|
| Fixed Classifier | Queue Match | | EXP | LLQ | SH | H | MH | L | 802.1p | DSCP | EXP | E2E | |
| EVPN - VPWS | ✓ | | ✓ | ✓ | -- | -- | -- | -- | ✓ | -- | ✓ | ✓ | |
| EVPN -FXC | ✓ | | ✓ | ✓ | -- | -- | -- | -- | ✓ | -- | ✓ | ✓ | |
| L2Circuit | ✓ | | ✓ | -- | -- | -- | -- | ✓ | ✓ | -- | ✓ | ✓ | |
| EVPN IRB VGA with L3VPN MH | ✓ | | ✓ | -- | -- | -- | ✓ | -- | ✓ | -- | ✓ | ✓ | |

Table 18: CoS Summarized Results (*Continued*)

| Traffic Scenario | Ingress Classification Mapped to FC | | | Scheduler Rates Honored | | | | | Codepoints Rewritten | | | Bits Preserved |
|---|-------------------------------------|------|-----|-------------------------|----|----|----|--------|----------------------|------|-----|----------------|
| Fixed Classifier | Queue Match | EXP | LLQ | SH | H | MH | L | 802.1p | DSCP | EXP | E2E | |
| BD IRB anyca st static MAC /IP with L3VP N | ✓ | | ✓ | -- | -- | -- | ✓ | -- | ✓ | -- | ✓ | ✓ |
| BA Classifier | 802.1p | DSCP | EXP | LLQ | SH | H | MH | L | 802.1p | DSCP | EXP | E2E |
| EVPN - VPWS | ✓ | -- | ✓ | -- | ✓ | ✓ | ✓ | ✓ | ✓ | -- | ✓ | ✓ |
| EVPN -FXC | ✓ | -- | ✓ | -- | ✓ | ✓ | ✓ | ✓ | ✓ | -- | ✓ | ✓ |
| EVPN - ELAN | ✓ | -- | ✓ | -- | ✓ | ✓ | ✓ | ✓ | ✓ | -- | ✓ | ✓ |

Table 18: CoS Summarized Results (*Continued*)

| Traffic Scenario | Ingress Classification Mapped to FC | | | Scheduler Rates Honored | | | | | Codepoints Rewritten | | | Bits Preserved |
|---|-------------------------------------|----|-----|-------------------------|----|----|----|---|----------------------|------|-----|----------------|
| Fixed Classifier | Queue Match | | EXP | LLQ | SH | H | MH | L | 802.1p | DSCP | EXP | E2E |
| EVPN IRB anycast with L3VPN MH | ✓ | -- | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | -- | ✓ | ✓ |
| L3VPN | -- | ✓ | ✓ | -- | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ |
| VPLS | ✓ | -- | ✓ | -- | -- | -- | -- | ✓ | -- | -- | ✓ | ✓ |
| BD IRB anycast static MAC/IP with L3VPN | ✓ | -- | ✓ | -- | ✓ | ✓ | ✓ | ✓ | ✓ | -- | ✓ | ✓ |
| Multifield Classifier | Packet Match | | EXP | LLQ | SH | H | MH | L | 802.1p | DSCP | EXP | E2E |

Table 18: CoS Summarized Results (*Continued*)

| Traffic Scenario | Ingress Classification Mapped to FC | | | Scheduler Rates Honored | | | | | Codepoints Rewritten | | | Bits Preserved |
|---------------------|-------------------------------------|--|-----|-------------------------|----|----|----|----|----------------------|------|-----|----------------|
| Fixed Classifier | Queue Match | | EXP | LLQ | SH | H | MH | L | 802.1p | DSCP | EXP | E2E |
| Traffic Type: eCPRI | ✓ | | ✓ | ✓ | -- | -- | -- | -- | -- | -- | ✓ | ✓ |
| Traffic Type: PTP | ✓ | | ✓ | ✓ | -- | -- | -- | -- | -- | -- | ✓ | ✓ |
| Traffic Type: OAM | ✓ | | ✓ | ✓ | -- | -- | -- | -- | -- | -- | ✓ | ✓ |

As shown above, BA services included multiple traffic classes with checkmarks across multiple priority queues. In this condition, there is contention for bandwidth resources. By including a shaping rate to prevent priority queues from starving the low-priority queues, you can ensure bandwidth reserves are available to service all queues.

Multifield classifiers are used to match specific traffic based on packet header information to be mapped into a desired forwarding class. For example, eCPRI traffic is identified by EtherType 0xAEFE and PTP by EtherType 0x88F7. For more information about multifield classifier match conditions and results, including packet captures, see the Test Report.

For brevity, only one low-priority (L) queue is shown in the table above. However, the complete model includes four low-priority queues. After priority queues with shaping rate deduct the expected bandwidth, the remaining bandwidth defines the port speed that is utilized for the low priority queues using transmit rate percentages. The percentage allocation remains consistent based on this updated port speed.

Congestion Scenarios

The validation includes a variety of congestion scenarios, which are outlined in the ["Test Objectives" on page 17](#) section. Congestion constitutes one or more conditions where traffic exceeds the configured scheduler transmit rate, shaped rate, or port speed and results in an expected traffic loss. The major objective is to ensure critical priority traffic is uninterrupted even during periods of congestion and that during congestion, high-priority delay-sensitive traffic is given preference.

During key congestion events, the following observations are recorded. For more information, see the full test report:

- Low-latency queue (FC-LLQ) is serviced ahead of strict-high, high, medium-high, medium-low, and low priority queues and given preferential treatment to preserve the latency budget. FC-LLQ does not exceed the defined allowance when a shaping rate is configured.
- Strict-high queue (FC-SIGNALING) is serviced ahead of high, medium-high, medium-low, and low queues. FC-SIGNALING does not exceed the defined allowance when a shaping rate is configured.
- High queue (FC-CONTROL) is serviced ahead of medium-high, medium-low, and low queues. FC-CONTROL does not exceed the defined allowance when a shaping rate is configured.
- Medium-high queue (FC-REALTIME) is serviced ahead of medium-low and low queues. FC-REALTIME does not exceed the defined allowance when a shaping rate is configured.
- Low-priority queues (FC-HIGH, FC-MEDIUM, and FC-LOW) are serviced based on configured transmit rate percentages of the remaining port speed. Low-priority queues are serviced as WFQ when operating in excess of the transmit rate and bandwidth is available.
- Low-priority remainder queue (FC-BEST-EFFORT) is serviced only when unused bandwidth is available based on the intentionality of the scheduler configuration.
- During high congestion periods, queues operating within their configured bandwidth (in-profile) are guaranteed without packet drops.
- Critical eCPRI or ORAN traffic flows are assigned to priority queues and meet the given SLAs.
- Scheduler percentages correctly inherit the configured port-shaper as port speed.
- Queue shaping rate is deducted from total port speed with transmit-rates applied to the remaining bandwidth.
- Priority hierarchies are honored across and within VPN services that share common links.
- Traffic is transmitted up to 99% port speed under normal conditions where multi-level priorities are validated.
- Latency results are always measured in microseconds (μs).

Latency Budget Validation

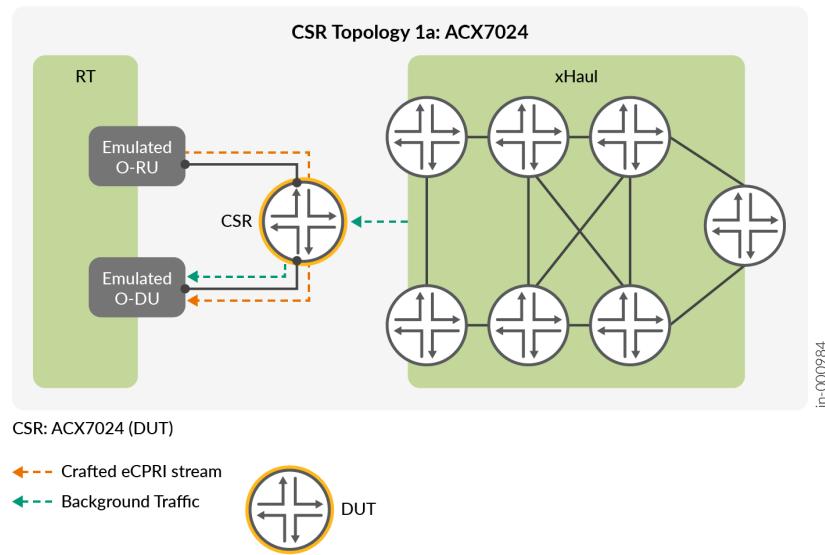
5G xHaul infrastructure defines strict latency budgets, particularly in the fronthaul segment. O-RAN WG9 technical specification (O-RAN.WG9.XTRP-REQ-v01.00) for transport requirements provides guidelines for a maximum one-way latency budget of $100\mu\text{s}$ for the fronthaul. This budget includes latency incurred by the fiber runs at approximately $4.9\mu\text{s}/\text{km}$ and transit nodes. The objective of the JVD is to deliver a solution with latency performance within the acceptable range from O-RU to O-DU, approximately $\leq 10\mu\text{s}$ per node. The ACX7000 platforms selected for CSR and HSR roles are shown to consistently achieve an average transit latency of $\leq 4\text{-}6\mu\text{s}$. In addition, the ability to preserve an objective latency budget for priority flows during conditions of congestion. This is a significant paradigm shift from the requirements of earlier 4G architectures.

When measuring a traffic delay, there should not be any traffic dropping from the queue under test since this invalidates the goal of capturing transit latency. The best measure of individual device performance is shown by the standalone topology 1. Additional topologies provide key insight into the overall network performance. Three permutations for each CSR and HSR device type are validated.

Topology 1a: ACX7024

Topology 1a ([Figure 28 on page 75](#)) is used to validate the performance of the ACX7024 platform in the CSR role. The test equipment emulates O-RU and O-DU devices and does not include self-latency. For each scenario, the results are provided across different frame sizes with either continuous or burst traffic patterns. Latency is measured in microseconds (μs).

Figure 28: Topology 1a ACX7024 Standalone DUT



The following table describes the ACX7024 performance when traffic is sent only to the low-latency queue without background traffic. No queues are congested during this time.

Table 19: ACX7024 Low-Latency without Background Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 5.83 | 6.18 | 9.54 | 64 | continuous | 10G |
| FC-LLQ | Low-Latency | 5.78 | 6.13 | 9.55 | 512 | continuous | 10G |
| FC-LLQ | Low-Latency | 5.09 | 5.31 | 6.61 | 1500 | continuous | 10G |

The following table includes continuous background traffic up to 99%-line rate with encapsulation overhead considerations for Topology 1a with ACX7024. The performance of the low-latency queue (FC-LLQ) is measured in conjunction with multi-level priority queues. From the results, you can see that LLQ is given preferential treatment for latency constraints. In these scenarios, lower priority queues are also maintained below 10μs average latency goal.

Table 20: ACX7024 Low-Latency with Continuous Background Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 5.09 | 5.43 | 12.95 | 64 | continuous | 10G |
| FC-SIGNALING | Strict-High | 5.09 | 6.94 | 52.72 | 64 | continuous | 10G |
| FC-CONTROL | High | 5.01 | 6.36 | 20.40 | 64 | continuous | 10G |
| FC-REALTIME | Medium-High | 5.08 | 6.24 | 24.66 | 64 | continuous | 10G |
| FC-HIGH | Low | 5.09 | 6.68 | 39.37 | 64 | continuous | 10G |
| FC-MEDIUM | Low | 5.09 | 7.05 | 51.91 | 64 | continuous | 10G |
| FC-LOW | Low | 5.09 | 6.94 | 52.72 | 64 | continuous | 10G |
| | | | | | | | |
| FC-LLQ | Low-Latency | 5.77 | 5.96 | 10.10 | 512 | continuous | 10G |
| FC-SIGNALING | Strict-High | 5.77 | 6.10 | 12.40 | 512 | continuous | 10G |
| FC-CONTROL | High | 5.76 | 6.26 | 17.33 | 512 | continuous | 10G |
| FC-REALTIME | Medium-High | 5.78 | 6.43 | 21.05 | 512 | continuous | 10G |
| FC-HIGH | Low | 5.78 | 6.40 | 19.31 | 512 | continuous | 10G |
| FC-MEDIUM | Low | 5.78 | 6.10 | 25.61 | 512 | continuous | 10G |
| FC-LOW | Low | 5.78 | 6.32 | 23.38 | 512 | continuous | 10G |

Table 20: ACX7024 Low-Latency with Continuous Background Traffic *(Continued)*

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 5.83 | 6.06 | 11.23 | 1500 | continuous | 10G |
| FC-SIGNALING | Strict-High | 5.84 | 6.08 | 13.48 | 1500 | continuous | 10G |
| FC-CONTROL | High | 5.84 | 6.10 | 16.25 | 1500 | continuous | 10G |
| FC-REALTIME | Medium-High | 5.81 | 6.36 | 22.66 | 1500 | continuous | 10G |
| FC-HIGH | Low | 5.87 | 6.07 | 34.26 | 1500 | continuous | 10G |
| FC-MEDIUM | Low | 5.87 | 6.06 | 30.42 | 1500 | continuous | 10G |
| FC-LOW | Low | 5.87 | 6.07 | 32.43 | 1500 | continuous | 10G |

The following table includes burst traffic type rather than continuous for Topology 1a with ACX7024, measuring the performance of the low-latency queue (FC-LLQ) in conjunction with multi-level priority queues. From the results, you can see that LLQ is given preferential treatment for latency constraints. In these scenarios, lower priority queues are also maintained well below the 10μs average latency goal.

Table 21: ACX7024 Low-Latency with Burst Background Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 5.82 | 5.91 | 6.38 | 64 | burst | 10G |
| FC-SIGNALING | Strict-High | 5.82 | 5.99 | 9.69 | 64 | burst | 10G |
| FC-CONTROL | High | 5.86 | 5.91 | 7.61 | 64 | burst | 10G |

Table 21: ACX7024 Low-Latency with Burst Background Traffic (Continued)

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-REALTIME | Medium-High | 5.83 | 5.91 | 8.18 | 64 | burst | 10G |
| FC-HIGH | Low | 5.87 | 5.91 | 6.71 | 64 | burst | 10G |
| FC-MEDIUM | Low | 5.87 | 5.91 | 6.46 | 64 | burst | 10G |
| FC-LOW | Low | 5.87 | 5.91 | 9.13 | 64 | burst | 10G |
| <hr/> | | | | | | | |
| FC-LLQ | Low-Latency | 5.77 | 5.80 | 6.21 | 512 | burst | 10G |
| FC-SIGNALING | Strict-High | 5.77 | 5.87 | 11.39 | 512 | burst | 10G |
| FC-CONTROL | High | 5.69 | 5.87 | 7.72 | 512 | burst | 10G |
| FC-REALTIME | Medium-High | 5.77 | 5.87 | 9.55 | 512 | burst | 10G |
| FC-HIGH | Low | 5.77 | 5.87 | 6.45 | 512 | burst | 10G |
| FC-MEDIUM | Low | 5.77 | 5.87 | 6.69 | 512 | burst | 10G |
| FC-LOW | Low | 5.77 | 5.87 | 6.64 | 512 | burst | 10G |
| <hr/> | | | | | | | |
| FC-LLQ | Low-Latency | 5.08 | 5.12 | 6.31 | 1500 | burst | 10G |
| FC-SIGNALING | Strict-High | 5.08 | 5.30 | 11.10 | 1500 | burst | 10G |

Table 21: ACX7024 Low-Latency with Burst Background Traffic (Continued)

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|-------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-CONTROL | High | 5.00 | 5.30 | 8.95 | 1500 | burst | 10G |
| FC-REALTIME | Medium-High | 5.08 | 5.30 | 8.72 | 1500 | burst | 10G |
| FC-HIGH | Low | 5.08 | 5.30 | 6.79 | 1500 | burst | 10G |
| FC-MEDIUM | Low | 5.08 | 5.30 | 8.35 | 1500 | burst | 10G |
| FC-LOW | Low | 5.08 | 5.30 | 6.57 | 1500 | burst | 10G |

For the next table, observe the impact on the LLQ when creating congestion to the point of traffic discarding in the oversubscribed strict-high queue. The intentionality of the LLQ design prevents such scenarios from disrupting the low-latency queue. From the results, you can observe that the average latency budget is $\leq 6\mu\text{s}$ and well below the goal of $10\mu\text{s}$.

Table 22: ACX7024 Low-Latency with Congestion Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Congested Queue |
|------------|----------------|------------------|------------------|------------------|------------|-----------------|-----------------|
| FC-LLQ | Low-Latency | 5.83 | 6.17 | 12.44 | 64 | continuous | Strict-High |
| FC-LLQ | Low-Latency | 5.78 | 6.02 | 10.95 | 512 | continuous | Strict-High |
| FC-LLQ | Low-Latency | 5.09 | 5.26 | 11.34 | 1500 | continuous | Strict-High |
| FC-LLQ | Low-Latency | 5.09 | 5.16 | 11.39 | 64 | burst | Strict-High |
| FC-LLQ | Low-Latency | 5.78 | 5.99 | 11.00 | 512 | burst | Strict-High |

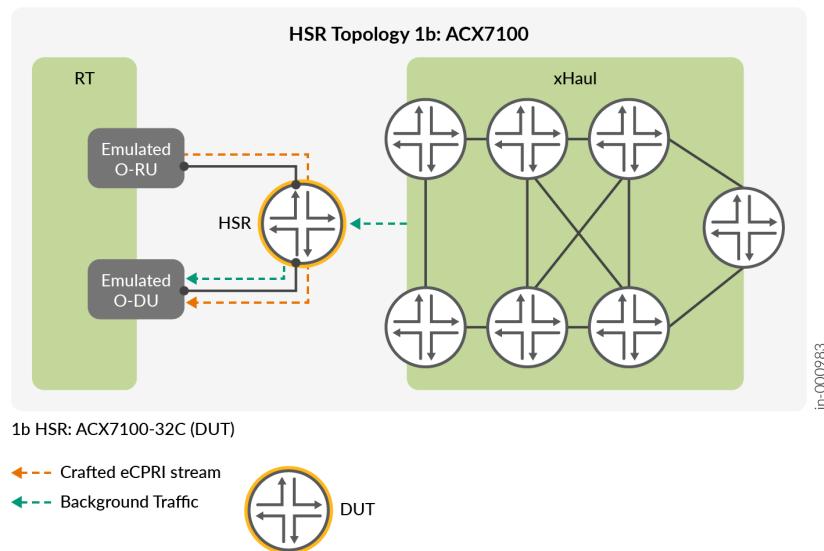
Table 22: ACX7024 Low-Latency with Congestion Traffic (*Continued*)

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Congested Queue |
|------------|----------------|------------------|------------------|------------------|------------|-----------------|-----------------|
| FC-LLQ | Low-Latency | 5.84 | 6.00 | 12.44 | 1500 | burst | Strict-High |

Topology 1b: ACX7100-32C

In Topology 1b, the same test scenarios are executed while moving to the HSR role with the examination of ACX7100-32C performance. Once again, the test equipment emulates O-RU and O-DU without self-latency. There is no physical CSR. In this case, the performance of the single HSR device is captured. For each scenario, the results are provided across different frame sizes with either continuous or burst traffic patterns. Latency is measured in microseconds (μs).

Figure 29: Topology 1b ACX7100-32C Standalone DUT



The following table describes the Topology 1b ACX7100-32C performance when traffic is sent only to the LLQ without background traffic. No queues are congested during this time.

Table 23: ACX7100-32C Low-Latency without Background Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|------------|----------------|---------------------|---------------------|---------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 4.35 | 4.44 | 7.82 | 64 | continuous | 10G |
| FC-LLQ | Low-Latency | 5.05 | 5.18 | 5.57 | 512 | continuous | 10G |
| FC-LLQ | Low-Latency | 4.68 | 4.88 | 5.86 | 1500 | continuous | 10G |

The following table includes continuous background traffic up to 99%-line rate with encapsulation overhead considerations for Topology 1b with ACX7100-32C. The performance of the low-latency queue is measured in conjunction with multi-level priority queues. From the results, you can see that LLQ is given preferential treatment for latency constraints. All queues are maintained below 10μs average latency, with LLQ operating at <6μs.

Table 24: ACX7100-32C Low-Latency with Continuous Background Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|---------------------|---------------------|---------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 4.57 | 4.78 | 9.30 | 64 | continuous | 10G |
| FC-SIGNALING | Strict-High | 4.57 | 4.83 | 11.84 | 64 | continuous | 10G |
| FC-CONTROL | High | 4.58 | 4.86 | 14.89 | 64 | continuous | 10G |
| FC-REALTIME | Medium-High | 4.57 | 5.20 | 21.32 | 64 | continuous | 10G |
| FC-HIGH | Low | 4.58 | 4.94 | 27.61 | 64 | continuous | 10G |
| FC-MEDIUM | Low | 4.61 | 4.96 | 31.12 | 64 | continuous | 10G |
| FC-LOW | Low | 4.58 | 4.95 | 25.48 | 64 | continuous | 10G |

Table 24: ACX7100-32C Low-Latency with Continuous Background Traffic (Continued)

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 5.29 | 5.53 | 8.85 | 512 | continuous | 10G |
| FC-SIGNALING | Strict-High | 5.29 | 5.70 | 12.28 | 512 | continuous | 10G |
| FC-CONTROL | High | 5.34 | 5.83 | 18.32 | 512 | continuous | 10G |
| FC-REALTIME | Medium-High | 5.29 | 5.95 | 17.89 | 512 | continuous | 10G |
| FC-HIGH | Low | 5.29 | 5.94 | 28.39 | 512 | continuous | 10G |
| FC-MEDIUM | Low | 5.34 | 6.47 | 21.83 | 512 | continuous | 10G |
| FC-LOW | Low | 5.35 | 5.68 | 28.15 | 512 | continuous | 10G |
| FC-LLQ | Low-Latency | 4.91 | 5.29 | 10.27 | 1500 | continuous | 10G |
| FC-SIGNALING | Strict-High | 4.91 | 5.67 | 15.49 | 1500 | continuous | 10G |
| FC-CONTROL | High | 4.96 | 6.14 | 18.20 | 1500 | continuous | 10G |
| FC-REALTIME | Medium-High | 4.91 | 6.03 | 24.50 | 1500 | continuous | 10G |
| FC-HIGH | Low | 4.96 | 6.45 | 39.09 | 1500 | continuous | 10G |
| FC-MEDIUM | Low | 4.96 | 7.52 | 50.54 | 1500 | continuous | 10G |
| FC-LOW | Low | 4.96 | 6.76 | 55.82 | 1500 | continuous | 10G |

The following table includes burst traffic type rather than continuous for Topology 1b with ACX7100-32C, measuring the performance of the low-latency queue in conjunction with multi-level

priority queues. From the results, you can see that LLQ is given preferential treatment for latency constraints. In these scenarios, all queues are transmitted $< 6\mu\text{s}$, with LLQ having slightly better performance.

Table 25: ACX7100-32C Low-Latency with Burst Background Traffic

| Queue Name | Queue Priority | Min (μs) | Ave (μs) | Max (μs) | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|-----------------------|-----------------------|-----------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 4.61 | 4.65 | 4.79 | 64 | burst | 10G |
| FC-SIGNALING | Strict-High | 4.61 | 4.67 | 5.04 | 64 | burst | 10G |
| FC-CONTROL | High | 4.61 | 4.67 | 5.04 | 64 | burst | 10G |
| FC-REALTIME | Medium-High | 4.61 | 4.66 | 5.04 | 64 | burst | 10G |
| FC-HIGH | Low | 4.61 | 4.67 | 5.04 | 64 | burst | 10G |
| FC-MEDIUM | Low | 4.61 | 4.67 | 5.04 | 64 | burst | 10G |
| FC-LOW | Low | 4.61 | 4.67 | 5.04 | 64 | burst | 10G |
| FC-LLQ | Low-Latency | 5.31 | 5.38 | 5.49 | 512 | burst | 10G |
| FC-SIGNALING | Strict-High | 5.31 | 5.42 | 8.41 | 512 | burst | 10G |
| FC-CONTROL | High | 5.31 | 5.42 | 5.78 | 512 | burst | 10G |
| FC-REALTIME | Medium-High | 5.31 | 5.42 | 7.82 | 512 | burst | 10G |
| FC-HIGH | Low | 5.31 | 5.42 | 5.74 | 512 | burst | 10G |
| FC-MEDIUM | Low | 5.31 | 5.42 | 7.73 | 512 | burst | 10G |

Table 25: ACX7100-32C Low-Latency with Burst Background Traffic (Continued)

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LOW | Low | 5.31 | 5.42 | 5.74 | 512 | burst | 10G |
| FC-LLQ | Low-Latency | 4.94 | 5.01 | 5.87 | 1500 | burst | 10G |
| FC-SIGNALING | Strict-High | 4.94 | 5.13 | 7.81 | 1500 | burst | 10G |
| FC-CONTROL | High | 4.94 | 5.13 | 8.42 | 1500 | burst | 10G |
| FC-REALTIME | Medium-High | 4.94 | 5.13 | 8.55 | 1500 | burst | 10G |
| FC-HIGH | Low | 4.94 | 5.13 | 5.69 | 1500 | burst | 10G |
| FC-MEDIUM | Low | 4.94 | 5.13 | 5.65 | 1500 | burst | 10G |
| FC-LOW | Low | 4.94 | 5.13 | 5.76 | 1500 | burst | 10G |

In the next table, look at the impact on the LLQ when creating congestion to the point of traffic discarding in the oversubscribed strict-high queue. The intentionality of the LLQ design prevents such scenarios from disrupting the LLQ. From the results, you can observe that the latency budget is preserved at ~4-5μs and well below the goal of 10μs.

Table 26: ACX7100-32C Low-Latency with Congestion Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Congested Queue |
|------------|----------------|------------------|------------------|------------------|------------|-----------------|-----------------|
| FC-LLQ | Low-Latency | 4.58 | 4.89 | 11.04 | 64 | continuous | Strict-High |
| FC-LLQ | Low-Latency | 5.34 | 5.57 | 11.54 | 512 | continuous | Strict-High |

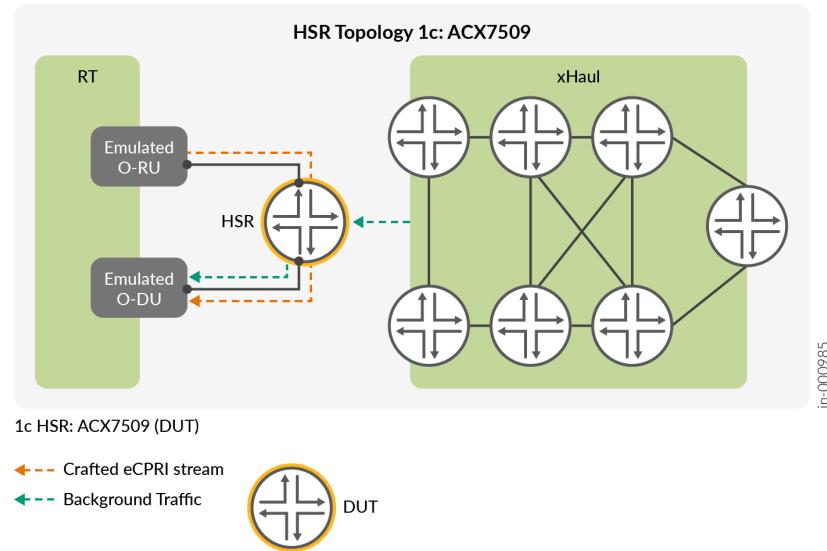
Table 26: ACX7100-32C Low-Latency with Congestion Traffic *(Continued)*

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Congested Queue |
|------------|-----------------|---------------------|---------------------|---------------------|------------|-----------------|-----------------|
| FC-LLQ | Low- Latency | 4.93 | 5.13 | 11.46 | 1500 | continuous | Strict-High |
| FC-LLQ | Low- Latency | 4.59 | 4.70 | 11.02 | 64 | burst | Strict-High |
| FC-LLQ | Low- Latency | 5.30 | 5.50 | 10.14 | 512 | burst | Strict-High |
| FC-LLQ | Low- Latency | 4.92 | 5.03 | 11.09 | 1500 | burst | Strict-High |

Topology 1c: ACX7509

In Topology 1c, same test scenarios are executed while moving to the second HSR device with the examination of ACX7509 performance. Once again, the test equipment emulates O-RU and O-DU without self-latency. There is no physical CSR in this case as the performance of the single device is captured. For each scenario, the results are provided across different frame sizes with either continuous or burst traffic patterns. Latency is measured in microseconds (μs).

Figure 30: Topology 1c ACX7509 Standalone DUT



The following table describes the Topology 1c ACX7509 performance when traffic is sent only to the LLQ without background traffic. No queues are congested during this time.

Table 27: ACX7509 Low-Latency without Background Traffic

| Queue Name | Queue Priority | Min (μs) | Ave (μs) | Max (μs) | Frame Size | Traffic Pattern | Port Speed |
|------------|----------------|----------|----------|----------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 4.72 | 5.16 | 8.49 | 64 | continuous | 10G |
| FC-LLQ | Low-Latency | 5.30 | 6.12 | 9.22 | 512 | continuous | 10G |
| FC-LLQ | Low-Latency | 5.10 | 6.03 | 9.50 | 1500 | continuous | 10G |

The following table includes continuous background traffic up to 99%-line rate with encapsulation overhead considerations for Topology 1c with ACX7509. Performance of the FC-LLQ is measured in conjunction with multi-level priority queues. From the results, you can see that LLQ is given preferential treatment for latency constraints. All priority queues are maintained below 10μs average latency, with LLQ operating ≤6μs.

Table 28: ACX7509 Low-Latency with Continuous Background Traffic

| Queue Name | Queue Priority | Min (μs) | Ave (μs) | Max (μs) | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|----------|----------|----------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 4.93 | 5.24 | 9.85 | 64 | continuous | 10G |
| FC-SIGNALING | Strict-High | 4.95 | 5.31 | 12.62 | 64 | continuous | 10G |
| FC-CONTROL | High | 4.97 | 5.34 | 15.32 | 64 | continuous | 10G |
| FC-REALTIME | Medium-High | 4.92 | 6.04 | 21.86 | 64 | continuous | 10G |
| FC-HIGH | Low | 4.97 | 5.41 | 36.48 | 64 | continuous | 10G |
| FC-MEDIUM | Low | 4.97 | 5.47 | 32.47 | 64 | continuous | 10G |
| FC-LOW | Low | 4.94 | 5.43 | 36.44 | 64 | continuous | 10G |
| FC-LLQ | Low-Latency | 5.69 | 6.18 | 10.24 | 512 | continuous | 10G |
| FC-SIGNALING | Strict-High | 5.69 | 6.54 | 14.42 | 512 | continuous | 10G |
| FC-CONTROL | High | 5.72 | 6.65 | 19.17 | 512 | continuous | 10G |
| FC-REALTIME | Medium-High | 5.69 | 7.69 | 22.86 | 512 | continuous | 10G |
| FC-HIGH | Low | 5.69 | 7.32 | 37.08 | 512 | continuous | 10G |
| FC-MEDIUM | Low | 5.73 | 8.01 | 35.80 | 512 | continuous | 10G |
| FC-LOW | Low | 5.73 | 7.28 | 36.19 | 512 | continuous | 10G |

Table 28: ACX7509 Low-Latency with Continuous Background Traffic *(Continued)*

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 5.34 | 5.91 | 12.29 | 1500 | continuous | 10G |
| FC-SIGNALING | Strict-High | 5.32 | 6.75 | 16.29 | 1500 | continuous | 10G |
| FC-CONTROL | High | 5.34 | 7.28 | 18.77 | 1500 | continuous | 10G |
| FC-REALTIME | Medium-High | 5.34 | 8.00 | 26.09 | 1500 | continuous | 10G |
| FC-HIGH | Low | 5.34 | 8.85 | 45.10 | 1500 | continuous | 10G |
| FC-MEDIUM | Low | 5.34 | 10.19 | 57.32 | 1500 | continuous | 10G |
| FC-LOW | Low | 5.34 | 9.30 | 60.16 | 1500 | continuous | 10G |

The following table includes burst traffic type rather than continuous for Topology 1c with ACX7509, measuring the performance of the low-latency queue in conjunction with multi-level priority queues. From the results, you can see that LLQ is given preferential treatment for latency constraints. In these scenarios, all queues are transmitted within $\leq 6\mu\text{s}$, with LLQ having slightly better performance.

Table 29: ACX7509 Low-Latency with Burst Background Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-LLQ | Low-Latency | 4.92 | 5.01 | 5.12 | 64 | burst | 10G |
| FC-SIGNALING | Strict-High | 4.92 | 5.07 | 8.70 | 64 | burst | 10G |
| FC-CONTROL | High | 4.92 | 5.07 | 9.44 | 64 | burst | 10G |

Table 29: ACX7509 Low-Latency with Burst Background Traffic (Continued)

| Queue Name | Queue Priority | Min (μs) | Ave (μs) | Max (μs) | Frame Size | Traffic Pattern | Port Speed |
|--------------|----------------|----------|----------|----------|------------|-----------------|------------|
| FC-REALTIME | Medium-High | 4.92 | 5.27 | 8.71 | 64 | burst | 10G |
| FC-HIGH | Low | 4.92 | 5.10 | 8.70 | 64 | burst | 10G |
| FC-MEDIUM | Low | 4.92 | 5.10 | 8.70 | 64 | burst | 10G |
| FC-LOW | Low | 4.92 | 5.10 | 8.70 | 64 | burst | 10G |
| FC-LLQ | Low-Latency | 5.71 | 5.75 | 5.85 | 512 | burst | 10G |
| FC-SIGNALING | Strict-High | 5.71 | 6.00 | 9.73 | 512 | burst | 10G |
| FC-CONTROL | High | 5.71 | 6.00 | 11.56 | 512 | burst | 10G |
| FC-REALTIME | Medium-High | 5.71 | 6.19 | 10.54 | 512 | burst | 10G |
| FC-HIGH | Low | 5.71 | 6.00 | 9.43 | 512 | burst | 10G |
| FC-MEDIUM | Low | 5.71 | 6.00 | 9.43 | 512 | burst | 10G |
| FC-LOW | Low | 5.71 | 6.00 | 9.43 | 512 | burst | 10G |
| FC-LLQ | Low-Latency | 5.30 | 5.38 | 7.37 | 1500 | burst | 10G |
| FC-SIGNALING | Strict-High | 5.30 | 6.10 | 9.07 | 1500 | burst | 10G |
| FC-CONTROL | High | 5.30 | 6.10 | 9.21 | 1500 | burst | 10G |

Table 29: ACX7509 Low-Latency with Burst Background Traffic (Continued)

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Port Speed |
|-------------|----------------|------------------|------------------|------------------|------------|-----------------|------------|
| FC-REALTIME | Medium-High | 5.30 | 6.13 | 12.20 | 1500 | burst | 10G |
| FC-HIGH | Low | 5.30 | 6.10 | 9.06 | 1500 | burst | 10G |
| FC-MEDIUM | Low | 5.30 | 6.10 | 11.82 | 1500 | burst | 10G |
| FC-LOW | Low | 5.30 | 6.10 | 9.02 | 1500 | burst | 10G |

The next table provides the impact on the low-latency queue when creating congestion to the point of traffic discarding in the oversubscribed strict-high queue. The intentionality of the design prevents such scenarios from disrupting the LLQ. From the results, you can observe that the latency budget is preserved at <6μs and well below the goal of 10μs.

Table 30: ACX7509 Low-Latency with Congestion Traffic

| Queue Name | Queue Priority | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Traffic Pattern | Congested Queue |
|------------|----------------|------------------|------------------|------------------|------------|-----------------|-----------------|
| FC-LLQ | Low-Latency | 4.93 | 5.26 | 11.40 | 64 | continuous | Strict-High |
| FC-LLQ | Low-Latency | 5.71 | 5.94 | 11.89 | 512 | continuous | Strict-High |
| FC-LLQ | Low-Latency | 5.30 | 5.52 | 11.50 | 1500 | continuous | Strict-High |
| FC-LLQ | Low-Latency | 4.94 | 5.05 | 11.36 | 64 | burst | Strict-High |
| FC-LLQ | Low-Latency | 5.67 | 5.91 | 10.52 | 512 | burst | Strict-High |
| FC-LLQ | Low-Latency | 5.32 | 5.40 | 11.48 | 1500 | burst | Strict-High |

eCPRI Validation

In this section, the validation is summarized and performed with real eCPRI traffic patterns, O-RAN conformance testing, and handling of eCPRI message types. For all eCPRI and O-RAN conformance validation results with flow data and measurements, see the Test Report. This section captures the most critical data to explain the network and device performance. Latency is measured in microseconds (μ s).

The O-RAN and eCPRI test scenarios include several functional permutations with emulated O-DU and RRU. Innovative steps are taken to validate the performance and assurance of 5G eCPRI communications, U-Plane message-type behaviors, and O-RAN conformance. These steps ensure that the featured DUTs are correctly and consistently transporting critical 5G services.

Results evolve across multiple topologies to capture a range of DUT performance characteristics in non-congested and congested scenarios.

Here is the summary of O-RAN and eCPRI test scenarios:

- eCPRI O-RAN Emulation: Leverages a standard IQ Sample file to generate flows.
- O-RAN Conformance: Analyzes messages for conformance to O-RAN specification.
- Crafted eCPRI O-RAN: Produces variable eCPRI payload for comparing latency performance.
- eCPRI Services Validation: Emulates User Plane messages, performing functional and integrity analysis.
- eCPRI Remote Memory Access (Type 4): Performs read or write from/to a specific memory address on the opposite eCPRI node and validates expected success or failure conditions.
- eCPRI Delay Measurement Message (Type 5): Estimates one-way delay between two eCPRI ports.
- eCPRI Remote Reset Message (Type 6): Requests reset of one eCPRI node from another node. This validates expected sender and receiver operations.
- eCPRI Event Indication Message (Type 7): An event notification message to convey information about raised or ceased faults or general notifications. Validation confirms events raised as expected.

Topology 1: Latency Performance with eCPRI

The proceeding test scenarios provide results across topology 1 for single-DUT performance, topology 2 for fronthaul EVPN-VPWS, and topology 3 for fronthaul active-active EVPN Multihoming. The focus of this section is on measuring the LLQ performance in non-congested and heavy congestion scenarios using real eCPRI traffic streams. Background traffic is sent up to 99%-line rate or exceeds the configured

bandwidth rate to the point of discarding frames. The validation aims to measure the LLQ performance in these conditions.

For brevity, the data is summarized here with crucial results. For additional test details, including latency measured for all queues, see the Test Report or contact your Juniper representative for more information.

In the following table, real eCPRI traffic is forwarded across topology 1a with ACX7024, topology 1b with ACX7100-32C, and topology 1c with ACX7509 for 10GbE and 100GbE port speeds with continuous or burst traffic patterns. In all cases, the average latency is measured as less than the 10 μ s objective, with the majority $\leq 6\mu$ s across all three DUTs (ACX7024, ACX7100-32C, and ACX7509).

Table 31: eCPRI ORAN Latency Performance without Congestion

| DUT | Scenario | Min (μ s) Latency | Ave (μ s) Latency | Max (μ s) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|-------------------|---------------------------|---------------------------|---------------------------|---------------|---------------|--------------------|-----------------|
| ACX7024 | Non- Congested | 5.71 | 5.86 | 12.72 | 64 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | Non- Congested | 5.56 | 5.76 | 12.33 | 512 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | Non- Congested | 4.88 | 5.20 | 12.45 | 1500 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | Non- Congested | 5.70 | 5.82 | 6.36 | 64 | 10GbE | Burst | eCPRI- ORAN |
| ACX7024 | Non- Congested | 5.45 | 5.68 | 6.40 | 512 | 10GbE | Burst | eCPRI- ORAN |
| ACX7024 | Non- Congested | 4.95 | 4.99 | 6.45 | 1500 | 10GbE | Burst | eCPRI- ORAN |
| ACX7024 | Non- Congested | 5.71 | 6.01 | 6.55 | 64 | 100Gb E | Continuo us | eCPRI- ORAN |
| ACX7024 | Non- Congested | 5.74 | 5.85 | 6.36 | 512 | 100Gb E | Continuo us | eCPRI- ORAN |
| ACX7024 | Non- Congested | 5.78 | 5.80 | 6.43 | 1500 | 100Gb E | Continuo us | eCPRI- ORAN |

Table 31: eCPRI ORAN Latency Performance without Congestion (Continued)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|---------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX7024 | Non-Congested | 5.74 | 6.00 | 6.54 | 64 | 100GbE | Burst | eCPRI-ORAN |
| ACX7024 | Non-Congested | 5.73 | 5.84 | 6.35 | 512 | 100GbE | Burst | eCPRI-ORAN |
| ACX7024 | Non-Congested | 5.77 | 5.79 | 6.40 | 1500 | 100GbE | Burst | eCPRI-ORAN |
| ACX7100 | Non-Congested | 4.94 | 5.81 | 12.40 | 64 | 10GbE | Continuous | eCPRI-ORAN |
| ACX7100 | Non-Congested | 5.68 | 6.36 | 12.40 | 512 | 10GbE | Continuous | eCPRI-ORAN |
| ACX7100 | Non-Congested | 5.33 | 6.26 | 11.81 | 1500 | 10GbE | Continuous | eCPRI-ORAN |
| ACX7100 | Non-Congested | 4.91 | 4.98 | 5.15 | 64 | 10GbE | Burst | eCPRI-ORAN |
| ACX7100 | Non-Congested | 5.69 | 5.77 | 5.92 | 512 | 10GbE | Burst | eCPRI-ORAN |
| ACX7100 | Non-Congested | 5.33 | 5.38 | 6.29 | 1500 | 10GbE | Burst | eCPRI-ORAN |
| ACX7100 | Non-Congested | 4.44 | 4.47 | 4.71 | 64 | 100GbE | Continuous | eCPRI-ORAN |
| ACX7100 | Non-Congested | 4.47 | 4.53 | 4.69 | 512 | 100GbE | Continuous | eCPRI-ORAN |
| ACX7100 | Non-Congested | 4.44 | 4.60 | 4.71 | 1500 | 100GbE | Continuous | eCPRI-ORAN |
| ACX7100 | Non-Congested | 4.43 | 4.60 | 4.71 | 64 | 100GbE | Burst | eCPRI-ORAN |

Table 31: eCPRI ORAN Latency Performance without Congestion (Continued)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|-------------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX7100 | Non- Congested | 4.47 | 4.53 | 4.73 | 512 | 100Gb E | Burst | eCPRI- ORAN |
| ACX7100 | Non- Congested | 4.44 | 4.46 | 4.70 | 1500 | 100Gb E | Burst | eCPRI- ORAN |
| ACX7509 | Non- Congested | 4.59 | 4.66 | 8.55 | 64 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | Non- Congested | 5.30 | 5.43 | 8.92 | 512 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | Non- Congested | 4.92 | 5.12 | 10.40 | 1500 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | Non- Congested | 4.59 | 4.63 | 4.82 | 64 | 10GbE | Burst | eCPRI- ORAN |
| ACX7509 | Non- Congested | 4.93 | 5.16 | 5.88 | 512 | 10GbE | Burst | eCPRI- ORAN |
| ACX7509 | Non- Congested | 4.93 | 5.01 | 5.88 | 1500 | 10GbE | Burst | eCPRI- ORAN |
| ACX7509 | Non- Congested | 5.71 | 6.01 | 6.55 | 64 | 100Gb E | Continuo us | eCPRI- ORAN |
| ACX7509 | Non- Congested | 5.74 | 5.85 | 6.36 | 512 | 100Gb E | Continuo us | eCPRI- ORAN |
| ACX7509 | Non- Congested | 5.78 | 5.80 | 6.43 | 1500 | 100Gb E | Continuo us | eCPRI- ORAN |
| ACX7509 | Non- Congested | 4.12 | 4.37 | 4.47 | 64 | 100Gb E | Burst | eCPRI- ORAN |
| ACX7509 | Non- Congested | 4.19 | 4.24 | 4.55 | 512 | 100Gb E | Burst | eCPRI- ORAN |

Table 31: eCPRI ORAN Latency Performance without Congestion (Continued)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|-------------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX7509 | Non- Congested | 4.15 | 4.16 | 4.58 | 1500 | 100Gb E | Burst | eCPRI- ORAN |

The following table validates the same scenario as above, introducing heavy congestion with the strict-high queue oversubscribed and dropping traffic. The intentionality of the LLQ design prevents such scenarios from causing significant disruption to the low-latency queue while sending eCPRI traffic streams. From the results, you can observe that the average latency budget is preserved below 10μs. As the port speed increases, the transmission delay associated with congestion typically decreases. As a result, slightly lower latency can be seen on 100GbE port speeds.

Table 32: eCPRI ORAN Latency Performance with 1 Queue Congested

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|------------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX7024 | 1Q- Congested | 5.88 | 7.24 | 15.96 | 64 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 6.18 | 7.50 | 16.27 | 512 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 6.28 | 7.78 | 19.60 | 1500 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 5.69 | 6.07 | 12.23 | 64 | 10GbE | Burst | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 5.64 | 5.87 | 10.56 | 512 | 10GbE | Burst | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 4.95 | 5.12 | 12.88 | 1500 | 10GbE | Burst | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 5.70 | 6.02 | 7.02 | 64 | 100GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 5.72 | 5.87 | 6.84 | 512 | 100GbE | Continuo us | eCPRI- ORAN |

Table 32: eCPRI ORAN Latency Performance with 1 Queue Congested (*Continued*)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|------------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX7024 | 1Q- Congested | 5.78 | 5.81 | 6.89 | 1500 | 100GbE | Continuo us | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 5.77 | 6.02 | 6.95 | 64 | 100GbE | Burst | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 5.71 | 5.86 | 6.86 | 512 | 100GbE | Burst | eCPRI- ORAN |
| ACX7024 | 1Q- Congested | 5.77 | 5.80 | 6.83 | 1500 | 100GbE | Burst | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 5.03 | 7.16 | 16.29 | 64 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 6.14 | 7.69 | 15.61 | 512 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 6.49 | 8.61 | 21.39 | 1500 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 4.91 | 5.28 | 11.45 | 64 | 10GbE | Burst | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 5.70 | 5.96 | 12.00 | 512 | 10GbE | Burst | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 5.30 | 5.42 | 12.01 | 1500 | 10GbE | Burst | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 4.44 | 4.49 | 5.08 | 64 | 100GbE | Continuo us | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 4.47 | 4.55 | 5.07 | 512 | 100GbE | Continuo us | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 4.42 | 4.61 | 5.48 | 1500 | 100GbE | Continuo us | eCPRI- ORAN |

Table 32: eCPRI ORAN Latency Performance with 1 Queue Congested (*Continued*)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|------------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX7100 | 1Q- Congested | 4.41 | 4.61 | 5.41 | 64 | 100GbE | Burst | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 4.46 | 4.54 | 5.07 | 512 | 100GbE | Burst | eCPRI- ORAN |
| ACX7100 | 1Q- Congested | 4.44 | 4.48 | 5.23 | 1500 | 100GbE | Burst | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 4.64 | 6.63 | 14.89 | 64 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 5.71 | 7.14 | 14.85 | 512 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 6.09 | 8.04 | 20.25 | 1500 | 10GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 4.59 | 4.92 | 11.10 | 64 | 10GbE | Burst | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 5.30 | 5.57 | 11.17 | 512 | 10GbE | Burst | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 4.59 | 4.68 | 12.39 | 1500 | 10GbE | Burst | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 4.13 | 4.17 | 4.94 | 64 | 100GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 4.18 | 4.24 | 4.76 | 512 | 100GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 4.10 | 4.36 | 4.96 | 1500 | 100GbE | Continuo us | eCPRI- ORAN |
| ACX7509 | 1Q- Congested | 4.12 | 4.37 | 5.02 | 64 | 100GbE | Burst | eCPRI- ORAN |

Table 32: eCPRI ORAN Latency Performance with 1 Queue Congested (*Continued*)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|---------|--------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX7509 | 1Q-Congested | 4.19 | 4.25 | 4.87 | 512 | 100GbE | Burst | eCPRI-ORAN |
| ACX7509 | 1Q-Congested | 4.15 | 4.19 | 5.07 | 1500 | 100GbE | Burst | eCPRI-ORAN |

The following table continues to increase very heavy congestion with multiple queues simultaneously dropping traffic. In this scenario, both strict-high and high priority queues are oversubscribed. The LLQ implementation prevents disruption to delay sensitive eCPRI traffic, maintaining latency budget. Despite the increased congestion with all queues sending traffic and two queues dropping, latency is preserved below the 10μs objective in all cases. As the port speed increases, the transmission delay associated with congestion typically decreases. As a result, you can see even more benefit as the congestion increases, with measurements consistently in the range of 4-5μs.

Table 33: eCPRI ORAN Latency Performance with Heavy Congestion (2Q)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|----------|--------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX702 4 | 2Q-Congested | 5.90 | 7.71 | 22.88 | 64 | 10GbE | Continuous | eCPRI-ORAN |
| ACX702 4 | 2Q-Congested | 6.18 | 7.37 | 16.89 | 512 | 10GbE | Continuous | eCPRI-ORAN |
| ACX702 4 | 2Q-Congested | 6.29 | 7.39 | 23.58 | 1500 | 10GbE | Continuous | eCPRI-ORAN |
| ACX702 4 | 2Q-Congested | 5.70 | 6.52 | 16.33 | 64 | 10GbE | Burst | eCPRI-ORAN |
| ACX702 4 | 2Q-Congested | 5.64 | 6.17 | 14.18 | 512 | 10GbE | Burst | eCPRI-ORAN |
| ACX702 4 | 2Q-Congested | 4.95 | 5.35 | 15.74 | 1500 | 10GbE | Burst | eCPRI-ORAN |

Table 33: eCPRI ORAN Latency Performance with Heavy Congestion (2Q) (Continued)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|-------------|------------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX702 4 | 2Q- Congested | 5.77 | 6.07 | 7.55 | 64 | 100Gb E | Continuous | eCPRI- ORAN |
| ACX702 4 | 2Q- Congested | 5.71 | 5.93 | 7.23 | 512 | 100Gb E | Continuous | eCPRI- ORAN |
| ACX702 4 | 2Q- Congested | 5.78 | 5.85 | 7.17 | 1500 | 100Gb E | Continuous | eCPRI- ORAN |
| ACX702 4 | 2Q- Congested | 5.77 | 6.06 | 7.54 | 64 | 100Gb E | Burst | eCPRI- ORAN |
| ACX702 4 | 2Q- Congested | 5.70 | 5.92 | 7.15 | 512 | 100Gb E | Burst | eCPRI- ORAN |
| ACX702 4 | 2Q- Congested | 5.77 | 5.84 | 7.16 | 1500 | 100Gb E | Burst | eCPRI- ORAN |
| ACX710 0 | 2Q- Congested | 4.98 | 7.09 | 22.41 | 64 | 10GbE | Continuous | eCPRI- ORAN |
| ACX710 0 | 2Q- Congested | 6.10 | 7.40 | 16.54 | 512 | 10GbE | Continuous | eCPRI- ORAN |
| ACX710 0 | 2Q- Congested | 6.49 | 8.55 | 26.38 | 1500 | 10GbE | Continuous | eCPRI- ORAN |
| ACX710 0 | 2Q- Congested | 4.91 | 5.81 | 15.77 | 64 | 10GbE | Burst | eCPRI- ORAN |
| ACX710 0 | 2Q- Congested | 5.70 | 6.31 | 13.40 | 512 | 10GbE | Burst | eCPRI- ORAN |
| ACX710 0 | 2Q- Congested | 5.33 | 5.97 | 16.20 | 1500 | 10GbE | Burst | eCPRI- ORAN |
| ACX710 0 | 2Q- Congested | 4.41 | 4.65 | 5.84 | 64 | 100Gb E | Continuous | eCPRI- ORAN |

Table 33: eCPRI ORAN Latency Performance with Heavy Congestion (2Q) (Continued)

| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|----------|--------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX710 0 | 2Q-Congested | 4.47 | 4.58 | 5.44 | 512 | 100GbE | Continuous | eCPRI-ORAN |
| ACX710 0 | 2Q-Congested | 4.44 | 4.52 | 5.68 | 1500 | 100GbE | Continuous | eCPRI-ORAN |
| ACX710 0 | 2Q-Congested | 4.12 | 4.37 | 5.02 | 64 | 100GbE | Burst | eCPRI-ORAN |
| ACX710 0 | 2Q-Congested | 4.19 | 4.25 | 4.87 | 512 | 100GbE | Burst | eCPRI-ORAN |
| ACX710 0 | 2Q-Congested | 4.15 | 4.19 | 5.07 | 1500 | 100GbE | Burst | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.60 | 6.54 | 19.92 | 64 | 10GbE | Continuous | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 5.69 | 6.93 | 15.62 | 512 | 10GbE | Continuous | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 6.10 | 7.70 | 21.52 | 1500 | 10GbE | Continuous | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.59 | 5.46 | 15.42 | 64 | 10GbE | Burst | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 5.30 | 5.88 | 15.93 | 512 | 10GbE | Burst | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.93 | 5.60 | 15.82 | 1500 | 10GbE | Burst | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.12 | 4.39 | 5.52 | 64 | 100GbE | Continuous | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.18 | 4.26 | 5.16 | 512 | 100GbE | Continuous | eCPRI-ORAN |

Table 33: eCPRI ORAN Latency Performance with Heavy Congestion (2Q) (Continued)

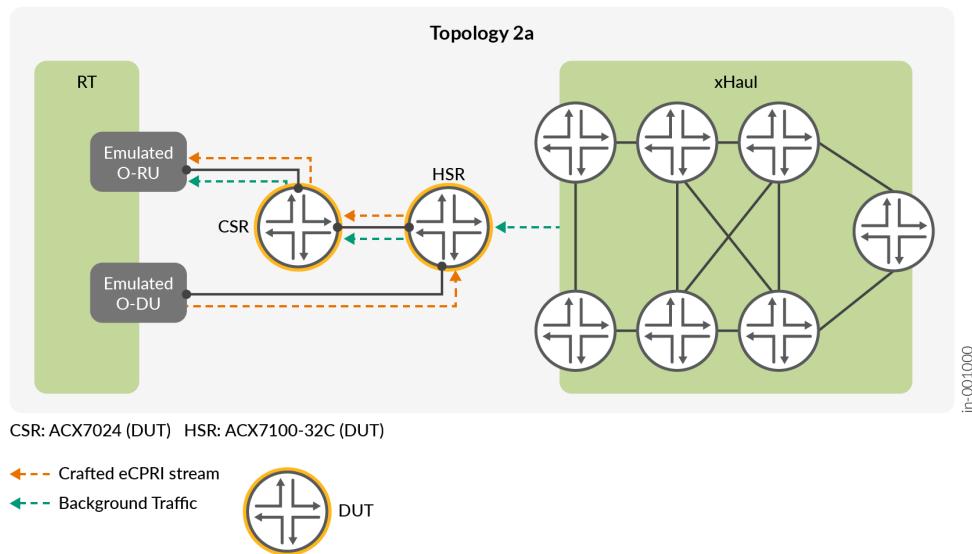
| DUT | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Traffic Type |
|----------|--------------|---------------------|---------------------|---------------------|---------------|---------------|--------------------|-----------------|
| ACX750 9 | 2Q-Congested | 4.13 | 4.21 | 5.38 | 1500 | 100GbE | Continuous | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.13 | 4.41 | 5.57 | 64 | 100GbE | Burst | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.19 | 4.27 | 5.18 | 512 | 100GbE | Burst | eCPRI-ORAN |
| ACX750 9 | 2Q-Congested | 4.15 | 4.22 | 5.37 | 1500 | 100GbE | Burst | eCPRI-ORAN |

Topology 2: Latency Performance with eCPRI over EVPN-VPWS

Topology 2 measures latency in the LLQ over point-to-point EVPN-VPWS services in the fronthaul network segment between CSR and HSR. In topology 2a, the ACX7024 is the CSR DUT, and ACX7100-32C is the HSR DUT. The test equipment emulates the O-RU and O-DU without self-latency. The focus of this section is on measuring the LLQ performance during conditions of non-congested and heavy congestion scenarios using real eCPRI traffic streams. Background traffic is sent up to 99%-line rate or exceeding the configured bandwidth rate to the point of discarding frames.

For brevity, the information provided here is summarized with results for EVPN-VPWS 5G fronthaul eCPRI performance. For fronthaul, midhaul and MBH workloads, see the Test Report of all data across multiple VPN services representing the traffic profiles explained in this document. In addition, the Test Report includes results for topology 2b, which features ACX7024 as the CSR and ACX7509 as the HSR.

Figure 31: Topology 2a ACX7024 CSR and ACX7100-32C HSR



Topology 2 provides single-homed performance data in a 2-hop scenario without a physical RU or DU element (both emulated by the test center). The following table measures the latency for eCPRI traffic sent over Topology 2, which includes two hops supporting point-to-point EVPN-VPWS services. The data shown is for Topology 2a, including ACX7024 as CSR and ACX7100-32C as HSR. Additional test results are found in the Test Report. Results are provided across different frame sizes, 10GbE and 100GbE port speeds, and continuous or burst traffic patterns. Latency is always measured in microseconds (μ s). In all cases, the per-device average latency is measured less than the 10 μ s objective with the majority $\leq 6\mu$ s.

Table 34: EVPN-VPWS Low-Latency without Congestion

| Traffic Type | Per-Hop Ave (μ s) | Scenario | Min (μ s) Latency | Ave (μ s) Latency | Max (μ s) Latency | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------------|---------------|------------------------|------------------------|------------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 5.29 | Non-Congested | 10.45 | 10.59 | 10.89 | 64 | 10GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.52 | Non-Congested | 10.89 | 11.04 | 11.40 | 512 | 10GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.30 | Non-Congested | 10.48 | 10.61 | 11.53 | 1500 | 10GbE | Continuous | 2 |

Table 34: EVPN-VPWS Low-Latency without Congestion (*Continued*)

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|---------------|------------------|------------------|------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 5.29 | Non-Congested | 10.46 | 10.59 | 10.84 | 64 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.52 | Non-Congested | 10.90 | 11.05 | 11.32 | 512 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.30 | Non-Congested | 10.46 | 10.62 | 11.47 | 1500 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.25 | Non-Congested | 10.00 | 10.50 | 10.94 | 64 | 100GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.10 | Non-Congested | 10.10 | 10.21 | 10.48 | 512 | 100GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.11 | Non-Congested | 10.12 | 10.23 | 10.60 | 1500 | 100GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.39 | Non-Congested | 10.40 | 10.80 | 11.17 | 64 | 100GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.29 | Non-Congested | 10.47 | 10.59 | 10.93 | 512 | 100GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.29 | Non-Congested | 10.43 | 10.58 | 11.01 | 1500 | 100GbE | Burst | 2 |

As shown, the average per-hop latency (2nd column) is kept below 6μs when the network is not experiencing congestion. For the test validation, eCPRI traffic is generated from CSR (ACX7024) to HSR (ACX7100-32C).

The following table describes the latency performance with congestion created on the strict-high queue and additional scenarios including heavy congestion with traffic loss across multiple queues (strict-high and high). The congestion is created at the CSR ACX7024, resulting in discards on the egress path toward ACX7100-32C HSR.

Table 35: EVPN-VPWS Low-Latency with Heavy Congestion

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|--------------|---------------------|---------------------|---------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 6.05 | 1Q-Congested | 10.71 | 12.11 | 20.12 | 64 | 10GbE | Continous | 2 |
| EVPN-VPWS (eCPRI) | 6.44 | 1Q-Congested | 11.50 | 12.89 | 21.10 | 512 | 10GbE | Continous | 2 |
| EVPN-VPWS (eCPRI) | 6.76 | 1Q-Congested | 11.87 | 13.53 | 22.96 | 1500 | 10GbE | Continous | 2 |
| EVPN-VPWS (eCPRI) | 5.37 | 1Q-Congested | 10.46 | 10.74 | 16.30 | 64 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.61 | 1Q-Congested | 10.90 | 11.23 | 16.20 | 512 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.38 | 1Q-Congested | 10.50 | 10.78 | 16.51 | 1500 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.29 | 1Q-Congested | 10.01 | 10.60 | 11.53 | 64 | 100GbE | Continous | 2 |
| EVPN-VPWS (eCPRI) | 5.12 | 1Q-Congested | 10.08 | 10.25 | 11.27 | 512 | 100GbE | Continous | 2 |
| EVPN-VPWS (eCPRI) | 5.12 | 1Q-Congested | 10.10 | 10.26 | 11.16 | 1500 | 100GbE | Continous | 2 |
| EVPN-VPWS (eCPRI) | 5.42 | 1Q-Congested | 10.41 | 10.84 | 11.94 | 64 | 100GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.32 | 1Q-Congested | 10.44 | 10.64 | 11.70 | 512 | 100GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.31 | 1Q-Congested | 10.43 | 10.62 | 11.58 | 1500 | 100GbE | Burst | 2 |

Table 35: EVPN-VPWS Low-Latency with Heavy Congestion (Continued)

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min Latency (μs) | Ave Latency (μs) | Max Latency (μs) | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|--------------|------------------|------------------|------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 5.78 | 2Q-Congested | 10.65 | 11.56 | 23.78 | 64 | 10GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 6.32 | 2Q-Congested | 11.46 | 12.64 | 22.78 | 512 | 10GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 6.53 | 2Q-Congested | 11.82 | 13.07 | 28.35 | 1500 | 10GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.48 | 2Q-Congested | 10.46 | 10.97 | 19.95 | 64 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.76 | 2Q-Congested | 10.92 | 11.52 | 20.39 | 512 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.53 | 2Q-Congested | 10.46 | 11.07 | 19.02 | 1500 | 10GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 6.05 | 2Q-Congested | 10.01 | 12.11 | 17.33 | 64 | 100GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.16 | 2Q-Congested | 10.08 | 10.33 | 11.98 | 512 | 100GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.15 | 2Q-Congested | 10.10 | 10.31 | 11.82 | 1500 | 100GbE | Continuous | 2 |
| EVPN-VPWS (eCPRI) | 5.47 | 2Q-Congested | 10.40 | 10.95 | 12.96 | 64 | 100GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.36 | 2Q-Congested | 10.43 | 10.73 | 12.29 | 512 | 100GbE | Burst | 2 |
| EVPN-VPWS (eCPRI) | 5.35 | 2Q-Congested | 10.44 | 10.70 | 12.23 | 1500 | 100GbE | Burst | 2 |

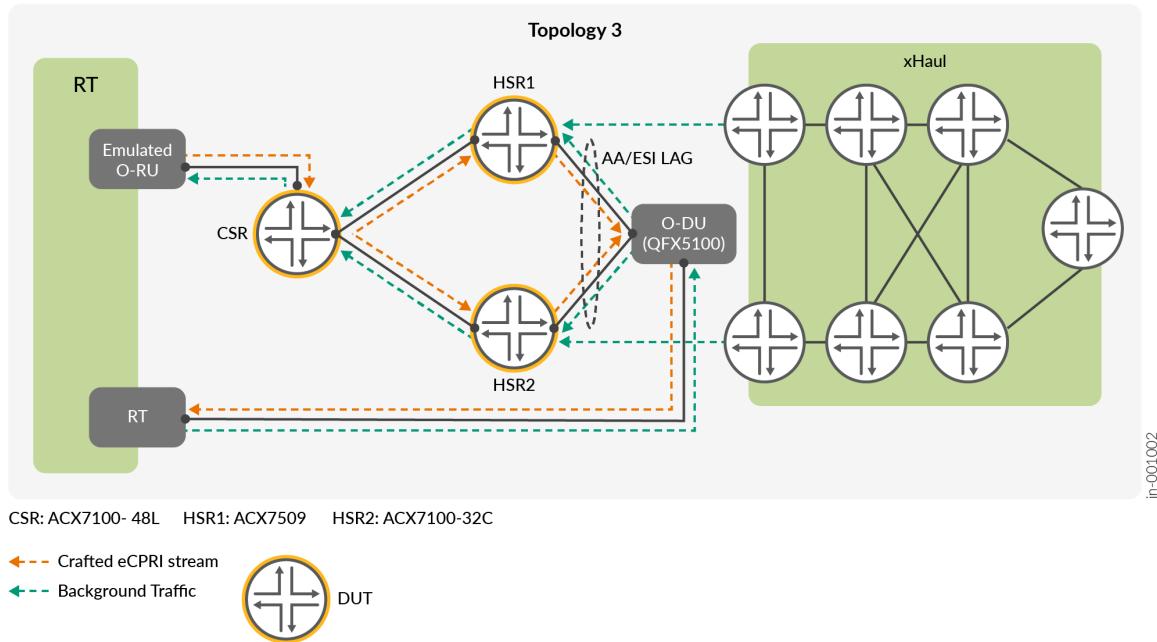
As shown, the average per-hop latency (2nd column) remains well below the overall goal of $\leq 10\mu\text{s}$ per device even during heavy network congestion, causing traffic loss in multiple priority queues (strict-high and high). In most cases, the per-hop latency results in less than $6\mu\text{s}$. As the port speed increases, the transmission delay associated with congestion typically decreases. As a result, slightly lower latency is seen. The low-latency queue is designed to preserve the strict latency budget. This is a critical function for 5G architectures.

Topology 3: Latency Performance with eCPRI over active-active EVPN-VPWS Multihoming

Topology 3 measures low-latency queue performance over active-active EVPN-VPWS Multihoming services in the fronthaul network segment between CSR and a pair of HSRs. ACX7024 is the CSR with ACX7100-32C and ACX7509 as HSRs, counting as a single hop as traffic is load-shared. In this scenario, the QFX5120 is the third device emulating the DU. This section measures LLQ performance in conditions of non-congested and heavy congestion scenarios using real eCPRI traffic streams. Background traffic is sent up to 99%-line rate or exceeding the configured bandwidth rate to the point of discarding frames.

For brevity, the information provided here is summarized with results for EVPN-VPWS 5G fronthaul eCPRI performance. For fronthaul, midhaul, and MBH workloads, see the test report of all data across multiple VPN services representing the traffic profiles explained in this document.

Figure 32: Topology 3 CSR to active-active HSR



Topology 3 provides the multihomed performance data in a 3-hop scenario, which includes QFX Series Switches as the DU connecting to the all-active ESI LAG. The test center emulates the RU.

The following table provides measurements of the latency for eCPRI traffic sent over topology 3, which includes three hops supporting active-active multihoming EVPN-VPWS services. Results are provided across different frame sizes, 10GbE and 100GbE port speeds, and continuous or burst traffic patterns. Latency is always measured in microseconds (μs). In all cases, the per-device average latency is measured less than the 10μs objective, with the majority ~3-4μs.

Table 36: Active-active Multihoming EVPN-VPWS Low-Latency without Congestion

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|---------------|------------------|------------------|------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 3.75 | Non-Congested | 11.09 | 11.24 | 11.79 | 64 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.98 | Non-Congested | 11.78 | 11.94 | 12.47 | 512 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 4.08 | Non-Congested | 12.09 | 12.23 | 13.54 | 1500 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.75 | Non-Congested | 11.09 | 11.25 | 11.77 | 64 | 10GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 3.98 | Non-Congested | 11.79 | 11.94 | 12.49 | 512 | 10GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 4.07 | Non-Congested | 11.80 | 12.22 | 13.56 | 1500 | 10GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 3.78 | Non-Congested | 10.95 | 11.35 | 11.66 | 64 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.73 | Non-Congested | 11.07 | 11.20 | 11.53 | 512 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.73 | Non-Congested | 11.11 | 11.18 | 11.63 | 1500 | 100GbE | Continuous | 3 |

Table 36: Active-active Multihoming EVPN-VPWS Low-Latency without Congestion (Continued)

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|---------------|------------------|------------------|------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 5.49 | Non-Congested | 64 | 10.99 | 11.39 | 64 | 100GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 5.56 | Non-Congested | 512 | 11.12 | 11.24 | 512 | 100GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 5.57 | Non-Congested | 1500 | 11.15 | 11.22 | 1500 | 100GbE | Burst | 3 |

As shown in the table, the average per-hop latency (2nd column) equates to ~3-4μs per device as traffic is load-shared and includes an additional low latency QFX Series Switches element as the 3rd hop DU. For the test validation, eCPRI traffic is generated from CSR (ACX7024) to active-active multihomed HSRs (ACX7100-32C and ACX7509) and traversing QFX5120 as the DU.

The following table describes the latency performance with congestion created on the strict-high queue and additional scenarios, including heavy congestion with traffic loss across multiple queues (strict-high and high).

Table 37: EVPN-VPWS Low-Latency with Heavy Congestion

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min (μs) Latency | Ave (μs) Latency | Max (μs) Latency | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|--------------|------------------|------------------|------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 4.24 | 1Q-Congested | 11.32 | 12.71 | 20.91 | 64 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 4.67 | 1Q-Congested | 12.44 | 14.01 | 21.78 | 512 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 5.12 | 1Q-Congested | 13.52 | 15.35 | 25.50 | 1500 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.80 | 1Q-Congested | 11.09 | 11.41 | 16.99 | 64 | 10GbE | Burst | 3 |

Table 37: EVPN-VPWS Low-Latency with Heavy Congestion (Continued)

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min Latency (μs) | Ave Latency (μs) | Max Latency (μs) | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|--------------|------------------|------------------|------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 4.04 | 1Q-Congested | 11.81 | 12.13 | 18.65 | 512 | 10GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 4.14 | 1Q-Congested | 12.09 | 12.41 | 18.08 | 1500 | 10GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 3.79 | 1Q-Congested | 10.83 | 11.38 | 12.51 | 64 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.73 | 1Q-Congested | 11.03 | 11.20 | 12.18 | 512 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.74 | 1Q-Congested | 11.06 | 11.21 | 12.22 | 1500 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.80 | 1Q-Congested | 11.03 | 11.41 | 12.65 | 64 | 100GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 3.75 | 1Q-Congested | 11.08 | 11.27 | 12.35 | 512 | 100GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 3.75 | 1Q-Congested | 11.11 | 11.26 | 11.29 | 1500 | 100GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 4.18 | 2Q-Congested | 11.33 | 12.54 | 25.68 | 64 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 4.57 | 2Q-Congested | 12.40 | 13.72 | 25.47 | 512 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 4.99 | 2Q-Congested | 13.51 | 14.98 | 27.73 | 1500 | 10GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 3.87 | 2Q-Congested | 11.08 | 11.62 | 20.63 | 64 | 10GbE | Burst | 3 |

Table 37: EVPN-VPWS Low-Latency with Heavy Congestion (Continued)

| Traffic Type | Per-Hop Ave (μs) | Scenario | Min Latency (μs) | Ave Latency (μs) | Max Latency (μs) | Frame Size | Port Speed | Traffic Pattern | Hops |
|-------------------|------------------|--------------|------------------|------------------|------------------|------------|------------|-----------------|------|
| EVPN-VPWS (eCPRI) | 4.14 | 2Q-Congested | 11.81 | 12.41 | 23.91 | 512 | 10GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 4.24 | 2Q-Congested | 12.12 | 12.71 | 20.88 | 1500 | 10GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 5.20 | 2Q-Congested | 11.08 | 15.61 | 17.93 | 64 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 5.13 | 2Q-Congested | 15.07 | 15.38 | 17.45 | 512 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 4.77 | 2Q-Congested | 11.07 | 14.32 | 17.39 | 1500 | 100GbE | Continuous | 3 |
| EVPN-VPWS (eCPRI) | 4.63 | 2Q-Congested | 11.10 | 13.91 | 18.05 | 64 | 100GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 3.78 | 2Q-Congested | 11.07 | 11.35 | 13.11 | 512 | 100GbE | Burst | 3 |
| EVPN-VPWS (eCPRI) | 4.25 | 2Q-Congested | 11.10 | 12.76 | 17.51 | 1500 | 100GbE | Burst | 3 |

As shown in the table, the average per-hop latency (2nd column) results in ~3-5μs per device even during heavy network congestion, causing traffic loss in multiple priority queues (strict-high and high). As the port speed increases, the associated transmission delay might decrease. The low-latency queue is designed to preserve the strict latency budget.

Buffer Occupancy Validation

As part of the validation, buffer utilization is verified throughout the test execution. For the complete details, see the full Test Report. The CoS buffer configuration outlined in the ["Scheduling" on page 49](#) section proposes the following buffer allocation per queue for the ACX7000 platform.

With the configuration implemented in the JVD, the following table displays the expected buffer allocation based on a 100GbE port with the 1250KB dedicated buffer.

Table 38: Buffer Allocation Estimation

| Queue | Rate | Buffer Size | Priority | Calculation | Buffer |
|-------|-------------------|-------------|-------------|----------------|--------|
| Q7 | 5% shaped rate | 5% | strict-high | $1250000*0.05$ | 62500 |
| Q6 | 40% shaped rate | 10% | low-latency | $1250000*0.1$ | 125000 |
| Q5 | 30% shaped rate | 20% | medium-high | $1250000*0.2$ | 250000 |
| Q4 | 40% transmit rate | 30% | low | $1250000*0.3$ | 375000 |
| Q3 | 5% shaped rate | 5% | high | $1250000*0.05$ | 62500 |
| Q2 | 30% transmit rate | 20% | low | $1250000*0.2$ | 250000 |
| Q1 | 20% transmit rate | 10% | low | $1250000*0.1$ | 125000 |
| Q0 | remainder | remainder | low | $1250000*0$ | *8192 |

** remaining buffer is 0%, so minimum buffer programmed*

The following outputs confirm that the buffer is allocated as expected. For more information related to buffer utilization during the test validation, see the Test Report.

```
jnx(/dev/pts/1)# sh cos voq buffer-profile ifd 1039
Rate Class Hardware Configuration
=====
  Ifd Index:1039  Unit:0  Port:1  Queue:0  Voq Index:40  Profile Id:15
  -----
  |          Tail Drop Config          |          WRED Drop Config
  |
  -----
  | colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold
  |
  -----
  | Green  |          8192 |      500000000 |          0 |          0 |          0
  |
  | Yellow |          8192 |      500000000 |          0 |          0 |          0
  |
  | Red    |          8192 |      500000000 |          0 |          0 |          0
  |
  -----
  =====
  Ifd Index:1039  Unit:0  Port:1  Queue:1  Voq Index:41  Profile Id:28
  -----
  |          Tail Drop Config          |          WRED Drop Config
  |
  -----
  | colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold
  |
  -----
  | Green  |          124928 |      500000000 |          0 |          0 |          0
  |
  | Yellow |          124928 |      500000000 |          0 |          0 |          0
  |
  | Red    |          124928 |      500000000 |          0 |          0 |          0
  |
  -----
  =====
  Ifd Index:1039  Unit:0  Port:1  Queue:2  Voq Index:42  Profile Id:29
  -----
  |          Tail Drop Config          |          WRED Drop Config
  |
  -----
  | colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold
  |
```

| | | | | | |
|--------|--------|-----------|---|---|---|
| Green | 249856 | 500000000 | 0 | 0 | 0 |
| Yellow | 249856 | 500000000 | 0 | 0 | 0 |
| Red | 249856 | 500000000 | 0 | 0 | 0 |

Ifd Index:1039 Unit:0 Port:1 Queue:3 Voq Index:43 Profile Id:14

| Tail Drop Config | | | WRED Drop Config | | |
|------------------|------------------|---------------|-------------------|---------------|---------------|
| colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold |
| Green | 62464 | 500000000 | 0 | 0 | 0 |
| Yellow | 62464 | 500000000 | 0 | 0 | 0 |
| Red | 62464 | 500000000 | 0 | 0 | 0 |

Ifd Index:1039 Unit:0 Port:1 Queue:4 Voq Index:44 Profile Id:30

| Tail Drop Config | | | WRED Drop Config | | |
|------------------|------------------|---------------|-------------------|---------------|---------------|
| colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold |
| Green | 374784 | 500000000 | 0 | 0 | 0 |
| Yellow | 374784 | 500000000 | 0 | 0 | 0 |
| Red | 374784 | 500000000 | 0 | 0 | 0 |

Ifd Index:1039 Unit:0 Port:1 Queue:5 Voq Index:45 Profile Id:29

| Tail Drop Config | | | WRED Drop Config | | |
|------------------|--|--|------------------|--|--|
| | | | | | |

| colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold |
|--|------------------|---------------|-------------------|---------------|---------------|
| <hr/> | | | | | |
| Green | 249856 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| Yellow | 249856 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| Red | 249856 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| <hr/> <hr/> | | | | | |
| Ifd Index:1039 Unit:0 Port:1 Queue:6 Voq Index:46 Profile Id:28 | | | | | |
| <hr/> | | | | | |
| Tail Drop Config | | | WRED Drop Config | | |
| <hr/> | | | | | |
| colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold |
| <hr/> | | | | | |
| Green | 124928 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| Yellow | 124928 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| Red | 124928 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| <hr/> <hr/> | | | | | |
| Ifd Index:1039 Unit:0 Port:1 Queue:7 Voq Index:47 Profile Id:14 | | | | | |
| <hr/> | | | | | |
| Tail Drop Config | | | WRED Drop Config | | |
| <hr/> | | | | | |
| colour | Dedicated Buffer | Shared Buffer | Dynamic Threshold | Min Threshold | Max Threshold |
| <hr/> | | | | | |
| Green | 62464 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| Yellow | 62464 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |
| Red | 62464 | 500000000 | 0 | 0 | 0 |
| <hr/> | | | | | |

Host Traffic Classification

For details on the results related to assigning host traffic to a forwarding class for this profile, see the Test Report.

Traffic Characteristics

For all details on traffic flows and types of streams created in the execution of this JVD, see the Test Report.

Recommendations

The 5G RAN imposes stringent demands on the MBH network infrastructure, necessitating a substantial increase in the number of nodes, performance, and feature richness. This elevated complexity poses new challenges that require innovative solutions. Juniper Networks presents an end-to-end 5G xHaul network infrastructure solution designed to support an existing 4G MBH while seamlessly evolving into 5G network infrastructure over the same physical network. This approach facilitates smooth transitions for operators from 4G to 5G without compromising their current services. The transition can be gradual, enabling an introduction of necessary changes that cater to new requirements such as enhanced bandwidth capabilities, sub-millisecond latency, improved synchronization for mission-critical applications, network slicing for customized services, and scalability to manage an increased number of connected devices. Juniper's solution ensures a seamless and efficient evolution of the MBH network infrastructure, meeting the evolving demands of 5G networks.

The 5G solution architecture demands strict QoS to deliver differentiated services and maintain low latency for critical traffic. Key benefits of the LLQ JVD include:

- **Custom Traffic Prioritization:** Tailored priorities for enhanced mobile broadband (eMBB), ultra-reliable low-latency communication (URLLC), and massive machine-type communications (mMTC).
- **Optimized Parameters:** Ensures appropriate bandwidth, latency, and reliability for various traffic classes.
- **Resource Efficiency:** Allocates bandwidth and prioritizes traffic efficiently based on application needs.
- **Enhanced User Experience:** Provides consistent and reliable quality for streaming, gaming, and remote work.

- **Low Latency Queuing (LLQ):** Reduces transmission delay for real-time applications by prioritizing crucial traffic.
- **High Reliability:** Minimizes data loss and transmission errors for stable communication in critical services.
- **Urgent Traffic Prioritization:** Ensures immediate attention and resources for emergency communications and public safety applications.

These features collectively enable 5G networks to meet modern application demands with robust, efficient, and flexible performance.

The ACX7024, ACX7100-32C, and ACX7509 platforms are presented as DUTs to support deterministic and effective QoS for supporting differentiated and delay-sensitive workloads. These platforms offer an advanced feature-set, enhanced performance, and support a common software architecture and feature roadmap. In addition, the ACX7024, ACX7100, and ACX7509 platforms have successfully completed MEF 3.0 certification, which carries rigid QoS requirements.

As part of this validation, multiple topologies are utilized to provide meaningful data in different circumstances. CoS functional operations are reported without issues, supporting the ability to classify traffic based on BA, Fixed, or Multifield Classification. In the validated design, an eight-queue 5G CoS model is created with four priority queues configured using shaping rate: Signaling, LLQ, Realtime, and Control. These queues carried the most critical traffic types, with the majority for the fronthaul network segment. Another four WFQs are used to allow proportional and dynamic bandwidth allocation based on the transmit rate. These low queues established High, Medium, Low, and Best-Effort service-type categories. Explicit instructions are created for all forwarding classes on the type of traffic to be serviced. Under all scenarios, the expected bandwidth allocation and queue priorities are honored. Codepoint preservation is achieved across all featured VLAN manipulation sequences covered in this JVD.

Through extensive validation, the performance of the ACX7000 platforms exceeded the objectives for latency preservation goals in both non-congested and congested scenarios, delivering a comprehensive multi-priority CoS solution while preserving a latency budget of $\leq 10\mu\text{s}$. During network congestion causing heavy packet loss, the ACX platforms averaged a transit latency of 4-7 microseconds (μs) while delivering services across all eight queues. The device architecture is designed to support delay sensitive applications, with the low-latency queue given further preferential treatment. These capabilities are critical functions for 5G solution architectures.

This JVD is based on 5G reference architectures. However, the CoS modeling, best practices, and performance results are applicable to many implementations. For more details on the Juniper JVD solution tested, see the Test Report or contact your Juniper Networks representative.

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